

# Clustering

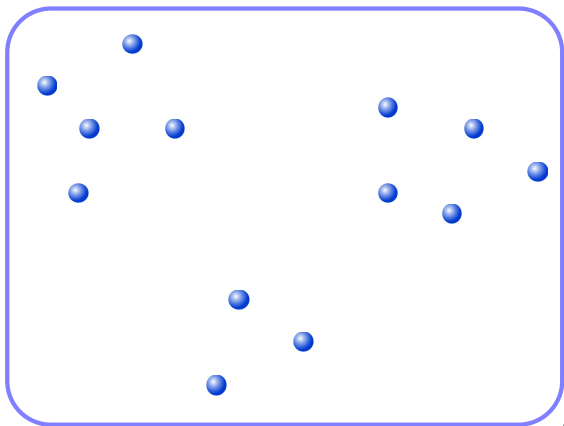
## How Bad Is The $k$ -Means++ Method?

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Maastricht University  
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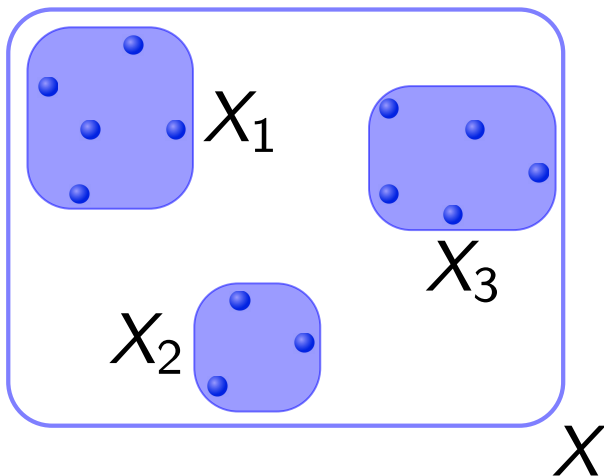
July 15, 2010

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Objective: Find clustering  $(X_1, \dots, X_k)$  and centers  $c_1, \dots, c_k$  with minimal potential  $\Phi(X)$



# The Challenge

$k$ -means is  $\mathcal{NP}$ -hard,

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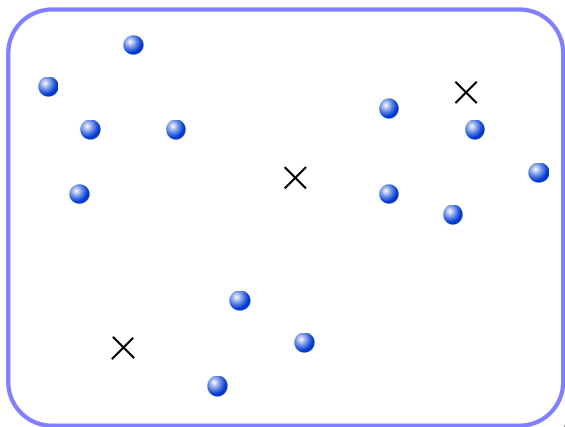
$\implies$  Approximation algorithms, heuristics

## Lloyd's Algorithm ( $k$ -Means Method, $k$ -Means)

Observations:

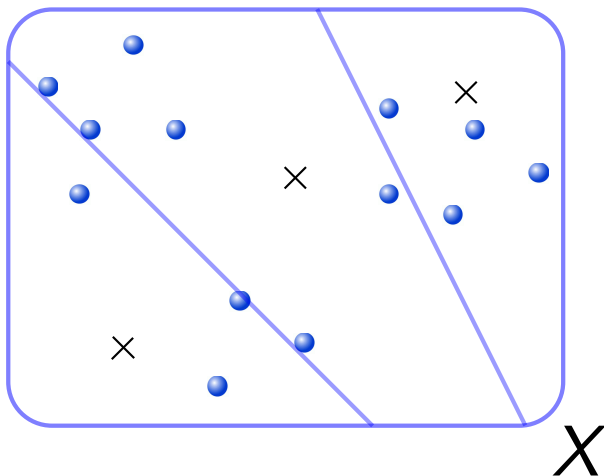
- The optimal centers  $c_i$  for given clusters  $X_i$  are their centers of mass
- The optimal clusters  $X_i$  for given centers  $c_i$  are the points nearest to  $c_i$

