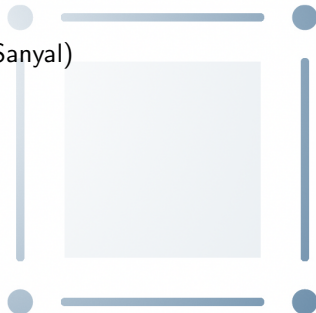


KALAI'S 3^d CONJECTURE FOR COORDINATE
SYMMETRIC POLYTOPES

Martin Winter
(joint work with Raman Sanyal)

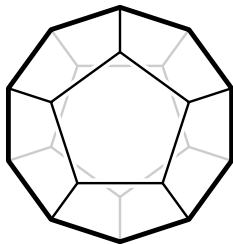
University of Warwick

17. January, 2024



CONVEX POLYTOPES

$$P = \text{conv}\{p_1, \dots, p_n\} \subset \mathbb{R}^d, \quad d \geq 1$$



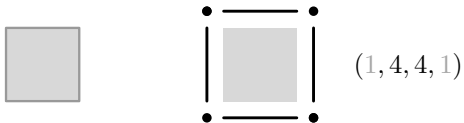
- ▶ general dimension $d \geq 1$
- ▶ general combinatorics (not simple/simplicial etc.)
- ▶ general geometry (not lattice, etc.)

POLYHEDRAL COMBINATORICS



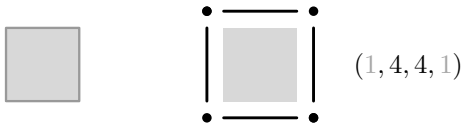
f -vector ... $\mathbf{f} = (f_{-1}, f_0, f_1, \dots, f_{d-2}, f_{d-1}, f_d)$

POLYHEDRAL COMBINATORICS



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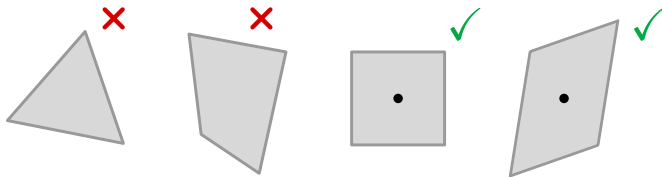
- ▶ Euler-Poincaré identity: $f_{-1} - f_0 + f_1 - \dots + (-1)^{d+1} f_d = 0$
- ▶ Dehn-Sommerville equations: (for simplicial polytopes)

$$\sum_{j=k}^{d-1} (-1)^j \binom{j+1}{k+1} f_j = (-1)^{d-1} f_k, \quad \text{for } k \in \{-1, \dots, d-2\}$$

- ▶ lower bound theorem / upper bound theorem / g -theorem

CENTRALLY SYMMETRIC POLYTOPES

centrally symmetric $:\Leftrightarrow P = -P$



KALAI'S 3^d CONJECTURE

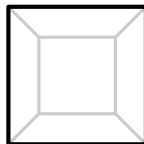
COUNTING FACES

$$s(P) := \cancel{f_{-1}} + f_0 + f_1 + \cdots + f_{d-1} + f_d = \#\underline{\text{non-empty faces}}$$

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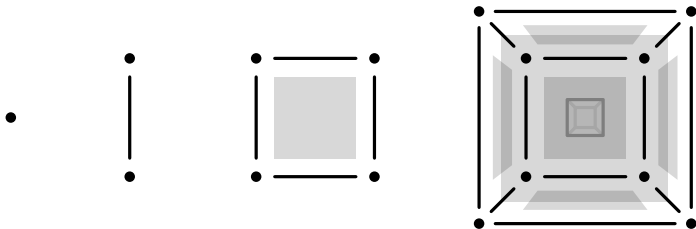
Example: d -cube $:= [-1, 1]^d$, $d \geq 0$



