Classes of languages generated by the Kleene star of a word

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Classes of rational languages

- closed under some operations
- generated by a basis

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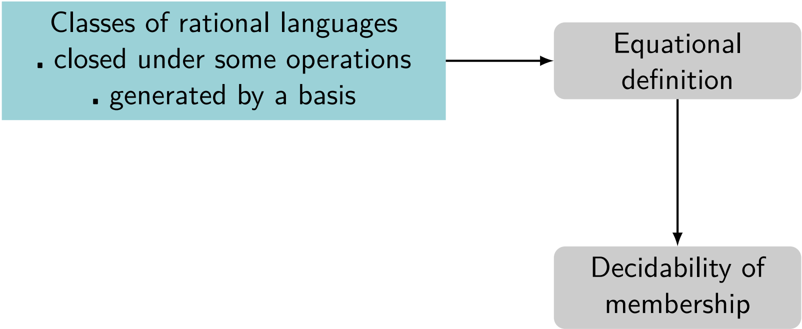
→ Equational definition



Classes of rational languages
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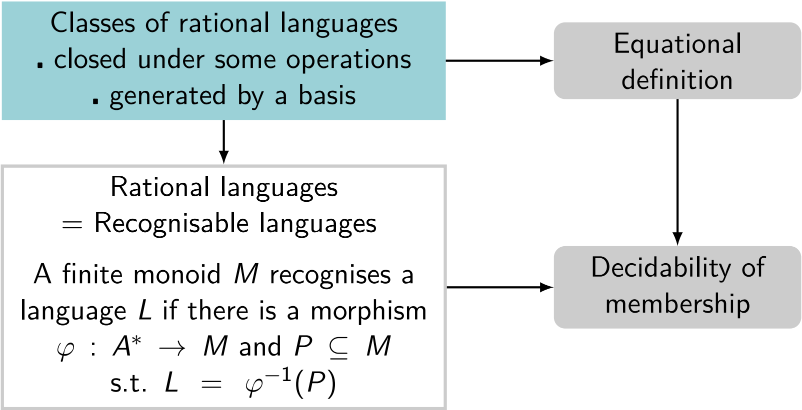
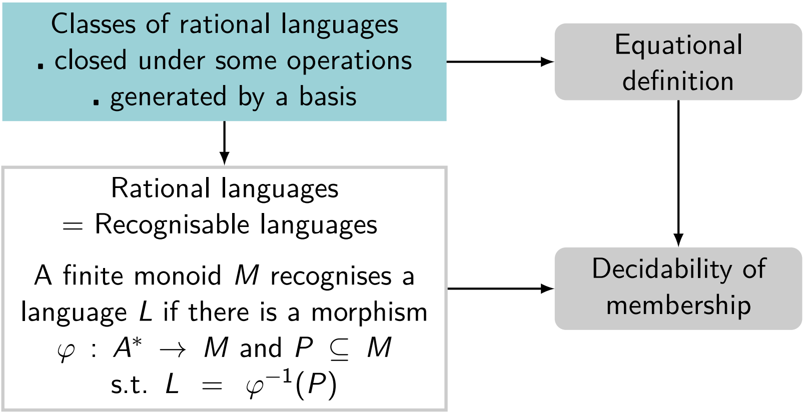
Equational definition
Decidability of membership


Classes of rational languages
. closed under some operations
. generated by a basis
Equational definition
Rational languages
= Recognisable languages
A finite monoid M recognises a language L if there is a morphism $\varphi : A^* \to M$ and $P \subseteq M$ s.t. $L = \varphi^{-1}(P)$
Decidability of membership



Classes of languages generated by $\{u^* \mid u \in A^*\}$

$$u^* = \{u^n \mid n \in \mathbb{N}\}$$

Profinite monoid

A monoid M separates u and v if:

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Definition

Profinite monoid $\widehat{A^*}$:
completion of A^* with respect to the distance d.

- Compact
- $\varphi : A^* \to M$ can be uniquely extended to a continuous morphism $\widehat{\varphi} : \widehat{A^*} \to M$

V.I.P. words (very important profinite words)

Idempotent power

$$u^\omega = \lim_{n\to\infty} u^{n!}$$

- For all morphisms $\varphi : A^* \to M$ (finite monoid): $\widehat{\varphi}(u^\omega)$ is the idempotent power of $\widehat{\varphi}(u)$ in M.

Zero (Reilly-Zhang 2000, Almeida-Volkov 2003)

$|A| \geqslant 2$

$u_0, u_1, \ldots$ an enumeration of the words of A^*

$v_0 = u_0, \quad v_{n+1} = (v_n u_{n+1} v_n)^{(n+1)!}$

$$\rho_A = \lim_{n\to\infty} v_n$$

- For all morphisms $\varphi : A^* \to M$ (finite monoid): if M has a zero then $\widehat{\varphi}(\rho_A) = 0$.

