

Applications of entropy compression method to graph colorings

Jan Volec

Probabilistic method in combinatorics

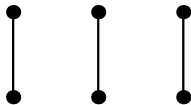
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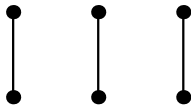
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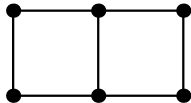
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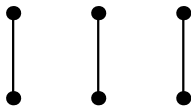
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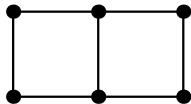
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$$\mathbf{P}[\text{Fail}] \leq \sum_{i=1}^m 1 - \mathbf{P}[P_i] = \frac{m}{k}, \text{ which is } < 1 \text{ for } k = m + 1$$

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- Example: graph G with maximum degree $\Delta \rightarrow \exists 4\Delta$ -coloring
- Well, simple greedy algorithm gives $(\Delta + 1)$ -coloring. . .

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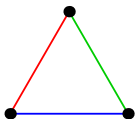
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- The method independently discovered by Schweitzer (2009)

Edge-colorings of graphs

- if we want to color edges s.t. incident edges \rightarrow different colors

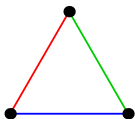
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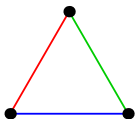
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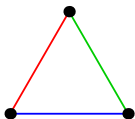
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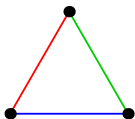
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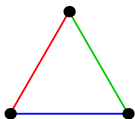
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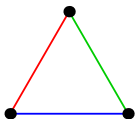
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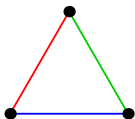
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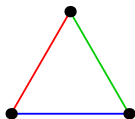
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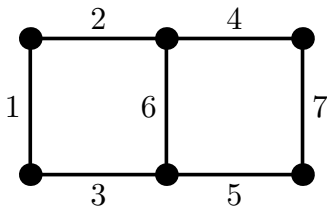
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colors from \mathcal{R}

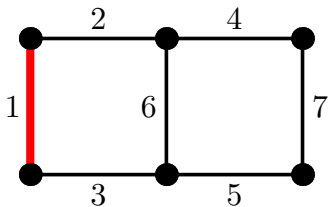


COLOR LOG:

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colors from \mathcal{R} **R**

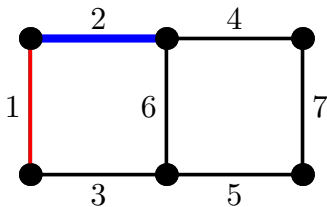


COLOR LOG: C

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colors from \mathcal{R} R **B**

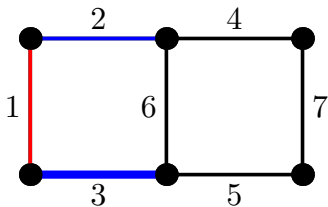


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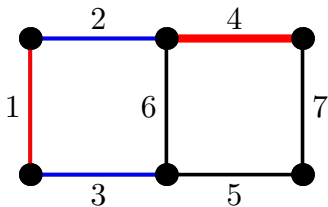


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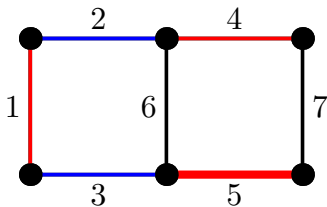


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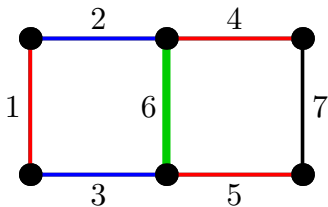


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colors from \mathcal{R} R B B R R **G**

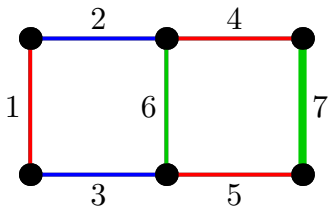


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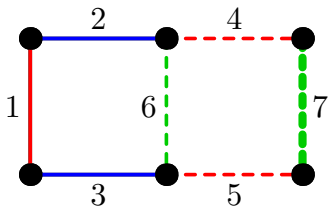


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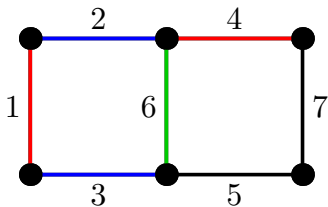
COLOR LOG: C C C C C C C C

