## (Random) Trees of Intermediate Volume Growth

 University of Warwick
## (Random) trees of intermediate VOLUME GROWTH

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## Volume Growth in Graphs

$$
\left|B_{v}(r)\right|
$$

## Volume growth

ball ... $B(v, r):=\{x \in V(G) \mid \operatorname{dist}(x, v) \leq r\}$


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EXAMPLES: POLYNOMIAL AND EXPONENTIAL


## Geometric group theory

Cayley graph of $\mathbb{Z}^{2}=\langle x, y \mid x y=y x\rangle$.


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$F_{3} /\left\langle x^{2}, y^{2}, z^{2}\right\rangle$

## Typical questions \& Results

Question: are there groups of intermediate growth?
$:=$ super-polynomial but sub-exponential e.g. $\exp \left(r^{1 / 2}\right)$ or $r^{\log r}$

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Yes: Grigorchuk group (1984)

$$
|B(e, r)| \sim \exp \left(r^{\alpha}\right) \quad \text { with } 0.504<\alpha<0.7675
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Theorem. (Gromov; 1981)
$G$ is of polynomial growth $\Longleftrightarrow G$ is virtually nilpotent.

Theorem. (Trofimov; 1985)
Polynomial growth of vertex-transitive graphs must have integer degree.

## Beyond Cayley graphs


$|B(v, r)| \sim r^{2}$

$|B(v, r)| \in \theta\left(r^{2}\right)$

## Uniform Growth

Fix a function $g: \mathbb{R}_{\geq 0} \rightarrow \mathbb{R}_{\geq 0}$.

## Definition.

A graph $G$ is of uniform volume growth $g$ if there are $c_{1}, c_{2}, C_{1}, C_{2}>0$ so that

$$
C_{1} \cdot g\left(c_{1} r\right) \leq|B(v, r)| \leq C_{2} \cdot g\left(c_{2} r\right), \quad \text { for all } v \in V(G) \text { and } r \geq 0
$$

## Planar triangulations



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In the same paper: $\sim r^{\alpha}$ for arbitrary $\alpha \geq 1$.

- Benjamini, Georgakopoulos (2021): $\sim r^{\alpha}$ with $\alpha<2$, then quasi-tree

Planar triangulations of Growth $r^{3 / 2}$


## PLANAR TRIANGULATIONS OF GROWTH $r^{3 / 2}$



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Trees

## Uniform growth of Trees

What kind of uniform growth can a tree have?

- linear $\checkmark$
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$$
|B(v, r)| \sim r^{\alpha}, \quad \text { where } \alpha=\frac{\log |E(T)|}{\log \operatorname{diam}(T)}=\frac{\log 5}{\log 3} \approx 1.464973 .
$$

## The Question

## super-polynomial: $e^{\omega(\log (r))}$ <br> sub-exponential: $e^{o(r)}$

Q: "Are there unimodular trees of uniform intermediate volume growth?"

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super-polynomial: $e^{\omega(\log (r))}$
sub-exponential: $e^{o(r)}$
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Idea: find them as spanning trees of known intermediate growth graphs.











## Does it work ... ?

## Question

Given a graph of uniform growth $g$. Is there a (spanning) tree $T \subseteq G$ of the same uniform growth $g$ ?

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Given a graph of uniform growth $g$. Is there a (spanning) tree $T \subseteq G$ of the same uniform growth $g$ ?

Turns out we don't need the ambient graph!

## The Construction

$$
T_{0} \subset T_{1} \subset T_{2} \subset T_{3} \subset \cdots
$$

## Construction - A sequence of trees

Given: sequence $\delta_{1}, \delta_{2}, \delta_{3}, \ldots \in \mathbb{N}, \delta_{n} \geq 1$
$T_{0}$

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\delta_{n}:=n+2 \quad 3 \quad 4 \quad 5
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## Heuristics argument




Properties of $T_{n}$ :

- number of vertices: $\left(\delta_{1}+1\right) \cdots\left(\delta_{n}+1\right)$

- distance from center to an apocentric vertex: $2^{n}-1$


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$$
|B(v, \underbrace{2^{n}-1}_{r})|=\left(\delta_{1}+1\right) \cdots\left(\delta_{n}+1\right) \sim n!\sim(\log r)!\sim r^{\log \log r}
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$$

## EXAMPLE: POLYNOMIAL GROWTH

$$
|B(v, r)| \stackrel{!}{=}(r+1)^{2}
$$



ExAMPLE: POLYNOMIAL GROWTH

$$
|B(v, r)| \stackrel{!}{=}(r+1)^{2} \Longrightarrow\left|B\left(v, 2^{n}-1\right)\right|=\left(2^{n}\right)^{2}=4^{n}=(3+1) \cdots(3+1)
$$



ExAMPLE: EXPONENTIAL GROWTH

$$
\delta_{n}:=2^{2^{n}}
$$



$$
\begin{aligned}
\left|B\left(v, 2^{n}-1\right)\right|=\left(\delta_{1}+1\right) \cdots\left(\delta_{n}+1\right)=\prod_{k=1}^{n}\left(2^{2^{k-1}}+1\right)=\sum_{i=0}^{2^{n}-1} 2^{i} & =2^{2^{n}}-1 \\
& \sim 2^{r+1}-1
\end{aligned}
$$

## Main Result

For every "nice" function $g: \mathbb{R}_{\geq 0} \rightarrow \mathbb{R}_{\geq 0}$ there is a tree of uniform growth $g$.

