

Coverability in 2-VASS with One Unary Counter

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Fun-Road-Trip Checklist

- ✓ at least one friend, and
- ✓ not have negative money!



Friends: 4
Money: 1000kr



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GOAL

Friends: ≥ 5
Money: $\geq 100kr$



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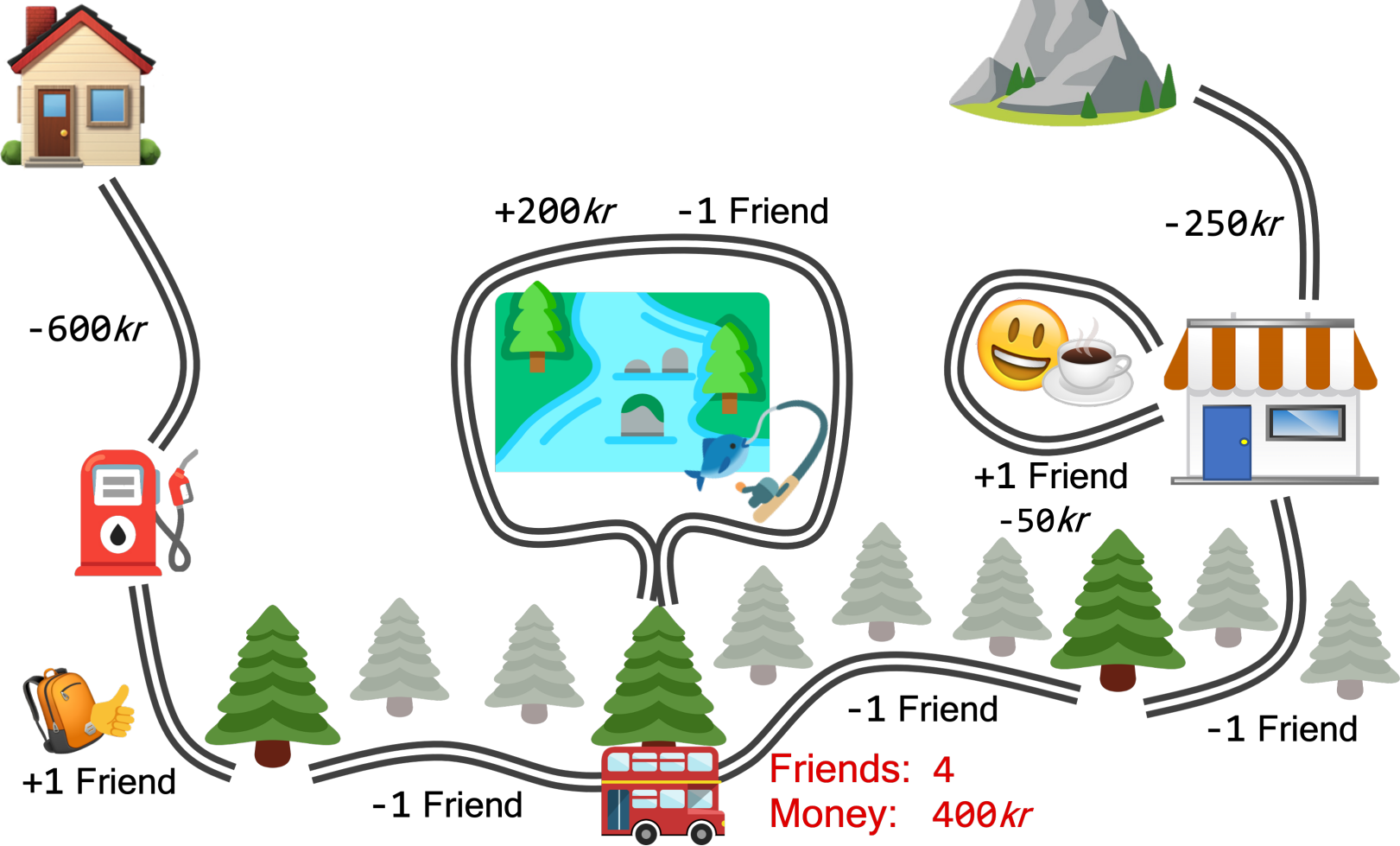
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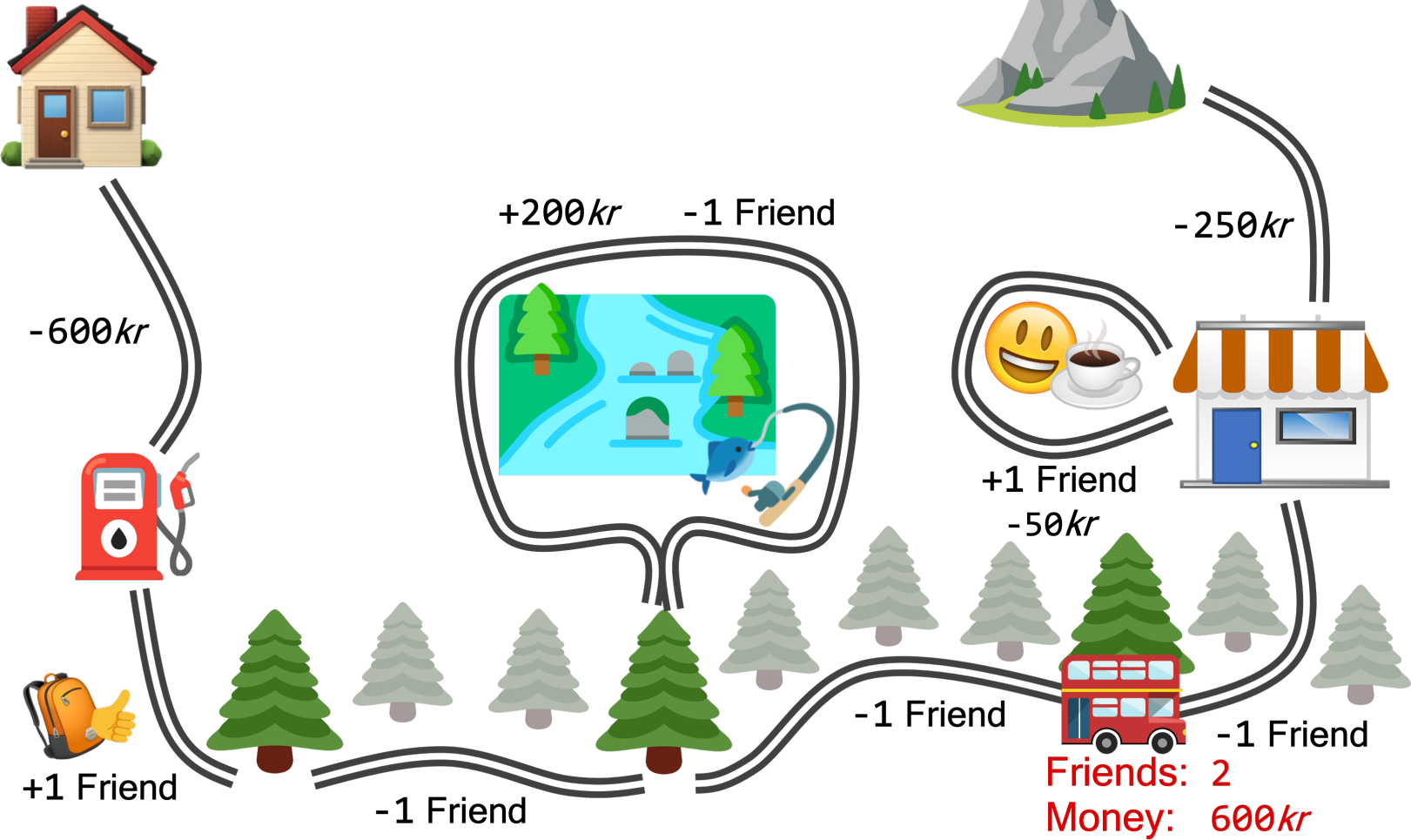