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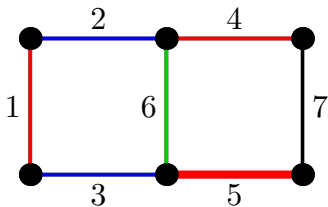
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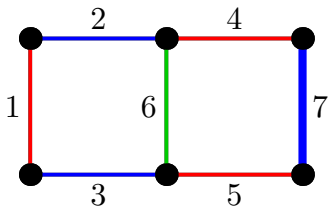
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colors from \mathcal{R} R B B R R G G R **B**



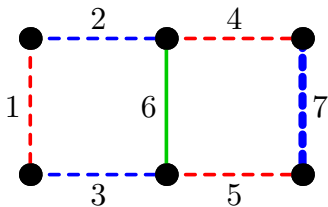
COLOR LOG: C C C C C C C U U C **C**

CYCLE LOG: 0 0

Claim: # of 2ℓ -cycles around fixed edge e is at most $(\Delta - 1)^{2\ell-2}$

- $e :=$ uncolored edge with minimum index
- Try to color e , does it create any 2-colored 2ℓ -cycle?
- If YES, then uncolor edges $c_3, c_4, \dots, c_{\ell-1}$ and e
- write LOG record about this step

colors from \mathcal{R} R B B R R G G R **B**



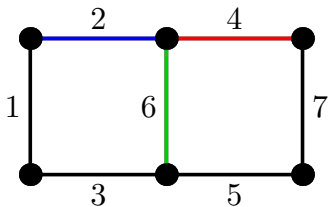
COLOR LOG: C C C C C C C U U C C

CYCLE LOG: 0 0

Claim: # of 2ℓ -cycles around fixed edge e is at most $(\Delta - 1)^{2\ell-2}$

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- If YES, then uncolor edges $c_3, c_4, \dots, c_{\ell-1}$ and e
- write LOG record about this step

colors from \mathcal{R} R B B R R G G R **B**



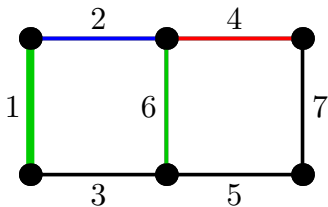
COLOR LOG: C C C C C C C U U C C **U U U U**

CYCLE LOG: 0 0 **0 0 0 0**

Claim: # of 2ℓ -cycles around fixed edge e is at most $(\Delta - 1)^{2\ell-2}$

- $e :=$ uncolored edge with minimum index
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- write LOG record about this step

colors from \mathcal{R} R B B R R G G R B **G**



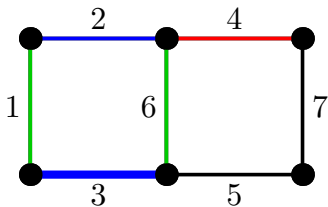
COLOR LOG: C C C C C C U U C C U U U U **C**

CYCLE LOG: 0 0 0 0 0 0

Claim: # of 2ℓ -cycles around fixed edge e is at most $(\Delta - 1)^{2\ell-2}$

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- Try to color e , does it create any 2-colored 2ℓ -cycle?
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- write LOG record about this step

colors from \mathcal{R} R B B R R G G R B G **B**



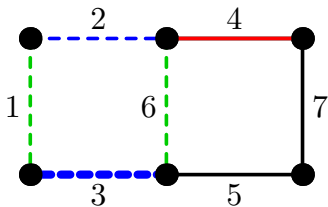
COLOR LOG: C C C C C C U U C C U U U U C **C**

CYCLE LOG: 0 0 0 0 0 0

Claim: # of 2ℓ -cycles around fixed edge e is at most $(\Delta - 1)^{2\ell-2}$

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- Try to color e , does it create any 2-colored 2ℓ -cycle?
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- write LOG record about this step

colors from \mathcal{R} R B B R R G G R B G **B**



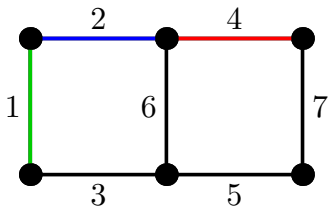
COLOR LOG: C C C C C C U U C C U U U C C

