

Integrable systems and duality

Based on Chen, de Gier, Wheeler

“Integrable stochastic dualities and the deformed Knizhnik–Zamolodchikov equation”

- ▶ The notion of the duality: operator / matrix formalism.
- ▶ Multi-species ASEP process.
- ▶ Hecke algebra, tKZ equation and trivial duality function.
- ▶ Non-symmetric Macdonald polynomials.
- ▶ ASEP polynomials: construction of non-trivial duality functions.
- ▶ Matrix product Ansatz and examples.
- ▶ Global duality.

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Duality for Markov processes: Definition

- ▶ Given two Markov processes on state spaces A and B ,

$$X = (\Omega_A, \mathcal{F}_A, (X_t)_{t \geq 0}, \{\mathbb{P}_a\}_{a \in A}), \quad Y = (\Omega_B, \mathcal{F}_B, (Y_t)_{t \geq 0}, \{\mathbb{P}^b\}_{b \in B}).$$

- ▶ A function $\Psi : A \times B \rightarrow \mathbb{C}$ is a duality function for X and Y iff:
- ▶ Let (P_t) and (Q_t) be semi-groups of (X_t) and (Y_t) . Recall: $P_t f(a) = \mathbb{E}_a[f(X_t)]$.
- ▶ Assume that the infinitesimal generators L and M are well-defined.

$$\begin{aligned} \forall a \in A, b \in B, \forall t \geq 0, \quad \mathbb{E}_a[\Psi(X_t, b)] &= \mathbb{E}^b[\Psi(a, Y_t)]. \\ \iff P_t \Psi(\cdot, b)(a) &= Q_t \Psi(a, \cdot)(b). \\ \iff L[\Psi(\cdot, b)](a) &= M[\Psi(a, \cdot)](b). \end{aligned}$$

- ▶ Write

$L[\Psi(\cdot, b)](a)$	$=$	$\sum_{a' \in A} \ell(a, a') \Psi(a', b),$	}	left actions	(Gen)
$M[\Psi(a, \cdot)](b)$	$=$	$\sum_{b' \in B} m(b, b') \Psi(a, b').$			

Duality for Markov processes: Definition

$\Psi : A \times B \rightarrow \mathbb{C}$ is a duality function.



$$\forall a \in A, b \in B, \quad L[\Psi(\cdot, b)](a) = M[\Psi(a, \cdot)](b)$$

where

$$L[\Psi(\cdot, b)](a) = \sum_{a' \in A} \ell(a, a') \Psi(a', b),$$

$$M[\Psi(a, \cdot)](b) = \sum_{b' \in B} m(b, b') \Psi(a, b').$$

(Gen)

Duality for Markov processes: Matrix formalism

- ▶ Consider vector spaces with elements of A and B as bases,

$$\mathbb{A} = \text{Span} \{|a\rangle, a \in A\}, \quad \mathbb{B} = \text{Span} \{|b\rangle, b \in B\}.$$

- ▶ Define an element in the tensor product $|\Psi\rangle \in \mathbb{A} \otimes \mathbb{B}$,

$$|\Psi\rangle = \sum_{\substack{a \in A \\ b \in B}} \Psi(a, b) |a\rangle \otimes |b\rangle.$$

- ▶ Define linear operators $\mathbb{L} \in \text{End}(\mathbb{A})$ and $\mathbb{M} \in \text{End}(\mathbb{B})$,

$$\mathbb{L}|a\rangle = \sum_{a' \in A} \ell(a', a) |a'\rangle, \quad \mathbb{M}|b\rangle = \sum_{b' \in B} m(b', b) |b'\rangle,$$

where coefficients $\ell(a', a)$ and $m(b', b)$ are given by (Gen).

- ▶ \mathbb{L} and \mathbb{M} are right actions.

Duality for Markov processes: Matrix formalism

Prop. 1.1 : Duality relation is equivalent to $\mathbb{L}|\Psi\rangle = \mathbb{M}|\Psi\rangle$.