## What are "Nice" Functions?

- $g$ is increasing
- $g$ grows at least linearly
- $g$ grows at most exponentially
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\begin{aligned}
g \text { super-additive } & \Longrightarrow g(a+b) \geq g(a)+g(b) \\
& \Longrightarrow g\left(2^{n+1}\right) \geq 2 g\left(2^{n}\right) \Longrightarrow \delta_{n} \approx \frac{g\left(2^{n+1}\right)}{g\left(2^{n}\right)}-1 \geq 1 .
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## What are "nice" Functions?

- $g$ is increasing
- $g$ grows at least linearly $\left(\delta_{n} \geq 1\right)$
- $g$ grows at most exponentially (bounded degree)
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## Main Result: $T$ has uniform growth

$$
\Delta(n):=\frac{\delta_{n}}{\delta_{1} \cdots \delta_{n-1}}, \quad \bar{\Delta}:=\sup _{n}\lceil\Delta(n)\rceil, \quad \Gamma:=\sup _{m \geq n}\left\lceil\frac{\Delta(m)}{\Delta(n)}\right\rceil
$$

Theorem. (Kontogeorgiou, W.; 2022)
For super-additive $g: \mathbb{R}_{\geq 0} \rightarrow \mathbb{R}_{\geq 0}$ exists a tree $T$ so that for all $v \in V(T)$ and $r \geq 0$

$$
\begin{aligned}
& |B(v, r)| \geq C_{1} \cdot g(r / 4) \\
& \text { if } \bar{\Delta}<\infty \text { then } \quad|B(v, r)| \leq C_{2} \cdot g(2 r)^{2} \\
& \text { if } \Gamma<\infty \text { then }|B(v, r)| \leq C_{3} \cdot g(4 r)
\end{aligned}
$$

In particular, if $\Gamma<\infty$, then $T$ is of uniform growth $g$.

## MAIn RESULT: $T$ HAS UNIFORM GROWTH

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\left.\begin{array}{rl} 
& |B(v, r)|
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In particular, if $\Gamma<\infty$, then $T$ is of uniform growth $g$.

## Theorem.

If $g$ is super-additive and (eventually) log-concave, then there is a tree of uniform volume growth $g$.

## Unimodular Trees

## The original question

Q: "Are there unimodular trees of uniform intermediate volume growth?" "unimodular random rooted trees"

- Itai Benjamini


## Alternative limits



## Alternative limits



## Apocentric limit



## Benjamini-Schramm Limits



## Benjamini-Schramm Limits



## Benjamini-Schramm Limits



## Benjamini-Schramm Limits



## Benjamini-Schramm Limits



## Benjamini-SCHRAMM Limits

$$
T_{0}, T_{1}, T_{2}, T_{3}, \ldots \quad \stackrel{\mathrm{BS}}{\longrightarrow} \mathcal{T}
$$

- Benjamini-Schramm limits are unimodular
- a set of graphs of uniformly bounded degree is compact
- every sequence of uniformly bounded degree has a convergent subsequence.


## Theorem.

If $g$ is super-additive and (eventually) log-concave, then there is a unimodular random rooted tree of uniform volume growth $g$.

## A THRESHOLD PHENOMENON

## Theorem. (structure theorem)

(i) if $g \in \Omega\left(r^{\log \log r}\right)$, then $\mathcal{T}$ is a.s. an apocentric limit.
(ii) if $g \in \mathcal{O}\left(r^{\alpha \log \log r}\right)$ for some $\alpha>1$, then $\mathcal{T}$ is a.s. a mixed limit.

- if growth is fast enough the Benjamini-Schramm limit can be a deterministic tree.

$$
\left|B_{T}(v, r)\right| \sim \exp \left(r^{\alpha}\right) \quad \text { where } \alpha=\log (\phi) \approx 0.6942
$$

## Question

Do general unimodular trees of uniform growth show a similar threshold phenomenon?

## Google

## Ähnliche Fragen

What are the stages of tree growth?

What is the growth of a tree called?

Why do plants have indeterminate growth?

What might different species of trees in a forest compete for?
® Bilder zu Trees of intermediate growth

apical meristem

tree trunk

canopy closure


## Thank you.


G. Kontogeorgiou and M. Winter (2022), arXiv "(Random) Trees of Intermediate Volume Growth"

