

Last week:

- Review of Problem Set 1

This week:

- Bit vs. phase oracles
- "Multi-output" Deutsch-Jozsa

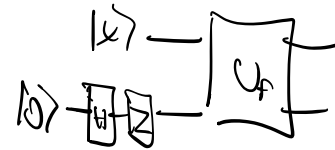
For a function $f: \{0, 1\}^n \rightarrow \{0, 1\}^n$, denote by U_f the unitary $|x, y\rangle \mapsto |x, (y \oplus f(x))\rangle$.

- (a) Suppose $n = 1$. Show how to build a circuit that computes the unitary $|x\rangle \mapsto (-1)^{f(x)} |x\rangle$ (known as the phase oracle). You may use Z gates, ancilla qubits initialized to $|0\rangle$, and two U_f gates. You must ensure that any ancilla qubits return to the state $|0\rangle$ so that they can be safely discarded. Prove that your circuit is correct.

$$U_f |x\rangle |y\rangle = |x\rangle |f(x) \oplus y\rangle$$

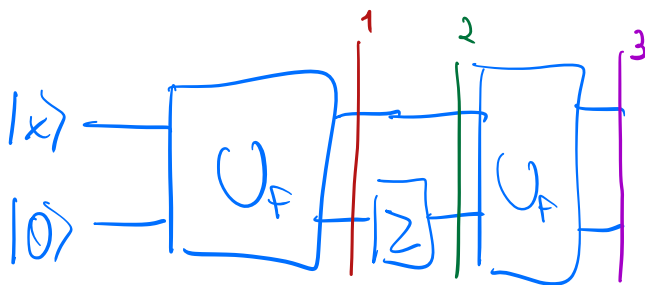
$$\Rightarrow |x\rangle |0\rangle \mapsto (-1)^{f(x)} |x\rangle |0\rangle$$

$$f: \{0, 1\} \rightarrow \{0, 1\}$$



$f: \{0, 1\}^n \rightarrow \{0, 1\}^n$ $\left\{ \begin{array}{l} \text{Condition 1: } f(x) = y \\ \text{Condition 2: } f \text{ is a bijection} \end{array} \right.$

1 query to U_f
 $\Omega(n)$ queries to f



$$\begin{cases} |x\rangle |0\rangle & \text{if } f(x)=0 \\ -|x\rangle |1\rangle & \text{if } f(x)=1 \end{cases}$$

$$f(x) \oplus f(x) = 0$$

$$|x\rangle |0\rangle \xrightarrow{1} |x\rangle |f(x)\rangle \xrightarrow{2} (-1)^{f(x)} |x\rangle |f(x)\rangle \xrightarrow{3} (-1)^{f(x)} |x\rangle |0\rangle$$

(b) Suppose now (and for the remaining parts of this question) that $n = 2$. The gate S maps $|0\rangle \mapsto |0\rangle$ and $|1\rangle \mapsto i|1\rangle$. Show that $S^2 = Z$.

$$S|0\rangle = |0\rangle$$

$$S|1\rangle = i|1\rangle$$

$$S = \begin{pmatrix} 1 & 0 \\ 0 & i \end{pmatrix}$$

$$S^2 = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} = Z$$

$$S^2 = Z$$

$$\boxed{S} \boxed{S} \equiv \boxed{Z}$$

$$S^2|0\rangle = |0\rangle = Z|0\rangle$$

$$S^2|1\rangle = S \cdot (i|1\rangle)$$

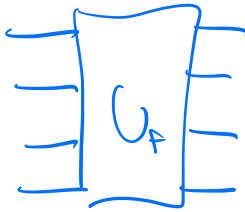
$$= i^2|1\rangle$$

$$= -|1\rangle = Z|1\rangle$$

(c) Show how to build a circuit that computes the unitary that maps $|x\rangle \mapsto i^{2f(x)_1+f(x)_2} |x\rangle$, where $f(x)_1, f(x)_2$ are the first and second bits of $f(x)$, respectively. You may use S gates, ancilla qubits initialised to $|0\rangle$, and **two** U_f gates. You must ensure that any ancilla qubits return to the state $|0\rangle$ so that they can be safely discarded. Prove that your circuit is correct.