Proof: Direct computation.

$$\begin{aligned}\mathbb{L}|\Psi\rangle &= \sum_{a' \in A} \sum_{\substack{a \in A \\ b \in B}} \ell(a', a) \Psi(a, b) |a'\rangle \otimes |b\rangle \\ &= \sum_{\substack{a \in A \\ b \in B}} \underbrace{\sum_{a' \in A} \ell(a, a') \Psi(a', b)}_{L[\Psi(\cdot, b)](a)} |a\rangle \otimes |b\rangle . \\ \mathbb{M}|\Psi\rangle &= \sum_{b' \in A} \sum_{\substack{a \in A \\ b \in B}} m(b', b) \Psi(a, b) |a\rangle \otimes |b'\rangle \\ &= \sum_{\substack{a \in A \\ b \in B}} \underbrace{\sum_{b' \in B} m(b, b') \Psi(a, b')}_{M[\Psi(a, \cdot)](b)} |a\rangle \otimes |b\rangle .\end{aligned}$$



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mASEP on the segment

The multi-species asymmetric simple exclusion process (mASEP):

- ▶ On the segment $\{1, \dots, n\}$.
- ▶ At most one particle per site.
- ▶ Particles of type $1, \dots, r$. Empty site is of type 0.
- ▶ **Higher-type** particle can switch with a **lower-type** particle.
- ▶ Switch rates are: τ_{right} to the right, τ_{left} to the left. For $i > j$,



- ▶ ASEP $\Leftrightarrow r = 1$.

mASEP on the segment

A configuration can be encoded in two ways.

Example for $n = 5$:



Monomial: $z^\nu \in \mathbb{C}[z_1, \dots, z_n]$. \rightsquigarrow Example: $z_2 z_3^2 z_5$.

$$\blacktriangleright A = \left\{ z^\nu := \prod_{i=1}^n z_i^{\nu_i}, 0 \leq \nu_i \leq r \right\}.$$

$\blacktriangleright \mathbb{A} = \{P \in \mathbb{C}[z_1, \dots, z_n] \text{ s.t. the degree in each variable is at most } r\}.$

Sequence: $\mu = (\mu_i)_{1 \leq i \leq n} \in \{0, \dots, r\}^n$. \rightsquigarrow Example: $|0, 1, 2, 0, 1\rangle$.

$$\blacktriangleright B = \left\{ |\mu_1 \dots \mu_n\rangle := \bigotimes_{i=1}^n |\mu_i\rangle_i, 0 \leq \mu_i \leq r \right\}$$

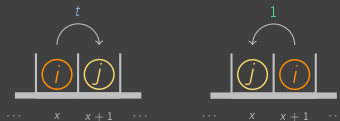
$\blacktriangleright \mathbb{B} = \bigotimes_{i=1}^n V_i$ where $V_i \cong \mathbb{C}^{r+1}$ and $\{|0\rangle, \dots, |r\rangle\} = \text{canonical basis of } \mathbb{C}^{r+1}$.

Particle content: $0^{m_0} 1^{m_1} \dots r^{m_r}$, $m_i = |\{k : \mu_k = i\}|$. \rightsquigarrow Example: $0^2 1^2 2^1$.

mASEP on the segment: Local generator

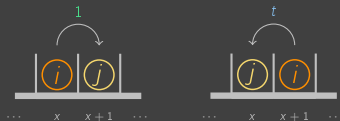
- ▶ Fix $t > 0$.
- ▶ Consider two processes:

▶ $\tau_{\text{right}} = t$ and $\tau_{\text{left}} = 1$:



for $i > j$.

▶ $\tau_{\text{right}} = 1$ and $\tau_{\text{left}} = t$:



for $i > j$.

- ▶ The forward dynamics of one is the backward dynamics of the other.
- ▶ Denote their generators by \mathbb{L} and \mathbb{M} .
- ▶ Initial conditions might be different!

mASEP on the segment: Local generator

- ▶ The local generators \mathbb{L}_i and \mathbb{M}_i act on particles at positions i and $i + 1$.

$$\mathbb{L} = \sum_{i=1}^{n-1} \mathbb{L}_i, \quad \mathbb{M} = \sum_{i=1}^{n-1} \mathbb{M}_i.$$

- ▶ The local operators L_i and M_i are,

$$L_i[f](\nu) = \sum_{\nu'} \underbrace{\ell_i(\nu, \nu')}_{\tau(\nu \rightarrow \nu')} f(\nu'), \quad M_i[f](\mu) = \sum_{\mu'} \underbrace{m_i(\mu, \mu')}_{\tau(\mu \rightarrow \mu')} f(\mu').$$