# Lloyd's Algorithm

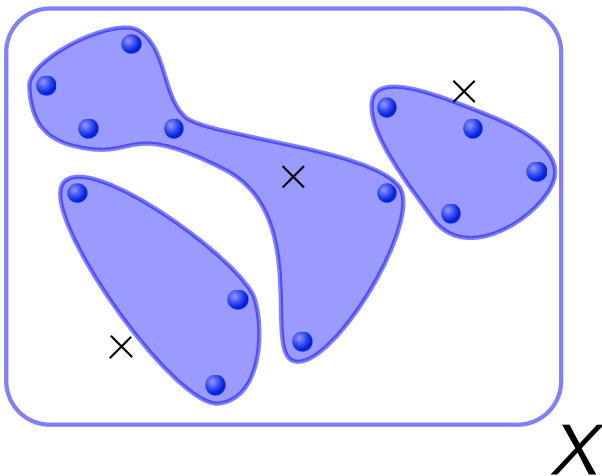


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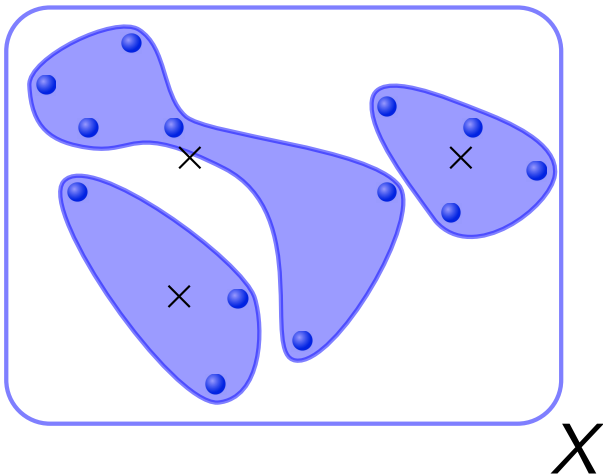
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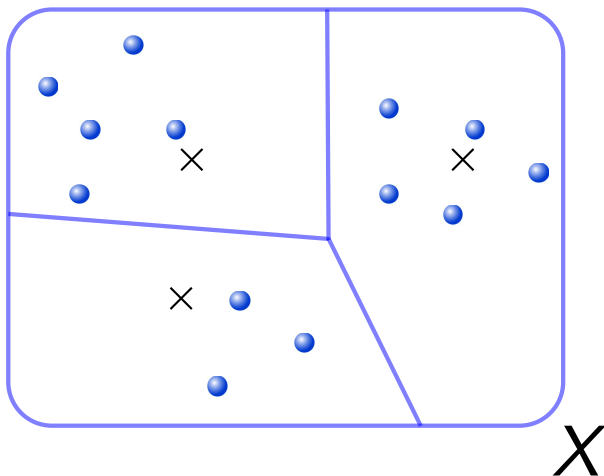
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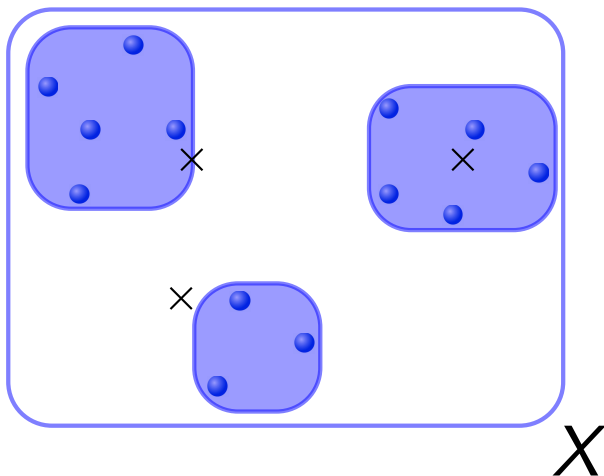