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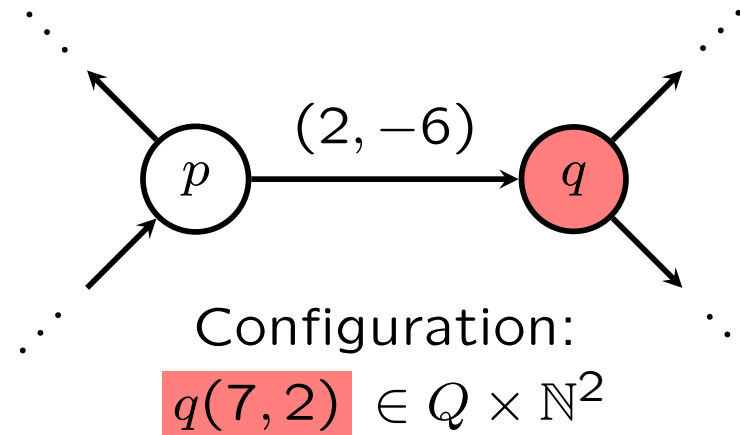
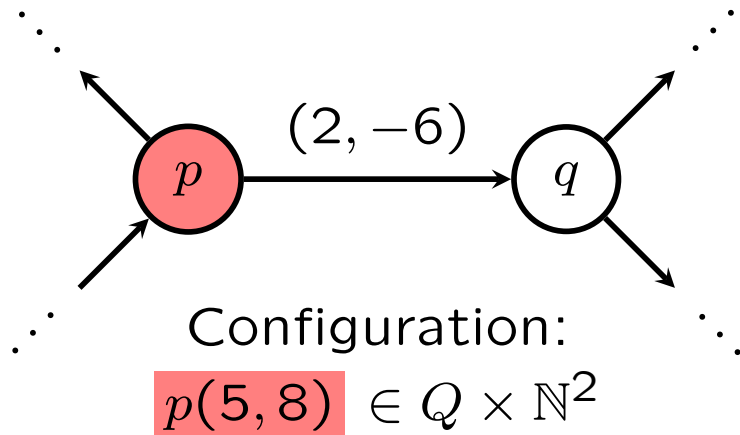
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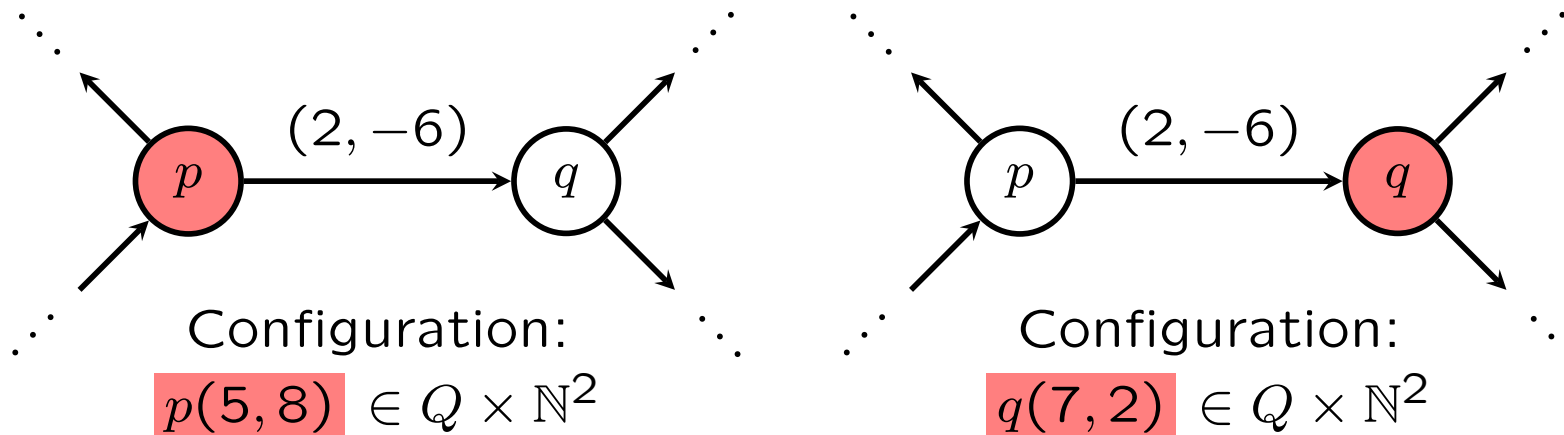
INTRODUCTION