- ▶ Example for ASEP ($r = 1$):

$$L_i = (\ell_i(\nu, \nu'))_{\nu, \nu'} \quad M_i = (m_i(\mu, \mu'))_{\mu, \mu'}$$

	00	01	10	11
00)	0	0	0
01		0	-1	1
10		0	t	-t
11		0	0	0

	00	01	10	11
00)	0	0	0
01		0	-t	t
10		0	1	-1
11		0	0	0

mASEP on the segment: $r = 1$, ASEP

- ▶ Define $\mathbb{L}_i = \left(\frac{tz_i - z_{i+1}}{z_i - z_{i+1}} \right) (s_i - 1)$, where s_i acts on polynomials by $z_i \leftrightarrow z_{i+1}$.

Prop. 2.1 : Let $|\nu\rangle = \prod_{i=1}^n z_i^{\nu_i}$. Then, $\mathbb{L}_i |\nu\rangle = \sum_{\nu' \in \mathcal{A}} \ell_i(\nu', \nu) |\nu'\rangle$.

Proof:

- ▶ $\mathbb{L}_i |1\rangle = \mathbb{L}_i |z_i z_{i+1}\rangle = 0$ due to the symmetry.

- ▶ $\mathbb{L}_i |z_i\rangle = \left(\frac{tz_i - z_{i+1}}{z_i - z_{i+1}} \right) (z_{i+1} - z_i) = z_{i+1} - tz_i$.

$$\frac{\textcircled{1}}{i \quad i+1} \rightsquigarrow 1 \cdot \frac{\quad \textcircled{1}}{i \quad i+1} - t \cdot \frac{\textcircled{1}}{i \quad i+1}$$

- ▶ $\mathbb{L}_i |z_{i+1}\rangle = \left(\frac{tz_i - z_{i+1}}{z_i - z_{i+1}} \right) (z_i - z_{i+1}) = tz_i - z_{i+1}$.

$$\frac{\quad \textcircled{1}}{i \quad i+1} \rightsquigarrow t \cdot \frac{\textcircled{1}}{i \quad i+1} - 1 \cdot \frac{\textcircled{1}}{i \quad i+1}$$

$$L_i = (\ell_i(\nu, \nu'))_{\nu, \nu'}$$

	00	01	10	11
00	0	0	0	0
01	0	-1	1	0
10	0	t	-t	0
11	0	0	0	0

□

mASEP on the segment: Local duality

- ▶ A function $\Psi : A \times B \rightarrow \mathbb{C}$ is a local duality function for X and Y iff:

$$\mathbb{L}_i |\Psi\rangle = \mathbb{M}_i |\Psi\rangle, \quad 1 \leq i \leq n-1, \quad (\text{LD})$$

where

$$|\Psi\rangle = \sum_{\nu \in A} \sum_{\mu \in B} \Psi(\nu, \mu) \prod_{k=1}^n z_k^{\nu_k} |\mu\rangle.$$

- ▶ Local duality \Rightarrow global duality. (The reverse is false in general.)
- ▶ We focus on the local duality, i.e. solutions of (LD).

mASEP on the segment: Local duality

- ▶ The following duality functions for ASEP are constructed in [BCS+14]¹.

Theorem 4.1 & 4.2 [Duality functions of ASEP on \mathbb{Z}]:

- ▶ ASEP occupation process $(\nu_x)_{x \in \mathbb{Z}}$ with rates $(\tau_{\text{right}}, \tau_{\text{left}}) = (1, t)$;
- ▶ Position process $\vec{x}(\mu) = (x_1(\mu) < \dots < x_n(\mu))$ with rates $(\tau_{\text{right}}, \tau_{\text{left}}) = (t, 1)$.

$$\Psi(\nu, \mu) = \prod_{x \in \vec{x}(\mu)} \left(\prod_{i < x} t^{\nu_i} \right) \nu_x,$$

$$\Psi(\nu, \mu) = \prod_{x \in \vec{x}(\mu)} \left(\prod_{i \leq x} t^{\nu_i} \right).$$

- ▶ The first one can be recovered using the framework of² but not the second one.