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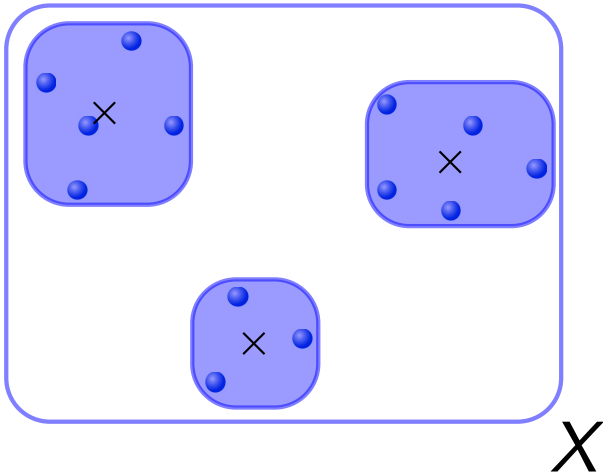




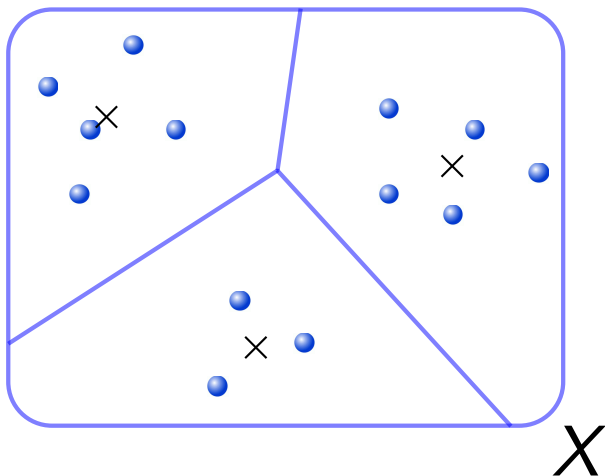
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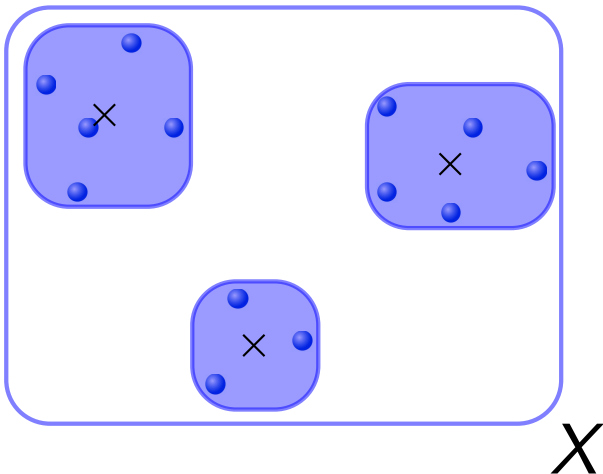
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# Advantages And Disadvantages

## Advantages

Simple to implement

Practitioners









Theoreticians











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Simple to implement		
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<b>Disadvantages</b>		
Exponential worst-case time		
Requires good initialization		



## Tackling The Disadvantages

- Polynomial time in the framework of *smoothed analysis*
- Approximation guarantee with  $k$ -means++ seeding technique

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- 1 Choose the first center uniformly at random
- 2 Choose point  $x \in X$  with probability  $\frac{D^2(x)}{\Phi(X)}$  as next center

$$\left( D^2(x) = \min_{c_i} \|x - c_i\|^2 \right)$$

## Asymptotic Bounds

Theorem (Arthur and Vassilvitskii, 2007)

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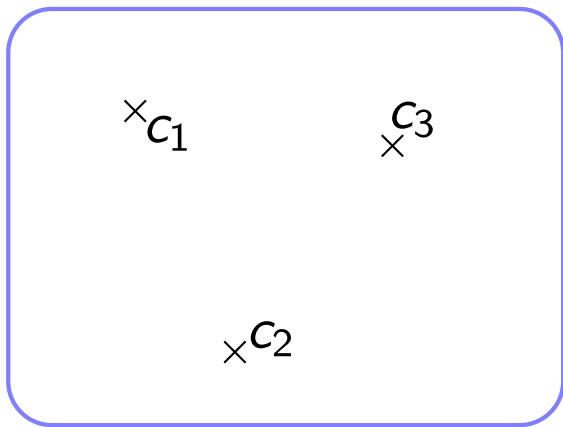
Observation

There is a family of instances on which the expected approximation ratio of  $k$ -means++ is  $\Omega(\log k)$ .