Vector Addition Systems with States (2-VASS)



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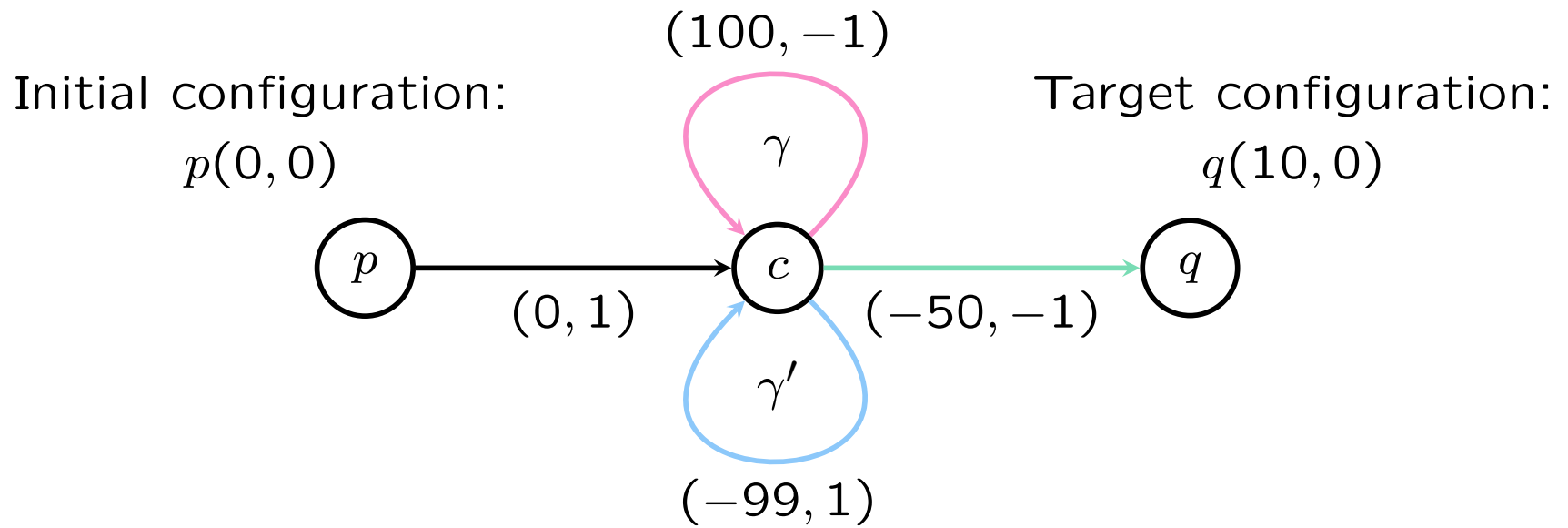
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