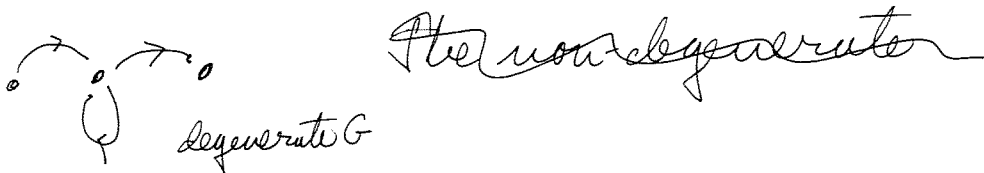


# Entropy of a topological Markov chain.

Let  $G$  be a directed graph. We want to show that the entropy of the shift  $\sigma: G_{\infty} \rightarrow G_{\infty}$  is given by the <sup>log of</sup> growth rate of the number of words of length  $n$ .

For this discussion I will make the technical hypothesis that each node of vertex of  $G$  has an entering edge and an exiting edge.

We will call such a  $G$  nondegenerate.



Let  $W_n(G)$  be the set of paths of length  $n$  in  $G$  also call words. If  $G$  is non-degenerate then every  $w \in W_n(G)$  is contained in ~~an~~ a bi-infinite word and in particular  $G_{\infty} \neq \emptyset$ .

If  $l$  is the # of vertices in  $G$  then  $A = A_G = (a_{ij})$  is the  $l \times l$  matrix where  $a_{ij}$  is the number of edges from  $v_i$  to  $v_j$ .

If we write  $A^n = (a_{ij}^n)$  then  $a_{ij}^n$  is the number of paths of length  $n$  from  $v_i$  to  $v_j$ .

Thus  $\#W_n(G) = \sum_{i,j} a_{ij}^n$ .

Define a norm on the vector space of  $l \times l$  matrices

$$\|M\| = \sum_{i,j} |m_{ij}|.$$

so

$$\#W_n(G) = \|A^n\|.$$

The spectral radius theorem says that

$$\lim_{n \rightarrow \infty} \|A^n\|^{1/n} = \rho(A)$$

where  $\rho(A) = \max_i |\lambda_i|$  and  $\lambda_i \in \mathbb{C}$  ranges over the complex eigenvalues of  $A$  is the solutions of  $AV = \lambda V$  for  $\lambda \in \mathbb{C}$  and  $V \in \mathbb{C}^l$ .

~~One consequence~~ This result is independent of our choice of norm.

Now taking logs of both sides we have

$$\log \left( \lim_{k \rightarrow \infty} \|A^k\|^{1/k} \right) = \log(P(A))$$

$$\lim_{k \rightarrow \infty} \frac{\log \|A^k\|}{k} = \log(P(A)).$$

Thus 
$$\lim_{k \rightarrow \infty} \frac{\log \#W_k(G)}{k} = \log(P(A_G)).$$

Proposition. Let  $G$  be a non-degenerate graph then and  $\sigma: G \rightarrow G$  the shift then:

$$h(\sigma) = \log(P(A_G)).$$

Remarks. We can consider 1-sided or 2-sided shifts and entropy is the same in either case. We will give the proof in the 2-sided case.

Proof. Let  $\lambda > 1$ . Consider the metric  $d_\lambda$  on  $G_{\mathbb{Z}}$ .

$$d_\lambda(x, x') = \max \frac{\varepsilon(x_i, x'_i)}{\lambda^{|i|}}$$

Pick  $k \geq 1$ . We wish to calculate  $h(\sigma, \varepsilon)$  where

$$\varepsilon = 1/\lambda^k.$$

Let  $\alpha_{-k} \dots \alpha_k \in W_{2k+1}(G)$ .

(4)

Let  $C_{-\alpha_k \dots \alpha_k} = \{ \alpha' \in G_{\infty} : \alpha'_j = \alpha_j \text{ for } -k \leq j \leq k \}$ .