## Open Question

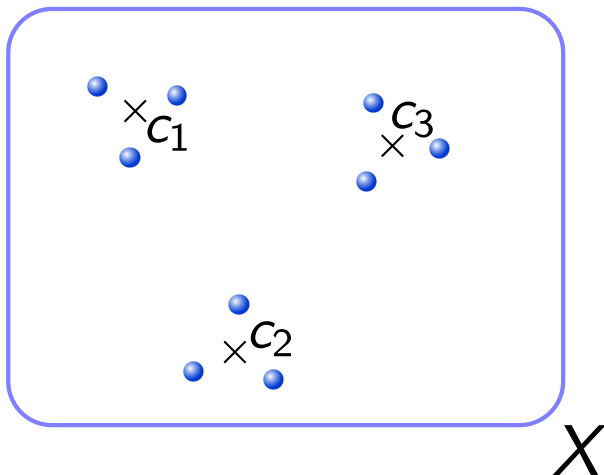
Does  $k$ -means++ yield an  $O(1)$ -approximation with constant probability?

## The Instance

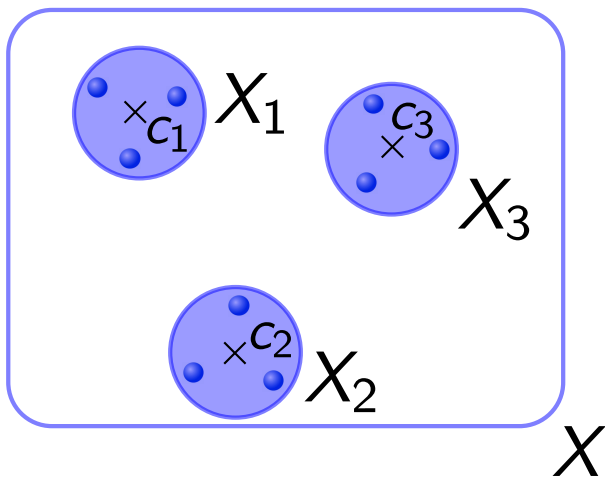




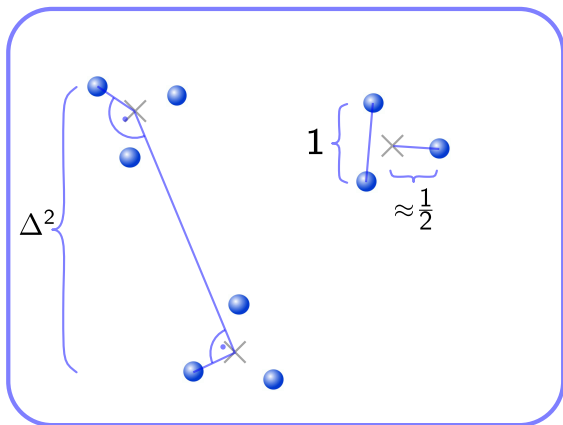
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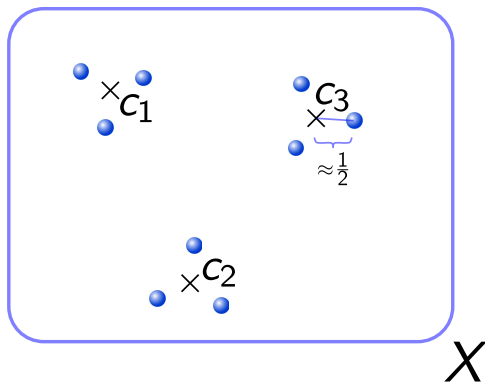
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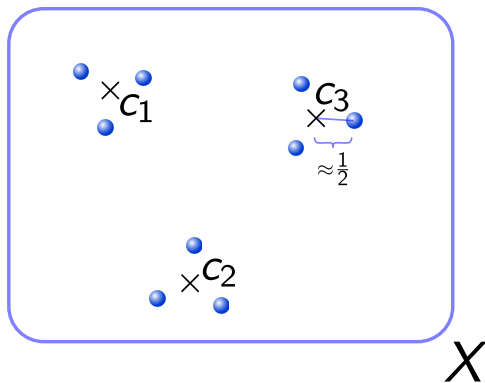
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## Optimal Clustering $C^*$