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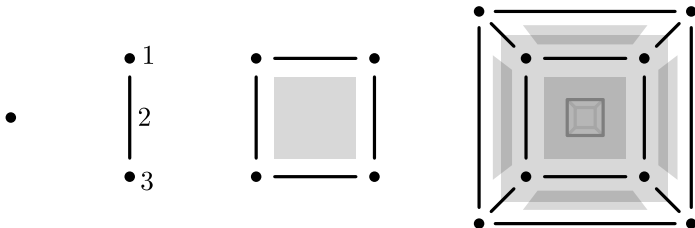
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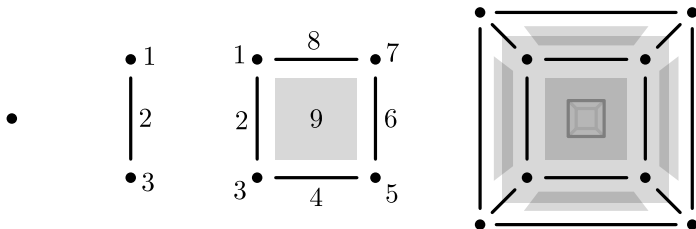
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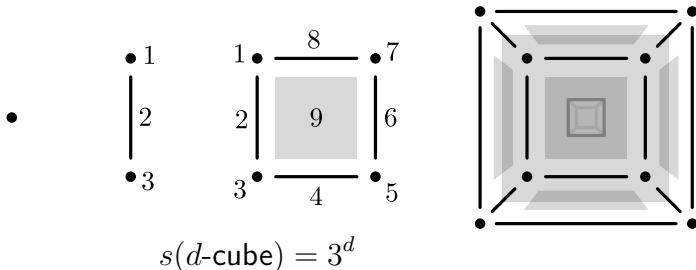
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Conjecture. (3^d conjecture, KALAI, 1989)

For every centrally symmetric d -polytope $P \subset \mathbb{R}^d$ holds

$$s(P) \geq s(d\text{-cube}) = 3^d.$$

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What is known ... ?

- ▶ dimension $d \leq 3$ ✓ easy
- ▶ dimension $d = 4$ ✓ not so easy (2007)
- ▶ simple/simplicial polytopes ✓ needs a lot of algebra
- ▶ without requiring central symmetry ✓ easy $\rightarrow s(d\text{-simplex}) = 2^d - 1$

HANNER POLYTOPES

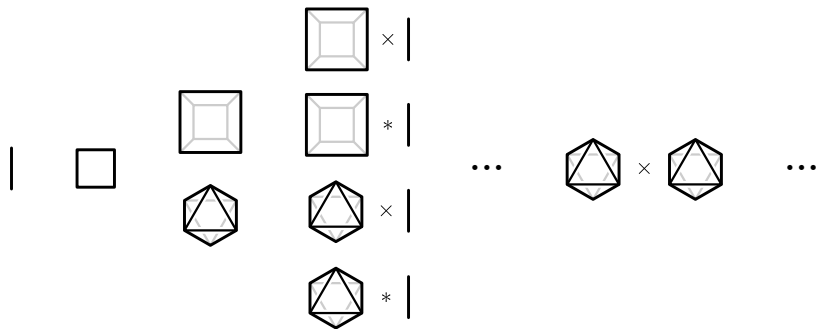
Hanner polytopes are defined recursively:

- (i) start from a line segment.
- (ii) recursively apply Cartesian products (\times) and sums ($*$)