Then  $B(\alpha, \lambda, \lambda^{-k}) = C_{\alpha_{-k} \dots \alpha_k}$   $\lambda^{-k-1} < \varepsilon < \lambda^{-k}$   
 $B(\alpha, \lambda^k)$   $B(\alpha, \varepsilon) = C_{-\alpha_k \dots \alpha_k}$

Case

Let  $d_{in}(\alpha, \alpha') = \max_{0 \leq l \leq n} d_x(\sigma^l(\alpha), \sigma^l(\alpha'))$

$\begin{matrix} -k & 0 & -k & k+n \\ [ & & ] & ] \end{matrix}$

$B(\alpha, n, \frac{\lambda^{k+n}}{\varepsilon})$  is the cylinder set  $C_{-\alpha_k \dots \alpha_{k+n}}$   $\lambda^{-k-1} < \varepsilon < \lambda^{-k}$

As  $C_{-\alpha_k \dots \alpha_{k+n}}$  ranges over the words in  $W_{2k+n+1}$  we get an ~~exhaustive~~ covering of  $G_{\infty}$  by disjoint  $\varepsilon$  balls so  $S(n, \varepsilon) \leq \# W_{2k+n+1}$ .  
 On the other hand we can find an  $\varepsilon$  separated set by choosing one point in each ball so  $N(n, \varepsilon) \geq \# W_{2k+n+1}$ .

Plus

$$\begin{aligned}
 h(\sigma, \varepsilon) &= \limsup_{n \rightarrow \infty} \frac{\log(S(n, \varepsilon))}{n} \geq \limsup_{n \rightarrow \infty} \frac{\log(\#W_{2k+1}^{(n)})}{n} \\
 &= \limsup_{n \rightarrow \infty} \frac{\log(\#W_{2k+1}^{(n)})}{2k+1} \cdot \lim_{n \rightarrow \infty} \frac{2k+1}{n} \\
 &= \log(p(A)).
 \end{aligned}$$

$$\begin{aligned}
 h(\sigma, \varepsilon) &= \limsup_{n \rightarrow \infty} \frac{\log(N(n, \varepsilon))}{n} \leq \limsup_{n \rightarrow \infty} \frac{\log(\#W_{2k+1}^{(n)})}{n} \\
 &= \log(p(A)).
 \end{aligned}$$

$$\text{So } \lim_{\varepsilon \rightarrow 0} h(\sigma) = \lim_{\varepsilon \rightarrow 0} h(\sigma, \varepsilon) = \log(p(A)).$$

Remark. We have seen that the growth rate of the number of fixed points of  $\sigma^n$  is  $T_0(A^n)$ .

$$T_0(A^n) = \lambda_1^n + \dots + \lambda_k^n.$$

If  $A$  is aperiodic then the Perron-Frobenius theorem says that  $A$  has a unique eigenvalue <sup>of max</sup> of maximum modulus and that this eigenvalue is real and positive.

In this case

$$\begin{aligned} \lim_{n \rightarrow \infty} \frac{1}{n} \log(\# \text{fixed pts of } \sigma^n) &= \lim_{n \rightarrow \infty} \frac{1}{n} \log \lambda_{\text{PF}}^n + \lambda_2^n + \dots \\ &= \lim_{n \rightarrow \infty} \frac{1}{n} \log \lambda_{\text{PF}}^n \\ &= \log \lambda_{\text{PF}} \\ &= P(A) \\ &= h(\sigma). \end{aligned}$$

Consider the linear hyperbolic linear toral automorphism  $f: \mathbb{T}^2 \rightarrow \mathbb{T}^2$  given by  $\begin{pmatrix} 2 & 1 \\ 1 & 1 \end{pmatrix}$ . We have a semi-conjugacy from a top. Markov chain  $(\sigma: \mathbb{S}^1 \rightarrow \mathbb{S}^1)$  to  $f: \mathbb{T}^2$ . A priori this means that  $h(\sigma) \geq h(f)$   
 $\log(\phi^2)$

by our result on semi-conjugacies.

In fact we can use our analysis of the sizes of rectangles to show the equality of entropy.

~~Let  $A$  be a  $2 \times 2$  integral matrix so~~

Let  $A$  be a hyperbolic integral matrix with  $\det A = \pm 1$ . Let  $f_A: \mathbb{T}^2 \rightarrow \mathbb{T}^2$  be the corresponding linear hyperbolic automorphism diffeomorphism.