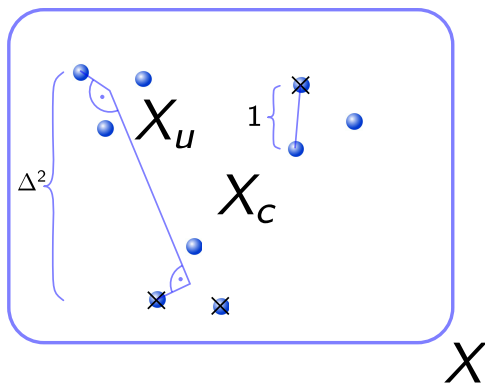


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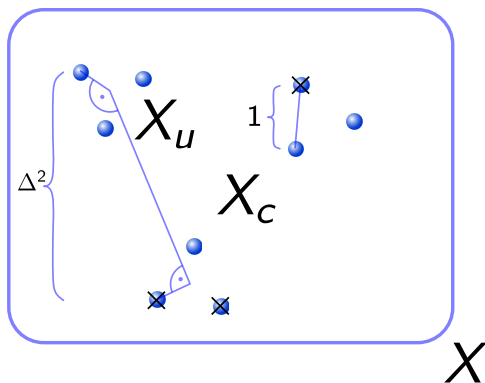


$$\Phi^*(X) \approx k \cdot k \cdot \frac{1}{2} = \frac{k^2}{2}$$

## Discrete Clustering With $s$ Covered Sets $X_i$

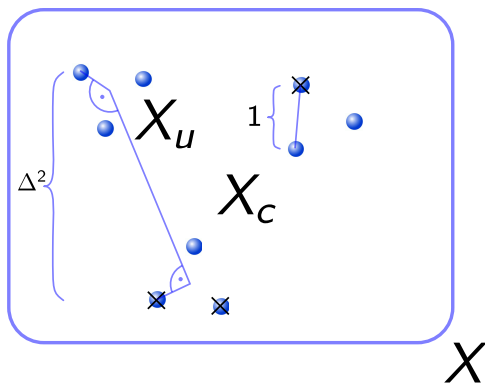


## Discrete Clustering With $s$ Covered Sets $X_i$



$$\Phi(X) = \Phi(X_c) + \Phi(X_u) \approx s \cdot k \cdot 1 + (k - s) \cdot k \cdot \Delta^2$$

## Discrete Clustering With $s$ Covered Sets $X_i$



Covering probability:  $\frac{\Phi(X_u)}{\Phi(X)} \approx \frac{1}{1 + \frac{1}{(k-s) \cdot \Delta^2}} =: p_s$



## How Many Sets To Cover?

In the end:

$$r \geq \frac{\Phi(X)}{\Phi^*(X)} \geq \frac{\Phi(X_u)}{\Phi^*(X)} \approx 2\Delta^2 \cdot \left(1 - \frac{s}{k}\right) \quad (r - \text{approximation factor})$$

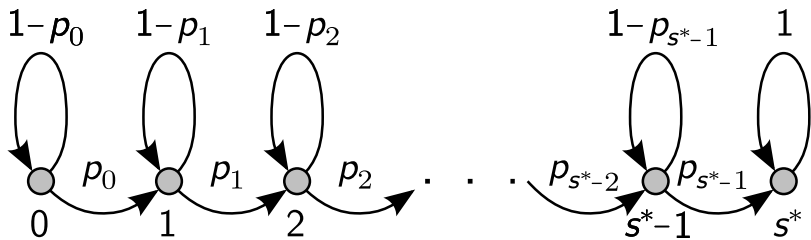
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$$\implies s \gtrsim k \cdot \left(1 - \frac{r}{2\Delta^2}\right) =: s^*$$

# Markov Chain



## Expected Number Of Steps $X$

$$\mathbf{E}[X] = \sum_{s=0}^{s^*-1} \frac{1}{p_s} \gtrsim k + \frac{k}{\Delta^2} \cdot \left( \ln \frac{\Delta^2}{r} - \frac{r}{2} \right)$$

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If  $r \in o(\log k)$ , then  $\Pr[X \leq k]$  is exponentially small in  $k$   
 (Hoeffding Inequality + workaround)

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- 1 Do  $k$ -means++ and  $k$ -means together yield an  $O(1)$ -approximation with constant probability?
- 2 Can we slightly modify  $k$ -means++ to guarantee better bounds?