¹From duality to determinants for q-TASEP and ASEP (2014)

²Integrable stochastic dualities and the deformed Knizhnik–Zamolodchikov equation (2020)

mASEP on the segment: Something wrong...

- ▶ Define $\mathbb{L}_i = \left(\frac{tz_i - z_{i+1}}{z_i - z_{i+1}} \right) (s_i - 1)$, where s_i acts on polynomials by $z_i \leftrightarrow z_{i+1}$.

Prop. 2.1 : Let $|\nu\rangle = \prod_{i=1}^n z_i^{\nu_i}$. Then, $\mathbb{L}_i |\nu\rangle = \sum_{\nu' \in \mathcal{A}} \ell_i(\nu', \nu) |\nu'\rangle$.

- ▶ \mathbb{L}_i acts on polynomials faithfully for $r = 1$.

- ▶ For $r \geq 2$, consider the local dynamics starting from $\frac{|\textcircled{2}\rangle}{i \ i+1}$.

$$\begin{aligned} \mathbb{L}_i |z_i^2\rangle &= \left(\frac{tz_i - z_{i+1}}{z_i - z_{i+1}} \right) (s_i - 1) \cdot z_i^2 = \left(\frac{tz_i - z_{i+1}}{z_i - z_{i+1}} \right) (z_{i+1}^2 - z_i^2) \\ &= z_{i+1}^2 + (1-t)z_i z_{i+1} - tz_i^2. \end{aligned}$$

- ▶ But the dynamics gives $\frac{|\textcircled{2}\rangle}{i \ i+1} \rightsquigarrow 1 \cdot \frac{|\textcircled{2}\rangle}{i \ i+1} - t \cdot \frac{|\textcircled{2}\rangle}{i \ i+1}$.

mASEP on the segment: Something wrong...

- ▶ Define $\mathbb{L}_i = \left(\frac{tz_i - z_{i+1}}{z_i - z_{i+1}} \right) (s_i - 1)$, where s_i acts on polynomials by $z_i \leftrightarrow z_{i+1}$.

Prop. 2.1 : Let $|\nu\rangle = \prod_{i=1}^n z_i^{\nu_i}$. Then, $\mathbb{L}_i |\nu\rangle = \sum_{\nu' \in \mathbb{A}} \ell_i(\nu', \nu) |\nu'\rangle$.

- ▶ \mathbb{L}_i acts on polynomials faithfully for $r = 1$.
- ▶ Conclusion: the basis $\left\{ |\nu\rangle = \prod_{i=1}^n z_i^{\nu_i} \right\}$ is not a good basis to work with.
- ▶ We look for a basis of polynomials $\{|\nu\rangle = f_\nu(z)\}$ such that

$$\mathbb{L}_i |\nu\rangle = \sum_{\nu' \in \mathbb{A}} \ell_i(\nu', \nu) |\nu'\rangle, \quad \forall \nu. \quad (\text{Adm})$$

We say that $\{|\nu\rangle = f_\nu(z)\}$ is admissible.

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mASEP on the segment: Hecke algebra

- Type A_{n-1} Hecke algebra with generators $\{T_i\}_{1 \leq i \leq n-1}$ and relations,

$$(T_i - t)(T_i + 1) = 0, \quad (\text{quadratic relation})$$

$$T_i T_{i+1} T_i = T_{i+1} T_i T_{i+1}, \quad (\text{braid relation})$$

$$T_i T_j = T_j T_i, \quad \forall i, j, |i - j| > 1. \quad (\text{commutativity})$$

- Generators and inverses can be realized as operators on the space of polynomials $\mathbb{C}[z_1, \dots, z_n]$,

$$\begin{cases} T_i &= t - \left(\frac{tz_i - z_{i+1}}{z_i - z_{i+1}} \right) (1 - s_i), & (\Rightarrow \mathbb{L}_i = T_i - t), \\ T_i^{-1} &= t^{-1} - t^{-1} \left(\frac{tz_i - z_{i+1}}{z_i - z_{i+1}} \right) (1 - s_i). \end{cases}$$

mASEP on the segment: tKZ equation

- ▶ Consider a family of polynomials $\{f_\nu(z) : \nu = (\nu_1, \dots, \nu_n)\}$ in $\mathbb{C}[z_1, \dots, z_n]$.
- ▶ We say that $\{f_\nu(z)\}$ is a solution to the ASEP exchange relations if

$$T_i f_\nu = \begin{cases} f_{s_i \nu}, & \text{if } \nu_i > \nu_{i+1}, \\ t f_\nu, & \text{if } \nu_i = \nu_{i+1}, \end{cases} \quad (\text{tKZ})$$

for all ν and $1 \leq i \leq n-1$.