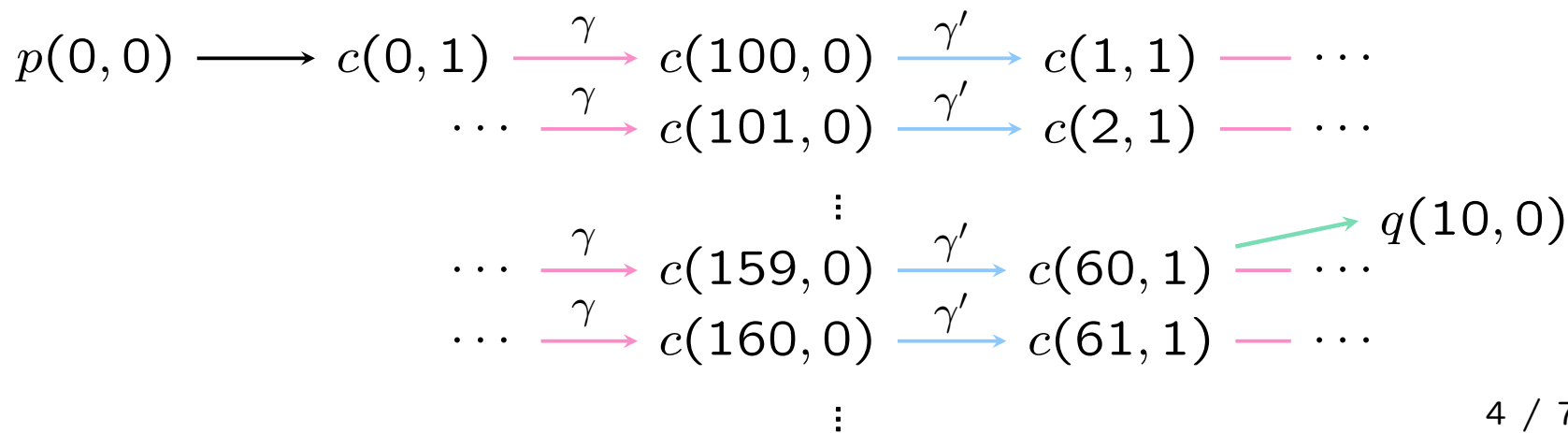
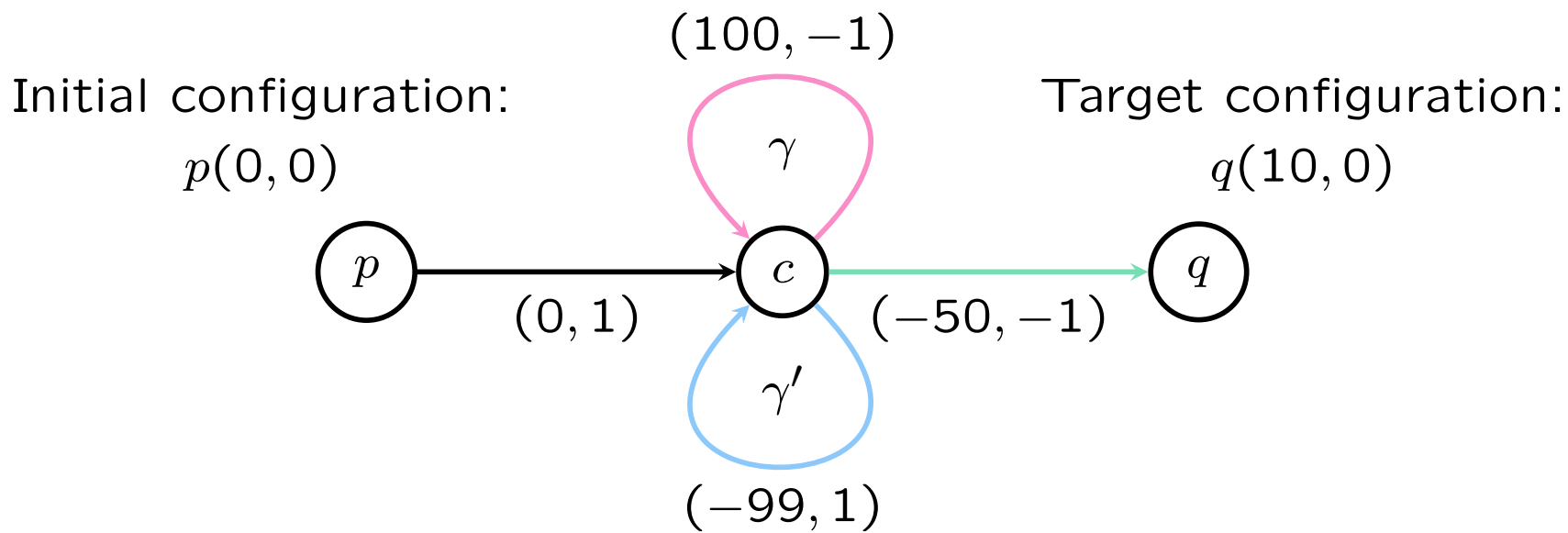
Reachability does there exist a *run* in V from $p(\mathbf{u})$ to $q(\mathbf{v})$?