$$| \times \text{---} = \square$$

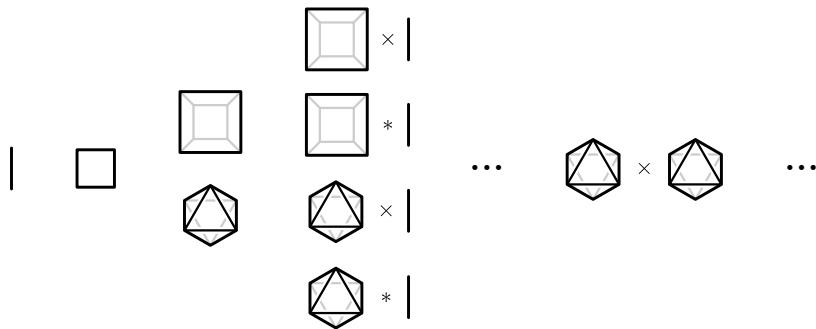
$$| * \text{---} = \diamond$$

HANNER POLYTOPES



#Hanner polytopes for $d \geq 1 = 1, 1, 2, 4, 8, 18, 40, 94, 224, 548, \dots$

HANNER POLYTOPES



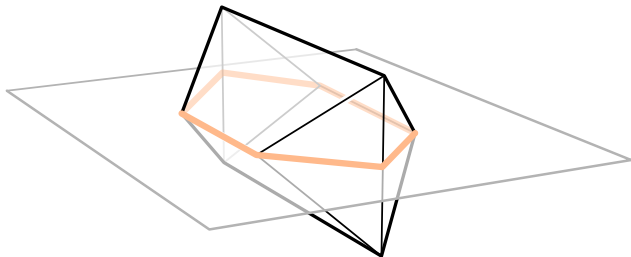
#Hanner polytopes for $d \geq 1 = 1, 1, 2, 4, 8, 18, 40, 94, 224, 548, \dots$

Note: Hanner polytopes do not minimize face numbers per dimension.

NAIVE APPROACH

Idea: induction by dimension

- ▶ slice your d -polytope P by a central hyperplane
→ this yields a $(d-1)$ -polytope P'
→ P' has 3^{d-1} faces by induction hypothesis
- ▶ somehow find three times as many faces in P



KALAI'S 3^d CONJECTURE

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MAHLER'S CONJECTURE

Mahler volume ... $M(P) := \text{vol}(P) \cdot \text{vol}(P^\circ)$

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For every centrally symmetric d -polytope $P \subset \mathbb{R}^d$ holds

$$\text{measures "roundness"} \rightarrow M(P) \geq M(d\text{-cube}) = \frac{4^d}{d!}.$$

But: cube is not the only minimizer! \rightarrow **Hanner polytopes**

What is known ... ?

- ▶ dimension $d \leq 3$ ✓ not so easy ($d = 2$: 1939, $d = 3$: 2020)
- ▶ dimension $d = 4$? out of reach
- ▶ cube is a local minimizer ✓ (2010)
- ▶ without requiring central symmetry ? open $\rightarrow M(d\text{-simplex}) = \frac{(d+1)^{d+1}}{(d!)^2}$

POLAR DUALITY AND MAHLER VOLUME

Mahler volume ... $M(P) := \text{vol}(P) \cdot \text{vol}(P^\circ)$.

polar dual ... $P^\circ := \{ x \in \mathbb{R}^d \mid \langle x, y \rangle \leq 1 \text{ for all } y \in P \}$

Examples:	line segment	1:1 \longleftrightarrow	line segment
	n -gon	\longleftrightarrow	n -gon (but rotated)
	cube	\longleftrightarrow	octahedron
	d -cube	\longleftrightarrow	d -crosspolytope
	Hanner polytope	\longleftrightarrow	(other) Hanner polytope

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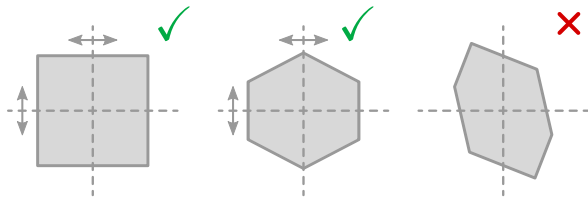
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$$M(d\text{-cube}) \stackrel{?}{\leq} M(P) \leq M(d\text{-ball}) \quad (\text{BLASCHKE, SANTALÓ})$$