- ▶ Using the quadratic relation $T_i^2 + (1-t)T_i - t = 0$, (tKZ) gives,

$$T_i f_\nu = (t-1)f_\nu + t f_{s_i \nu} \quad \text{if } \nu_i < \nu_{i+1}.$$

$$\begin{aligned} (s_i \nu)_i > (s_i \nu)_{i+1} &\implies T_i f_{s_i \nu} = f_\nu \implies \underbrace{T_i^2}_{(t-1)T_i + t} f_{s_i \nu} = T_i f_\nu \\ &\implies (t-1)f_\nu + t f_{s_i \nu} = T_i f_\nu. \end{aligned}$$

□

mASEP on the segment: tKZ equation

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- ▶ Using the quadratic relation $T_i^2 + (1-t)T_i - t = 0$, (tKZ) gives,

$$T_i f_\nu = (t-1)f_\nu + t f_{s_i \nu} \quad \text{if } \nu_i < \nu_{i+1}.$$

- ▶ Let $\mathbb{L}_i = T_i - t$. $\mathbb{L}_i f_\nu = \begin{cases} f_{s_i \nu} - t f_\nu, & \text{if } \nu_i > \nu_{i+1}, \\ 0, & \text{if } \nu_i = \nu_{i+1}, \\ t f_{s_i \nu} - f_\nu, & \text{if } \nu_i < \nu_{i+1}. \end{cases}$

\rightsquigarrow (Adm) is satisfied.

mASEP on the segment: tKZ equation

- ▶ Given a family of polynomials $\{f_\nu(z) : \nu = (\nu_1, \dots, \nu_n)\}$ in $\mathbb{C}[z_1, \dots, z_n]$.

$$T_i f_\nu = \begin{cases} f_{s_i \nu}, & \text{if } \nu_i > \nu_{i+1}, \\ t f_\nu, & \text{if } \nu_i = \nu_{i+1}, \end{cases} \quad (\text{tKZ})$$

for all ν and $1 \leq i \leq n-1$.



$$\mathbb{L}_i |\nu\rangle = \sum_{\nu' \in A} \ell_i(\nu', \nu) |\nu'\rangle, \quad \forall \nu. \quad (\text{Adm})$$

- ▶ Look for $\{f_\nu(z) : \nu = (\nu_1, \dots, \nu_n)\}$ satisfying tKZ.

From tKZ to duality

Prop. 3.2 : If $\{|\nu\rangle = f_\nu(z)\}$ satisfies (tKZ), then the function $\Psi(\nu, \mu) = \delta_{\nu, \mu}$ is a local mASEP duality function. In other words,

$$\mathbb{L}_i |\Psi\rangle = \mathbb{M}_i |\Psi\rangle, \quad 1 \leq i \leq n-1,$$

where $|\Psi\rangle = \sum_{\nu} \sum_{\mu} \Psi(\nu, \mu) f_\nu(z) |\mu\rangle = \sum_{\mu} f_\mu(z) |\mu\rangle$.

- ▶ Proof by direct computation.
- ▶ The trivial (diagonal) duality function is not interesting.
- ▶ Find a particular family of polynomials $\{f_\nu(z)\}$ satisfying (tKZ) and construct a non-trivial duality function.
- ▶ Key idea: construct functions $\{f_\nu(z)\}$ depending on additional parameters t and q and extract certain coefficients.