CYCLE LOG: 0 0 0 0 0 0

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- write LOG record about this step

colors from \mathcal{R} R B B R R G G R B G **B**



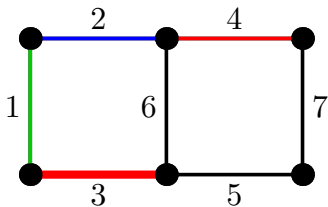
COLOR LOG: C C C C C C U U C C U U U U C C **U U**

CYCLE LOG: 0 0 0 0 0 0 **0 0**

Claim: # of 2ℓ -cycles around fixed edge e is at most $(\Delta - 1)^{2\ell-2}$

- $e :=$ uncolored edge with minimum index
- Try to color e , does it create any 2-colored 2ℓ -cycle?
- If YES, then uncolor edges $c_3, c_4, \dots, c_{\ell-1}$ and e
- write LOG record about this step

colors from \mathcal{R} R B B R R G G R B G B **R**



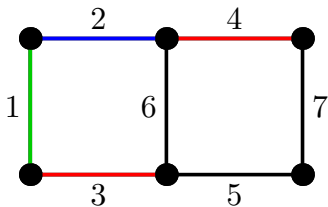
COLOR LOG: C C C C C C U U C C U U U U C C U U **C**

CYCLE LOG: 0 0 0 0 0 0 0 0

Claim: # of 2ℓ -cycles around fixed edge e is at most $(\Delta - 1)^{2\ell-2}$

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colors from \mathcal{R} R B B R R G G R B G B R...



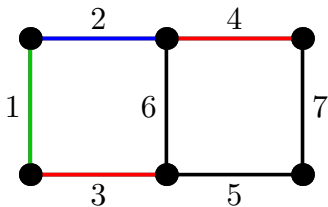
COLOR LOG: C C C C C C U U C C U U U U C C U U C ...

CYCLE LOG: 0 0 0 0 0 0 0 0 ...

Claim: # of 2ℓ -cycles around fixed edge e is at most $(\Delta - 1)^{2\ell-2}$

- $e :=$ uncolored edge with minimum index
- Try to color e , does it create any 2-colored 2ℓ -cycle?
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colors from \mathcal{R} R B B R R G G R B G B R...



COLOR LOG: C C C C C C U U C C U U U U C C U U C ...

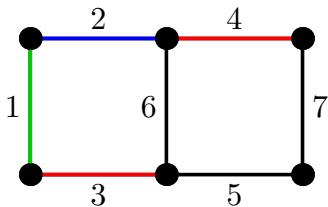
CYCLE LOG: 0 0 0 0 0 0 0 0 ...

Claim: # of 2ℓ -cycles around fixed edge e is at most $(\Delta - 1)^{2\ell-2}$

Claim: COLOR and CYCLE LOGs \rightarrow current set of colored edges

- $e :=$ uncolored edge with minimum index
- Try to color e , does it create any 2-colored 2ℓ -cycle?
- If YES, then uncolor edges $c_3, c_4, \dots, c_{\ell-1}$ and e
- write LOG record about this step

colors from \mathcal{R} R B B R R G G R B G B R...



COLOR LOG: C C C C C C U U C C U U U U C C U U C ...

CYCLE LOG: 0 0 0 0 0 0 0 0 ...

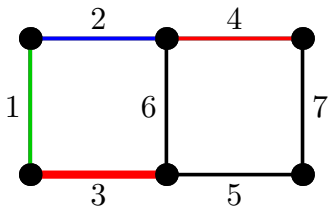
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Claim: current coloring and LOGs \rightarrow the whole sequence \mathcal{R}

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- write LOG record about this step

colors from \mathcal{R} R B B R R G G R B G B **R**...



COLOR LOG: C C C C C C U U C C U U U U C C U U **C** ...

CYCLE LOG: 0 0 0 0 0 0 0 0 ...

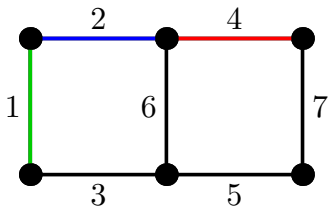
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colors from \mathcal{R} R B B R R G G R B G B **R**...



COLOR LOG: C C C C C C U U C C U U U U C C U U ~~X~~...

CYCLE LOG: 0 0 0 0 0 0 0 0 ...

