Distinct Sampling on Streaming Data with Near-Duplicates

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joint work with Qin Zhang (IUB)



Workshop on Data Summarisation, Mar/2018

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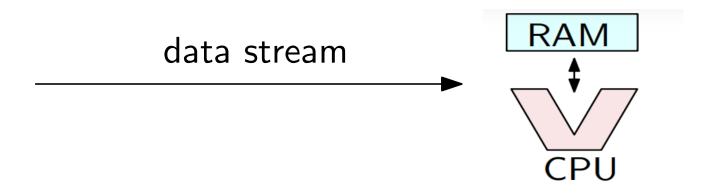


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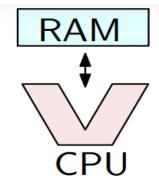


- "data summarization"
- "summarization of data"
- "the summarization of data"
- "summarization data"

queries of the same meaning sent to Google

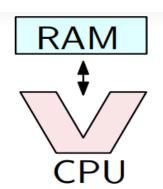


data stream



- high-speed online data
- want space/time efficient algorithms

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sliding window: only consider recent w items.

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So we could not apply existing streaming algorithms directly need some new ideas

Robust ℓ_0 -sampling

- data: data points in \mathbb{R}^d
- ℓ_0 -sampling: each distinct element is sampled with prob. $\frac{1}{F_0}$

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$$F_0 = 6$$







Formally ...

• $S \subset \mathbb{R}^d$ is (α, β) -sparse: either $d(u, v) \leq \alpha$ or $d(u, v) > \beta$ for all $u, v \in S$.

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- when $\beta/\alpha > 2$,

$$G(v) = \{ u \in S \mid d(u, v) \le \alpha \}$$

forms a group of v

- ullet the dataset S is well-shaped
- a natural partition exists for a well-shaped dataset
- F_0 is the number of groups

Our goal:

- ullet S is well-shaped, fed as a data stream
- $\mathcal{G} = \{G_1, G_2, \dots, G_{F_0}\}$ is the natural partition
- Goal: outputs a point u such that,

$$\forall i \in [F_0], \Pr[u \in G_i] = 1/F_0$$

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our algorithm also work with general datasets in $\mathbb{R}^{O(1)}$ (discuss later)

Basic idea

• over the data stream, mark one representative point for each group.

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---X--0--0--X---0--X----0---X-----

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Question: Can we identify (not necessarily store) the first arrived point of each group space-efficiently?

Unfortunately, $\Omega(F_0)$ space required to identify the representative points in noisy dataset

Solution: sample in advance

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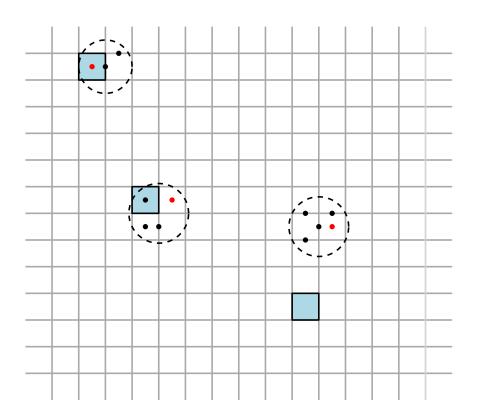
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- how to sample in advance?
 - place a random grid (side length $\frac{\alpha}{2}$) in \mathbb{R}^2 , sample cells before we see the data stream

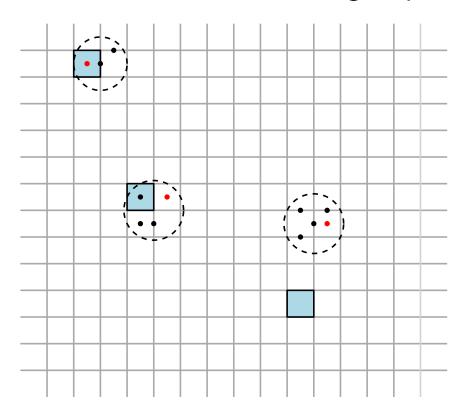
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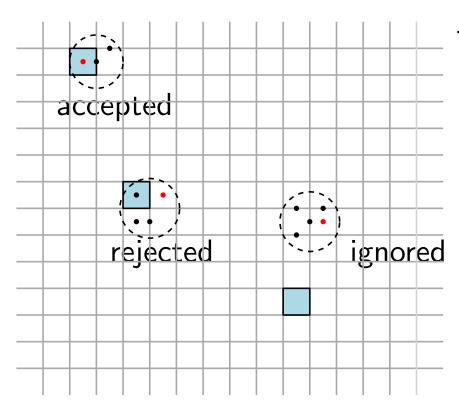
- how to sample in advance?
 - place a random grid (side length $\frac{\alpha}{2}$) in \mathbb{R}^2 , sample cells before we see the data stream
- how to decide the sample rate?
 - decrease when see more groups