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Non-symmetric Macdonald Polynomials: Compositions

- ▶ A composition μ is an n -tuple of non-negative integers (μ_1, \dots, μ_n) .
- ▶ A partition λ is a composition s.t. $\lambda_1 \geq \dots \geq \lambda_n \geq 0$.
- ▶ An anti-partition δ is a composition s.t. $0 \leq \delta_1 \leq \dots \leq \delta_n$.
- ▶ For a given composition μ , write μ^+ and μ^- its unique partition / anti-partition obtained by permutation μ . \rightsquigarrow Example: $(1, 0, 2)^+ = (2, 1, 0)$.
- ▶ Given two compositions μ and ν , define two orders:

$$\mu \geq \nu \iff \sum_{i=1}^j \mu_i \geq \sum_{i=1}^j \nu_i, \quad \forall 1 \leq j \leq n.$$

$$\mu \succ \nu \iff \mu^+ > \nu^+ \text{ or } \mu^+ = \nu^+, \mu > \nu.$$

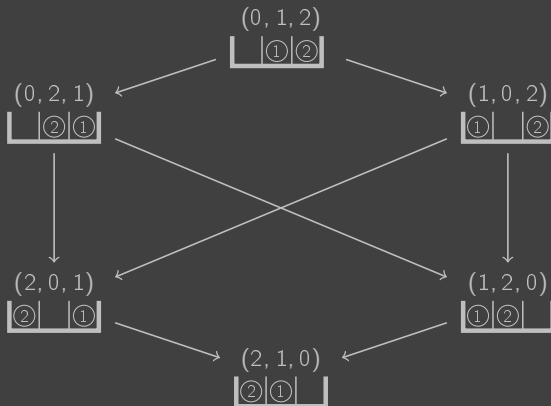
- ▶ Example: $(0, 0) \prec \boxed{(0, 1) \prec (1, 0)} \prec (1, 1) \prec \boxed{(0, 2) \prec (2, 0)}$.
- ▶ It is not a total order: $(2, 2) \not\prec (3, 0)$ and $(3, 0) \not\prec (2, 2)$

Non-symmetric Macdonald Polynomials: Compositions

For a composition μ , its composition sector $\sigma(\mu)$ is defined by

$$\sigma(\mu) := \{\nu \mid \mu^+ = \nu^+\} = S_n \cdot \mu.$$

Example: $n = 3$ with $\mu = (0, 1, 2)$.



Non-symmetric Macdonald Polynomials: Definition

- ▶ Multi-variate polynomial ring: $\mathbb{C}_{q,t}[z_1, \dots, z_n] = \mathbb{C}(q, t)[z_1, \dots, z_n]$ where t is the rate of ASEP and q is an additional parameter.
- ▶ Consider the Hecke algebra of type A_{n-1} as before. Define

$$(\omega g)(z_1, \dots, z_n) = g(qz_n, z_1, \dots, z_{n-1}).$$

- ▶ Define the Cherednik-Dunkl operators Y_i [Che91]¹ [Che95]²,

$$Y_i = T_i \dots T_{n-1} \omega T_1^{-1} \dots T_{i-1}^{-1}, \quad 1 \leq i \leq n.$$

- ▶ Fact: the operators (Y_i) commute, so can be jointly diagonalized.
- ▶ Non-symmetric Macdonald Polynomials (NSMP) are defined as “normalized” eigenfunctions of these operators.
- ▶ NSMP are indexed by compositions $\mu \in \mathbb{Z}_{\geq 0}^n$.

¹A unification of Knizhnik-Zamolodchikov and Dunkl operators via affine Hecke algebras (1991)

²Non-symmetric Macdonald's polynomials (1995)

Non-symmetric Macdonald Polynomials: Properties

- ▶ The change of basis w.r.t. the canonical basis is triangular:

$$E_\mu = z^\mu + \sum_{\nu \prec \mu} c_{\mu,\nu}(q, t) z^\nu, \quad c_{\mu,\nu}(q, t) \in \mathbb{Q}(q, t).$$

- ▶ Eigenvalues are given by, for all $1 \leq i \leq n$,

$$Y_i E_\mu = y_i(\mu; q, t) E_\mu,$$
$$y_i(\mu; q, t) = q^{\mu_i} t^{\rho(\mu)_i + n - i + 1}, \quad \rho(\mu) = -w_\mu \cdot (1, 2, \dots, n),$$

where $w_\mu \in S_n$ is the permutation with minimal length s.t. $\mu = w_\mu \cdot \mu^+$.