(LOCALLY)
COORDINATE SYMMETRIC
POLYTOPES

COORDINATE SYMMETRIC POLYTOPES

= UNCONDITIONAL POLYTOPES



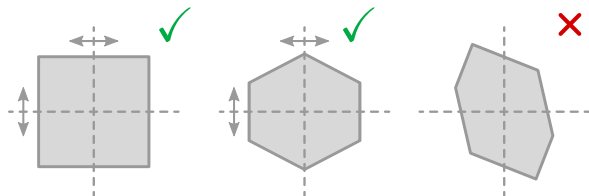
coordinate symmetric \iff symmetric w.r.t. all coordinate hyperplanes

Theorem. (SANYAL, W.; 2023+)

Kalai's conjecture holds for coordinate symmetric polytopes. + minimizers

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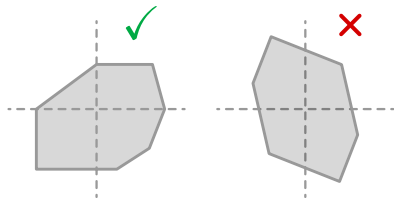
Theorem. (SAINT-RAYMOND, 1980; REISNER, 1987)

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LOCALLY COORDINATE SYMMETRIC POLYTOPES

= LOCALLY ANTI-BLOCKING POLYTOPES

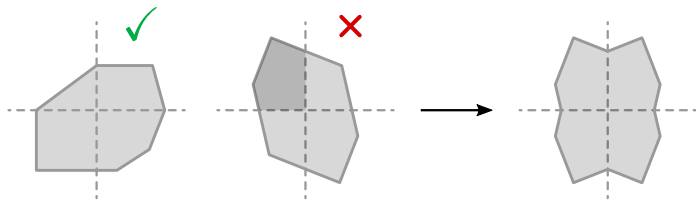
All discussed results hold more generally ...

**locally coordinate symmetric**: \iff “looks” coordinate symmetric in every orthant

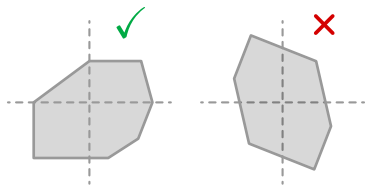
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All discussed results hold more generally ...

**locally coordinate symmetric**: \iff “looks” coordinate symmetric in every orthant**Surprise:** minimizers are still coordinate symmetric for Kalai's conjecture!

BUT NOTE ...

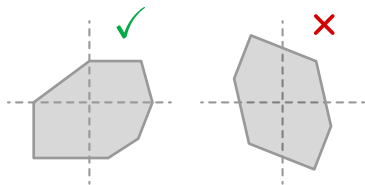


“(locally) coordinate symmetric” is a **geometric** notion

BUT

$s(P) \leq 3^d$ is a **combinatorial** notion

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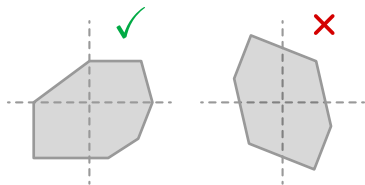
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$s(P) \leq 3^d$ is a **combinatorial** notion

Challenges:

- Find a polytope that is *provably not* combinatorially equivalent to a locally coordinate-symmetric polytope.

BUT NOTE ...



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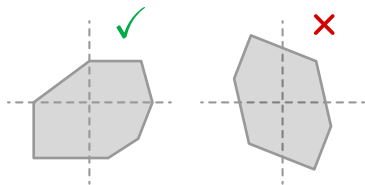
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Challenges:

- ▶ Find a polytope that is *provably not* combinatorially equivalent to a locally coordinate-symmetric polytope.
- ▶ Find many of them.
- ▶ Find many with few faces.