Coverability does there exist a *run* in V from $p(\mathbf{u})$ to $q(\mathbf{v}')$ for some $\mathbf{v}' \geq \mathbf{v}$?

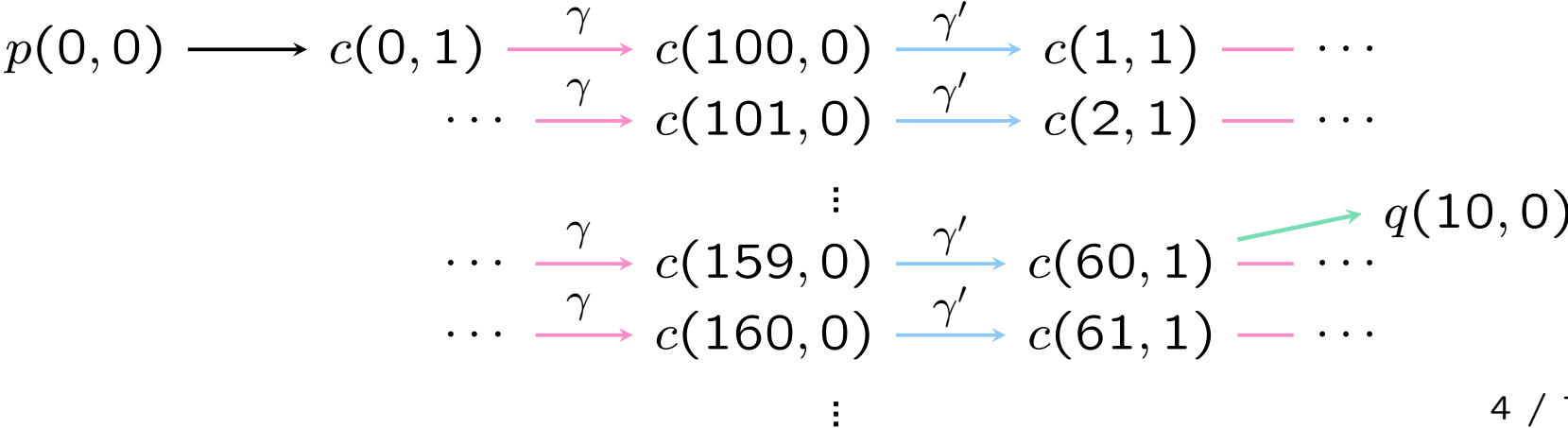
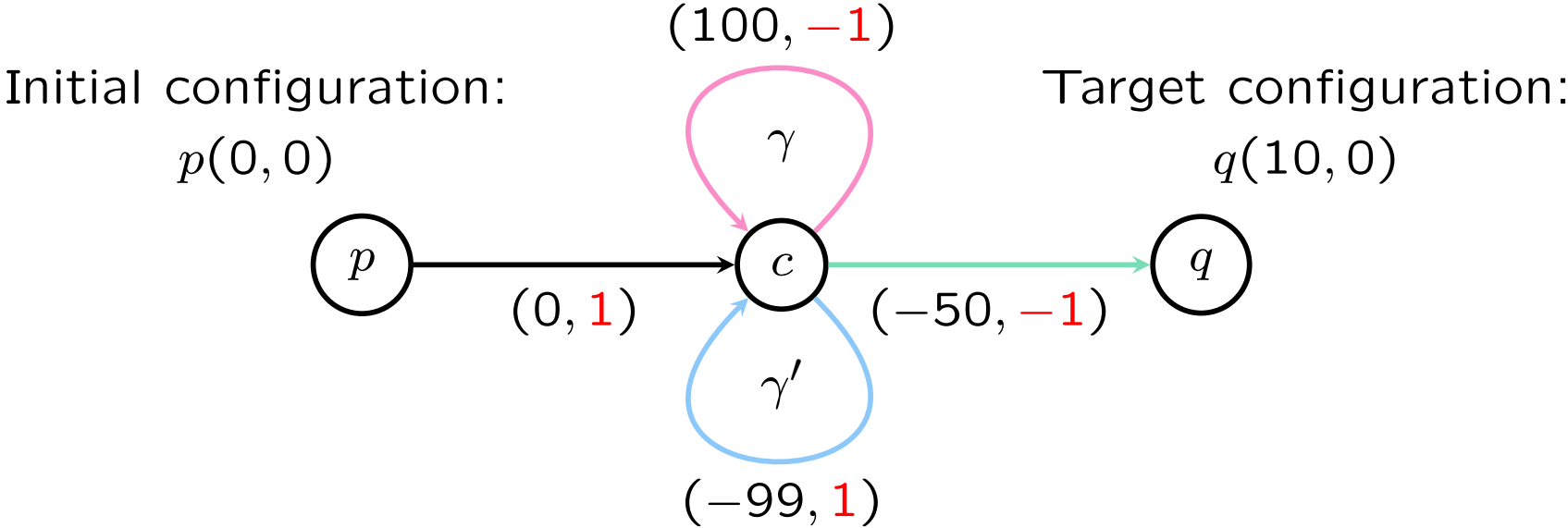
COVERABILITY EXAMPLE