Let  $\lambda$  be the larger eigenvalue of  $A$ .

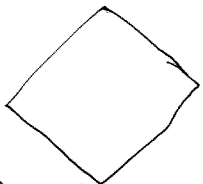
Thm.  $h(f_A) = \log(\lambda)$  where  $|\lambda_1| > 1 > |\lambda_2|$  are the eigenvalues of  $A$ .

Thm.  $h(f_A) = \log(\lambda)$ .

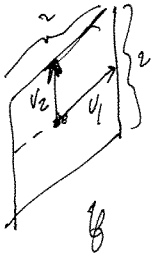
Proof. I want to construct a convenient norm on  $\mathbb{R}^2$ . Let  $v_1$  and  $v_2$  be unit length eigenvectors for  $A$  associated to  $\lambda_1$  and  $\lambda_2$ .

Define  $\|v\|_{\max} = |av_1 + bv_2|_{\max} = \max\{|a|, |b|\}$ .

The unit ball for this norm is a tipped parallelogram or diamond (rhombus?).



or



Let  $\pi: \mathbb{R}^2 \rightarrow \mathbb{R}^2 / \mathbb{Z}^2 = \mathbb{T}^2$ .

For  $\tilde{p}, \tilde{q} \in \mathbb{R}^2$

Define  $\tilde{d}_{\max}(\tilde{p}, \tilde{q}) = \|p - q\|_{\max}$ .

Define  $d_{\max}$  a metric on  $\mathbb{T}^2$  by

$d_{\max}(p, q) = \min_{\tilde{p}, \tilde{q} \in \mathbb{R}^2} \tilde{d}_{\max}(\tilde{p}, \tilde{q})$  where  
 $\pi(\tilde{p}) = p$   
 $\pi(\tilde{q}) = q$ .



The map  $\pi$  is a local <sup>isometry</sup> homeomorphism. ~~For~~

There is a  $C > 0$  so that  $\exists \tilde{C}(\tilde{p}, \tilde{q}) < C \Rightarrow d_{\max}(p, q) = d_{\max}(\tilde{p}, \tilde{q})$ .

$C \approx 1/4$ .

Now define  $d_n^{(A^e)} = \max_{0 \leq k \leq n} d_{\max}(f^k(p), f^k(q))$ ,  $\tilde{d}_n(\tilde{p}, \tilde{q}) = \max_{0 \leq k \leq n} d_{\max}(f^k(\tilde{p}), f^k(\tilde{q}))$

Claim that  $\pi$  is also a local isometry wrt.  $\tilde{d}_n$  and  $d_n$ .

If  $\tilde{d}_n(\tilde{p}, \tilde{q}) < C \Rightarrow \max_{0 \leq k \leq n} (d_{\max}(f^k(\tilde{p}), f^k(\tilde{q}))) < C$

$$\Rightarrow \tilde{d}_{\max}(A^e(\tilde{p}), A^e(\tilde{q})) < C$$

$$\Rightarrow d_{\max}(A^e(p), A^e(q)) < C$$

$$\Rightarrow \tilde{d}_{\max}(A^e(\tilde{p}), A^e(\tilde{q})) = d_{\max}(f_{A^e}^e(p), f_{A^e}^e(q))$$

$$\Rightarrow d_n(p, q) = \tilde{d}_n(\tilde{p}, \tilde{q}).$$

Now we see the advantage of dealing with  $\tilde{d}_n$  is that we can write it down more explicitly:

Let  $v = av_1 + bv_2$  then  $A^e(v) = a\lambda_1^e v_1 + b\lambda_2^e v_2$

$$\|A^e(v)\| = \max(|a\lambda_1^e|, |b\lambda_2^e|)$$

$$\|A^e(v)\|_n = \max\{|a|, |b|, |\lambda_1^e a|, |\lambda_2^e b|$$

$$\dots, |\lambda_1^{n-1} a|, |\lambda_2^{n-1} b|\}$$

$$|A^n(v)|_n = \max\{|r_i^{n-1} a|, |b|\}$$

Define  $L_n(av_1 + bv_2) = (r_i^{n-1} a, b)$   
 $= (r_i^{n-1} a v_1, b v_2)$ .