WE HAVE TWO PROOFS

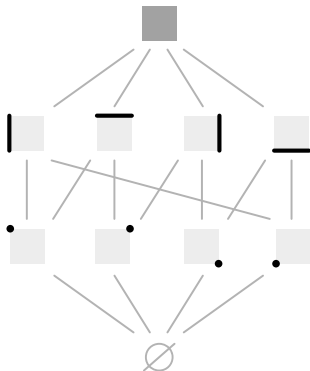
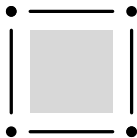
Proof I:

- ▶ mostly combinatorial
- ▶ makes the naive approach work \rightarrow uses induction
- ▶ counts faces in a clever way

Proof II:

- ▶ almost entirely geometric
- ▶ reminds of the original proof for Mahler for coordinate symmetric !
- ▶ no induction !!
- ▶ we used it to classify minimizers

THE FACE LATTICE



$$\mathcal{F}(P) := \{ \text{faces of } P \text{ ordered by inclusion} \}$$

THE FACE LATTICE

Meet, join, intervals, ...

- ▶ $\sigma \wedge \tau :=$ largest face that contained in σ and $\tau = \sigma \cap \tau$
- ▶ $\sigma \vee \tau :=$ smallest face that contains σ and τ
- ▶ $[\sigma, \tau] :=$ faces that are contained in τ and contain σ

Properties:

- ▶ $\mathcal{F}(P^\circ)$ is “upside down” $\mathcal{F}(P)$
- ▶ face lattices are **complemented**: for $\sigma \in \mathcal{F}(P)$ exists $\tau \in \mathcal{F}(P)$ s.t.

$$\sigma \wedge \tau = \emptyset \quad \text{and} \quad \sigma \vee \tau = P.$$

- ▶ intervals in face lattices are face lattice

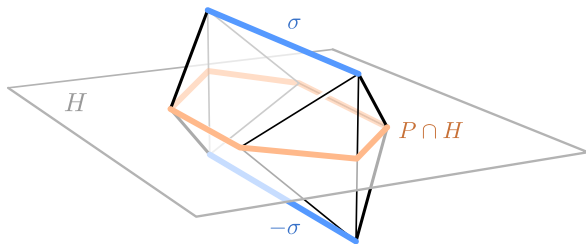
KALAI FOR COORDINATE SYMMETRIC POLYTOPES (I)

$$s(P) := f_{-1} + f_0 + f_1 + \cdots + f_{d-1} + f_d = \#\text{non-empty faces}$$

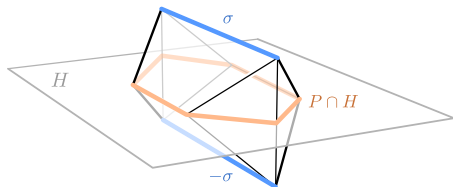
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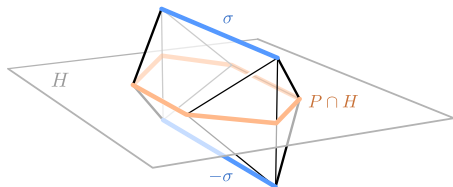


$$S_+ := \{\text{faces of } P \text{ that intersect } \sigma \text{ but not } -\sigma\}$$

$$S_0 := \{\text{faces of } P \text{ that intersect both or neither of } \pm\sigma\}$$

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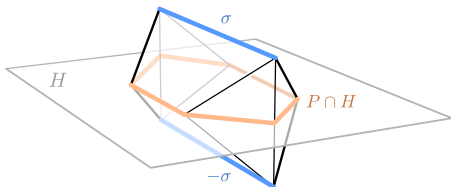
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$$\left. \begin{array}{l} S_+ \\ S_0 \\ S_- \end{array} \right\} |S_\bullet| \geq 3^{d-1}$$

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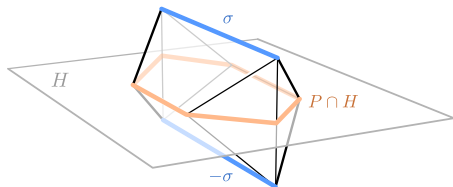
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$$|S_0| = s(P \cap H) \stackrel{\text{IH}}{\geq} 3^{d-1}$$