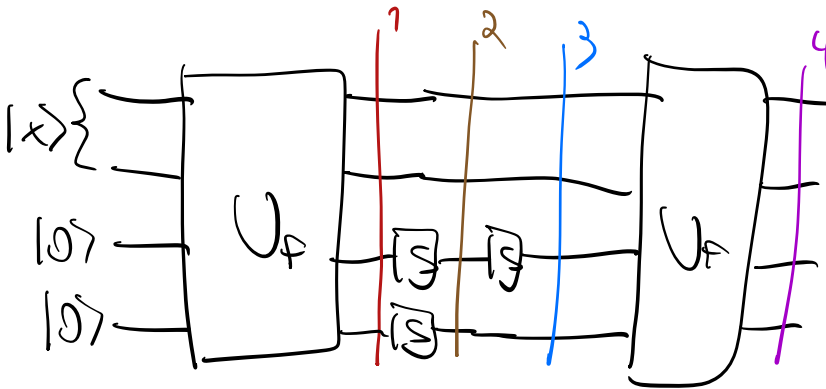
$$F: \{0,1\}^2 \rightarrow \{0,1\}^2$$

$$U_f \rightsquigarrow |x\rangle |00\rangle \xrightarrow{i^{2f(x)_1+f(x)_2}} |x\rangle |00\rangle$$



$$U_f |x_1\rangle |x_2\rangle |y_1\rangle |y_2\rangle$$

$$= |x_1\rangle |x_2\rangle |f(x)_1 \oplus y_1\rangle |f(x)_2 \oplus y_2\rangle$$



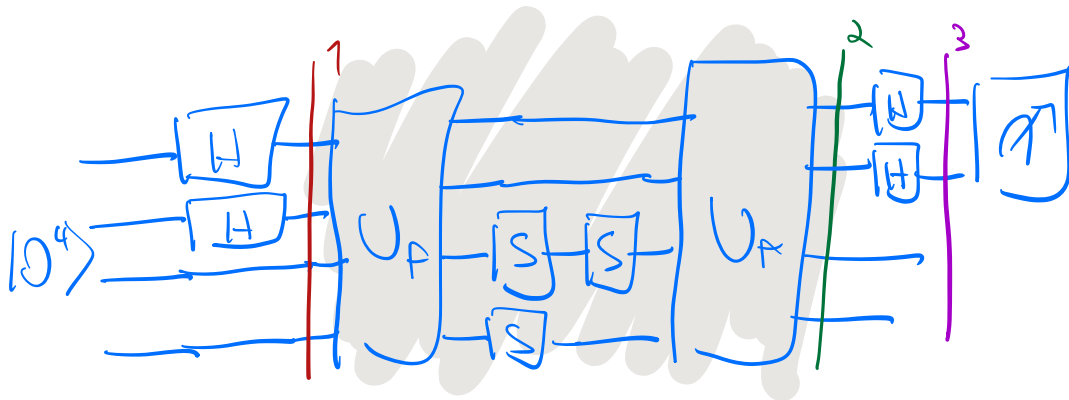
$$|x\rangle |00\rangle \xrightarrow{i^{2f(x)_1+f(x)_2}} |x\rangle |00\rangle$$

$f(x)_1$
 (-1)

(d) Design a circuit that determines whether f is constant or one-to-one. You may use:

- any number of qubits initialized to $|0\rangle$,
- Hadamard (H) gates,
- S gates,
- measurements in the computational basis, and
- **two** U_f gates.

Prove that your circuit is correct.



$$|0000\rangle \xrightarrow{1} \frac{1}{2} \sum_{x \in \{0,1\}^2} |x\rangle |00\rangle \xrightarrow{2} \frac{1}{2} \sum_x i^{2f(x)+f(x)} |x\rangle |00\rangle$$

$$H^{\otimes n} |x\rangle = \frac{1}{\sqrt{2^n}} \sum_y (-1)^{x \cdot y} |y\rangle$$

$$H^{\otimes n} |0^n\rangle = \frac{1}{\sqrt{2^n}} \sum_y |y\rangle$$

$$\frac{1}{4} \sum_x \sum_y (-1)^{x \cdot y + 2f(x)+f(x)} |y\rangle |00\rangle$$

$$\left(\frac{1}{4} \sum_x i^{2f(x)+f(x)} \right)^2 \begin{cases} 1 & \text{if } f \text{ is constant} \\ 0 & \text{if } f \text{ is 1-to-1} \end{cases}$$

$$i^0 + i^1 + i^2 + i^3 = 0$$

$ab \mapsto a^2b$

$00 \mapsto 0$

$01 \mapsto 1$

$10 \mapsto 2$

$11 \mapsto 3$