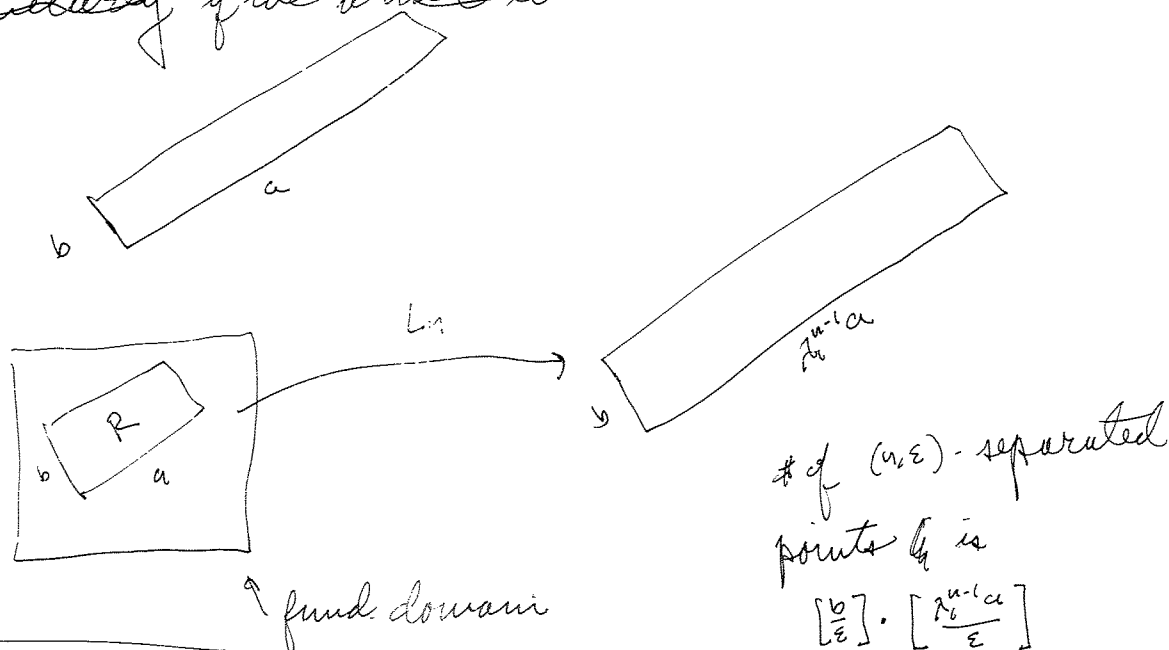
Then  $|L_n(av_1 + bv_2)|_{\max}$

$$|L_n(av_1 + bv_2)|_{\max} = |av_1 + bv_2|_n$$

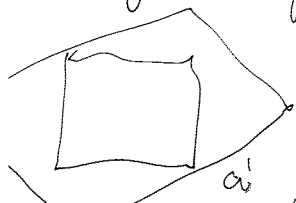
The advantage of describing  $| \cdot |_{\max}$  in terms of the map  $L_n$  ( $L_n?$ ) is the following.

If we want to see how many  $d_n^\varepsilon$  balls it takes to cover a given region  $R$  this will be the same as the number of  $d_{\max}$  balls it takes to cover the region  $L_n(R)$ . (similarly for  $(n, \varepsilon)$ -separated sets.)

~~similarly if we want to~~



# of  $(n, \varepsilon)$ -spanning points



$$\leq \left(\left\lceil \frac{b'}{\varepsilon} \right\rceil + 1\right) \left(\left\lceil \frac{r_i^{n-1} a'}{\varepsilon} \right\rceil + 1\right)$$

$$\limsup_{n \rightarrow \infty} \frac{1}{n} \log S(n, \varepsilon) \leq \log(r_i)$$

$$\limsup_{n \rightarrow \infty} \frac{1}{n} \log N(n, \varepsilon) \geq \lim_{n \rightarrow \infty} \frac{1}{n} \log \left\lceil \frac{b}{\varepsilon} \right\rceil \cdot \left\lceil \frac{r_i^{n-1} a}{\varepsilon} \right\rceil \geq \log(r_i)$$