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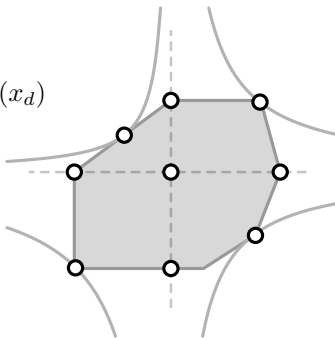
$$|S_0| = s(P \cap H) \stackrel{\text{IH}}{\geq} 3^{d-1}$$

$$\begin{aligned} |S_+| &\geq \left| \left\{ \text{compl}(\sigma; [\check{\sigma}, \hat{\sigma}]) \mid \emptyset \subset \check{\sigma} \subseteq \sigma \subseteq \hat{\sigma} \subset P \right\} \right| \\ &= s((\emptyset, \sigma]) \cdot s([\sigma, P)) = s(\sigma) \cdot s(\sigma^\diamond) \stackrel{\text{IH}}{\geq} 3^{\dim(\sigma)} \cdot 3^{d-1-\dim(\sigma)} = 3^{d-1} \end{aligned}$$

KALAI FOR COORDINATE SYMMETRIC POLYTOPES (II)

$$s \in \underbrace{\{-, 0, +\}^d}_{\#=3^d}, \quad s\text{-orthant} \dots \mathbb{R}_s^d := \{x \in \mathbb{R}^d \mid \text{sign}(x) = s\}$$

1. identify one *special point* per orthant
 = maximum of $s_1 \log(x_1) + \dots + s_d \log(x_d)$
 in $P \cap \mathbb{R}_s^d$
2. **show:** a face contains at most one special point



CLASSIFICATION OF MINIMIZERS

If P has $s(P) = 3^d$ then define

$$G(P) := \begin{cases} \{1, \dots, d\} \\ i \sim j \iff P \cap \{x_i, x_j = 0\} \text{ is an axis-aligned rectangle} \end{cases}$$

$G(P)$ does not contain a path of length three.

\implies such graphs are *cographs*. (graph analogues of Hanner polytopes)

One can reconstruct P from $G(P)$:

$$P(G) = \text{conv} \left\{ \sum_{i \in I} \pm e_i \mid I \subseteq \{1, \dots, d\} \text{ where } G[I] \text{ is a complete graph} \right\}.$$

If G is a cograph, then $P(G)$ is a Hanner polytope.



KALAI'S FLAG CONJECTURE

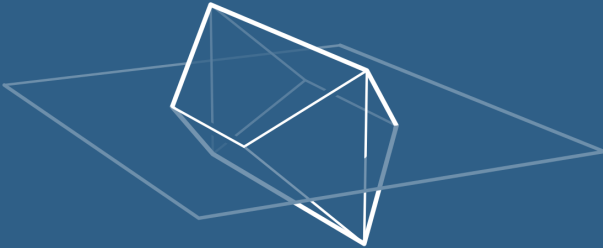
$$S(P) := \# \text{flags of } P$$

Conjecture.

For every centrally symmetric d -polytope $P \subset \mathbb{R}^d$ holds

$$S(P) \geq S(d\text{-cube}) = d! 2^d.$$

Thank you.



R. Sanyal, M. Winter. (arXiv:2308.02909)

"Kalai's 3^d -conjecture for unconditional and locally anti-blocking polytopes".

MAHLER FOR COORDINATE SYMMETRIC POLYTOPES

Mahler volume ... $M(P) := \text{vol}(P) \cdot \text{vol}(P^\circ)$.

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Theorem. (SAINT-RAYMOND, 1980; REISNER, 1987)

For every coordinate symmetric d -polytope $P \subset \mathbb{R}^d$ holds

$$M(P) \geq M(d\text{-cube}) = \frac{4^d}{d!}.$$

$$\text{vol}(P) = 2^d \text{vol}(P^+)$$