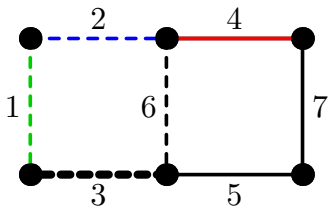
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colors from \mathcal{R} R B B R R G G R B G **B** R...



COLOR LOG: C C C C C C U U C C U U U U C **C U U** ~~X~~ ...

CYCLE LOG: 0 0 0 0 0 0 **0 0** ...

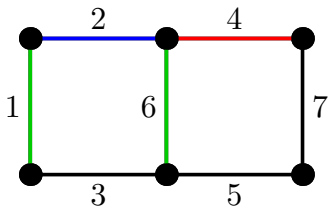
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- write LOG record about this step

colors from \mathcal{R} R B B R R G G R B G **B** R...



COLOR LOG: C C C C C C U U C C U U U U C ~~C~~ ~~C~~ ~~C~~ ~~C~~ ...

CYCLE LOG: 0 0 0 0 0 0 ~~0~~ ~~0~~ ...

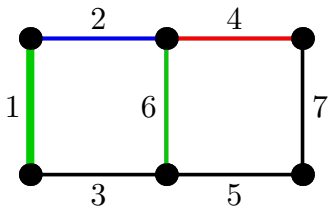
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- write LOG record about this step

colors from \mathcal{R} R B B R R G G R B **G** B R...



COLOR LOG: C C C C C C U U C C U U U U **C** ~~X~~ ~~X~~ ~~X~~ ~~X~~ ...

CYCLE LOG: 0 0 0 0 0 0 ~~X~~ ~~X~~ ...

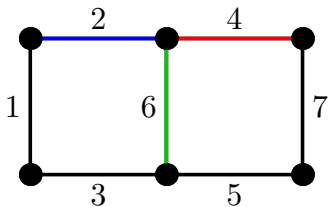
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colors from \mathcal{R} R B B R R G G R B **G** B R...



COLOR LOG: C C C C C C U U C C U U U U ~~X~~ ~~X~~ ~~X~~ ~~X~~ ~~X~~ ...

CYCLE LOG: 0 0 0 0 0 0 ~~X~~ ~~X~~ ...

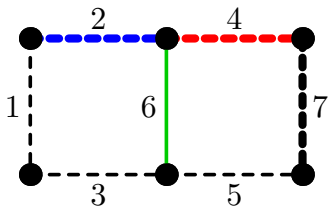
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- write LOG record about this step

colors from \mathcal{R} R B B R R G G R **B** G B R...



COLOR LOG: C C C C C C C U U C **C U U U U** ~~C~~ ~~C~~ ~~U~~ ~~U~~ ~~U~~ ~~U~~ ~~U~~ ...

CYCLE LOG: 0 0 **0 0 0 0** ~~0~~ ~~0~~ ...

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Claim: COLOR and CYCLE LOGs \rightarrow current set of colored edges

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- If YES, then uncolor edges $c_3, c_4, \dots, c_{\ell-1}$ and e
- write LOG record about this step

colors from \mathcal{R} R B B R R G G R B G B R...

COLOR LOG: ~~X~~...

CYCLE LOG: ~~O~~~~O~~~~O~~~~O~~~~O~~~~O~~~~O~~~~O~~...

Claim: # of 2ℓ -cycles around fixed edge e is at most $(\Delta - 1)^{2\ell-2}$

Claim: COLOR and CYCLE LOGs \rightarrow current set of colored edges

Claim: current coloring and LOGs \rightarrow the whole sequence \mathcal{R}

After t steps: # of bad \mathcal{R} 's \leq # of partial colorings \times # of LOGs

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- write LOG record about this step

colors from \mathcal{R} R B B R R G G R B G B R...

COLOR LOG: ~~X~~...

CYCLE LOG: ~~O~~~~O~~~~O~~~~O~~~~O~~~~O~~~~O~~~~O~~...

Claim: # of 2ℓ -cycles around fixed edge e is at most $(\Delta - 1)^{2\ell-2}$

Claim: COLOR and CYCLE LOGs \rightarrow current set of colored edges

Claim: current coloring and LOGs \rightarrow the whole sequence \mathcal{R}

After t steps: # of bad \mathcal{R} 's \leq # of partial colorings \times # of LOGs

But observe: $t - m \leq$ # of uncolorings $(u) \leq t$

Thank you for your attention!