Equations

Definition

Given two profinite words u, v, a rational language L satisfies

$$u \to v$$

if $u \in \overline{L}$ implies $v \in \overline{L}$

$a \in A$

Equation $ab \to aba$

$\{L \subseteq A^* \mid ab \notin L\} \cup \{L \subseteq A^* \mid ab, aba \in L\}$

Equations

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Given two profinite words u, v, a rational language L satisfies

$$u \leftrightarrow v$$

if $u \in \overline{L}$ if and only if $v \in \overline{L}$

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Equation $ab \leftrightarrow aba$

$\{L \subseteq A^* \mid ab, aba \notin L\} \cup \{L \subseteq A^* \mid ab, aba \in L\}$

Equations

Definition

Given two profinite words u, v, a rational language L satisfies

$$u \leqslant v$$

if for all $w, w' \in A^*, wuw' \in \bar{L}$ implies $wvw' \in \bar{L}$

$a \in A$

Equation $ab \leqslant aba$

$\{L \subseteq A^* \mid \text{for all } w, w', \text{ if } wabw' \in L \text{ then } wabaw' \in L\}$

Equations

Definition

Given two profinite words u, v, a rational language L satisfies

$$u = v$$

if for all $w, w' \in A^*, wuw' \in \overline{L}$ if and only if $wvw' \in \overline{L}$

$a \in A$

Equation $ab = aba$

$\{L \subseteq A^* \mid \text{for all } w, w',\ wabw' \in L \text{ iff } wabaw' \in L\}$

Equations

Theorem [Gehrke, Grigorieff, Pin 2008]

Classes of rational languages

- Lattice (union, intersection): $\rightarrow$
- Boolean algebra (lattice, complement): $\leftrightarrow$
- Lattice closed under quotient: $\leqslant$
- Boolean algebra closed under quotient: $=$

quotient : $u^{-1}Lv^{-1} = \{w \mid uwv \in L\}$

Existence of a zero: $\{\rho_A u = u\rho_A = \rho_A \mid u \in A^*\}$

Why u^*?

- Non trivial applications of the equational theory of rational languages.

- Characterise the variety of languages generated by the languages u^*. [Restivo]

- (long-term perspective) Boolean algebra generated by the languages F^* for F a finite set of words. [key result for the generalised star-height problem]

Equations for u^*

$$P_u = \bigcup_{p \text{ prefix of } u} u^* p \quad \text{and} \quad S_u = \bigcup_{s \text{ suffix of } u} s u^*$$

$$x^\omega y^\omega = 0 \text{ for } x, y \in A^* \text{ such that } xy \neq yx \tag{E_1}$$

$$x^\omega y = 0 \text{ for } x, y \in A^* \text{ such that } y \notin P_x \tag{E_2}$$

$$y x^\omega = 0 \text{ for } x, y \in A^* \text{ such that } y \notin S_x \tag{E_3}$$

$$x^\omega \leqslant 1 \text{ for } x \in A^* \tag{E_4}$$

$$x^\ell \leftrightarrow x^{\omega+\ell} \text{ for } x \in A^*,\ \ell > 0 \tag{E_5}$$

$$x \to x^\ell \text{ for } x \in A^*,\ \ell > 0 \tag{E_6}$$

$$x^\alpha \leftrightarrow x^\beta \text{ for all } (\alpha, \beta) \in \Gamma \tag{E_7}$$

DECIDABLE

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DECIDABLE Lattice closed under quotients

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DECIDABLE Boolean algebra closed under quotients

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DECIDABLE Boolean algebra

The Boolean algebra

An example:

$$(a^2)^* - (a^6)^* = (a^6)^* a^2 \cup (a^6)^* a^4$$

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Equivalence relation over the integers

$r \equiv_m s$ if and only if $\gcd(r, m) = \gcd(s, m)$
$(u^m)^* u^r \subseteq L$ if and only if $(u^m)^* u^s \subseteq L$

The Boolean algebra

An example:

$$(a^2)^* - (a^6)^* = (a^6)^* a^2 \cup (a^6)^* a^4$$

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Translation in terms of profinite numbers

Γ is the set of pairs of profinite numbers $(dz^{\mathcal{P}}, dpz^{\mathcal{P}})$ such that:

- $\mathcal{P}$ is a cofinite sequence of prime numbers $\{p_1, p_2, \ldots\}$
- $z^{\mathcal{P}} = \lim_n (p_1 p_2 \ldots p_n)^{n!}$
- $p \in \mathcal{P}$
- if q divides d then $q \notin \mathcal{P}$

$$x^\alpha \leftrightarrow x^\beta \text{ for all } (\alpha, \beta) \in \Gamma \qquad (E_7)$$

Conclusion and further questions

- Equational description of classes of rational languages generated by the languages u^*
- Decidability

- Variety of languages generated by the languages u^*
- Classes generated by F^* for F a finite set of words