- blue cells: sampled cell
- red points: first arrived point of its group



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three types of groups:

- accepted: first arrived point falls into a sampled cell
- ignored: no point falls into a sampled cell
- rejected: has point falling into a sampled cell, but not the first arrived point

How to maintain accepted groups?

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keep all first points of accepted groups

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keep it!

•

discard?

- keep all first points of accepted groups
 - keep it! discard?

If we discard the first arrived point of a rejected group...

is this the first arrived point?

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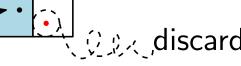
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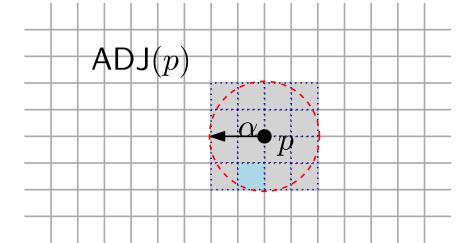
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if p is not in a sampled cell and $\mathsf{ADJ}(p)$ has cell sampled...

keep p!

- ullet $S^{\rm acc}$: first arrived points of accepted groups
- $\bullet \ \ S^{\rm rej} = \{ \text{first point } p \not \in S^{\rm acc} \ \text{and } \mathsf{ADJ}(p) \ \text{has sampled cell} \}$

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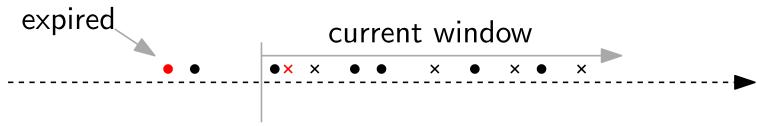
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how to decide the sample rate?

- if $|S^{\rm acc}| > \kappa \log m$, re-sample each sampled cell with prob. $\frac{1}{2}$
- ullet roughly, $S^{
 m acc}$ will drop half of its size
- $S^{\rm acc}$ is not empty w.h.p.
- space usage $O(\log m)$

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current window

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in S^{acc} , maintain pairs, (u,p) the latest point in G(u) the representative point

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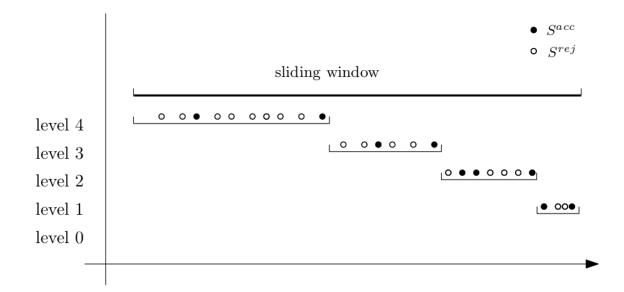
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- Level i samples cells with prob. $\frac{1}{2^i}$
- discard expired groups

- a level may sample $> \kappa \cdot \log m$
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 - Level i+1 now may have $> \kappa \cdot \log m$
 - ullet do the re-sampling again in Level i+1
 - this process may cascade to the top level
 - each level actually samples from a disjoint subwindow

the last point must be in $S^{\rm acc}$. Invariant during the split/merge process

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How to generate a sample at the point of query?

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How to generate a sample at the point of query?

- ullet Level T is the top non-empty level
- all groups in Level T-i is re-sampled with prob. $\frac{1}{2^i}$
- union all sampled groups
- then return a sample uniformly at random

Theorems

For well-shaped datasets in $\mathbb{R}^{O(1)}$

- infinite window: use $O(\log m)$ space, $O(\log m)$ processing time
- sliding window: use $O(\log m \log w)$ space, $O(\log m \log w)$ amortized processing time

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applying dimension reduction reduces d to $O(\log m)$

General datasets ...

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consistent with the well-shaped case

our algorithms in well-shaped dataset can achieve this goal in $\mathbb{R}^{O(1)}$

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• Other formulations for robust ℓ_0 -sampling on general data sets?

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How to extend to other metric spaces?

Questions?

Thank you!