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Coverability in 2-VASS with One Unary Counter



CONTRIBUTION

Binary **2-VASS coverability** is PSPACE-complete.

[Blondin, Finkel, Göller, Haase, and McKenzie '15]

Unary **2-VASS coverability** is NL-complete.

[Rackoff '78]

OUR CONTRIBUTION

Binary **2-VASS coverability** is PSPACE-complete.

[Blondin, Finkel, Göller, Haase, and McKenzie '15]

Theorem:

Coverability in 2-VASS with one unary counter is in NP.

Unary **2-VASS coverability** is NL-complete.

[Rackoff '78]

RESULTS

Theorem: Given a **2-VASS with one unary counter** V and suppose there exists a run in V from $p(\mathbf{u})$ to $q(\mathbf{v})$. Then there exists a compressed linear form path of *polynomial size* inducing a run from $p(\mathbf{u})$ to $q(\mathbf{v}')$ for some $\mathbf{v}' \geq \mathbf{v}$.

RESULTS

Theorem: Given a **2-VASS with one unary counter** V and suppose there exists a run in V from $p(\mathbf{u})$ to $q(\mathbf{v})$. Then there exists a compressed linear form path of *polynomial size* inducing a run from $p(\mathbf{u})$ to $q(\mathbf{v}')$ for some $\mathbf{v}' \geq \mathbf{v}$.

\Rightarrow Coverability in 2-VASS with one unary counter is in NP.

... just guess and check compressed linear form paths.

CONCLUSION

Coverability in 2-VASS with one unary counter is in NP.

Unfortunately, we lack a matching NP-hard lower bound.

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Conjecture: **coverability** in P.

Future Work: is **reachability** also in NP?

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THANK YOU!

Presented by Henry Sinclair-Banks