- ▶ Let μ be a composition such that $\mu_i < \mu_{i+1}$. Then

$$E_{s_i \mu} = t^{-1} \left(T_i + \frac{1-t}{1 - y_{i+1}(\mu)/y_i(\mu)} \right) E_\mu.$$

Non-symmetric Macdonald Polynomials: Examples

- ▶ Example for $n = 2$.

$$\begin{array}{rcl}
 E_{(0,0)} & = & 1 \\
 \hline
 E_{(0,1)} & = & z_2 \\
 \hline
 E_{(1,0)} & = & z_1 + \left(\frac{1-t}{1-qt} \right) z_2 \\
 \hline
 E_{(1,1)} & = & z_1 z_2 \\
 \hline
 E_{(0,2)} & = & \left(\frac{q(1-t)}{1-qt} \right) z_1 z_2 + z_2^2 \\
 \hline
 E_{(2,0)} & = & z_1^2 + \left(\frac{q(1-t)^2}{(1-q^2t)(1-qt)} + \frac{1-t}{1-qt} \right) z_1 z_2 + \left(\frac{1-t}{1-q^2t} \right) z_2^2
 \end{array}$$

- ▶ $(0, 0) \prec \boxed{(0, 1) \prec (1, 0)} \prec (1, 1) \prec \boxed{(0, 2) \prec (2, 0)}$.

- ▶ The order is defined in [CGW20]¹.

$$\mu \succcurlyeq \nu \iff \sum_{i=1}^j \mu_i \geq \sum_{i=1}^j \nu_i, \quad \forall 1 \leq j \leq n.$$

$$\mu \succ \nu \iff \mu^+ > \nu^+ \text{ or } \mu^+ = \nu^+, \mu > \nu.$$

¹Integrable stochastic dualities and the deformed Knizhnik–Zamolodchikov equation (2020)

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$$\begin{array}{rcl}
 E_{(0,0)} & = & 1 \\
 \hline
 E_{(1,0)} & = & z_1 \\
 \hline
 E_{(0,1)} & = & \left(\frac{1-t}{1-qt}\right) z_1 + z_2 \\
 \hline
 E_{(1,1)} & = & z_1 z_2 \\
 \hline
 E_{(2,0)} & = & z_1^2 + \left(\frac{q(1-t)}{1-qt}\right) z_1 z_2 \\
 \hline
 E_{(0,2)} & = & \left(\frac{1-t}{1-q^2t}\right) z_1^2 + \left(\frac{q(1-t)^2}{(1-q^2t)(1-qt)} + \frac{1-t}{1-qt}\right) z_1 z_2 + z_2^2
 \end{array}$$

- ▶ $(0, 0) \prec \boxed{(1, 0) \prec (0, 1)} \prec (1, 1) \prec \boxed{(2, 0) \prec (0, 2)}$.

- ▶ The order is defined in [HHL08]¹ also implemented in Sage.

$$\mu \succcurlyeq \nu \iff \sum_{i=1}^j \mu_{n-i+1} \geq \sum_{i=1}^j \nu_{n-i+1}, \quad \forall 1 \leq j \leq n.$$

$$\mu \succ \nu \iff \mu^+ > \nu^+ \text{ or } \mu^+ = \nu^+, \mu > \nu.$$

¹A combinatorial formula for nonsymmetric Macdonald polynomials (2008)

Non-symmetric Macdonald Polynomials: Coefficients

Conjecture 3.8 : Fix μ a composition. Let $m \in \mathbb{Q}_{>0}$ and $p \in \mathbb{N}$. If

$$\text{Coeff}_p [E_\mu, m] := \lim_{q \rightarrow t^{-m}} (1 - qt^m)^p E_\mu(z; q, t)$$

exists and is non-zero. Then, there exists a unique ν for which

$$\text{Coeff}_p [E_\mu, m] \propto E_\nu(z; t^{-m}, t) := \lim_{q \rightarrow t^{-m}} E_\nu(z; q, t).$$

- ▶ One can check some examples using Sage.
- ▶ Authors of [CGW20]¹ are able to show this conjecture only in some particular cases which is enough to construct some non-trivial duality functions.

¹Integrable stochastic dualities and the deformed Knizhnik–Zamolodchikov equation (2020)

Non-symmetric Macdonald Polynomials: Coefficients

Prop. 3.5 : Define generating series

$$Y(w) = \sum_{i=1}^n Y_i w^i, \quad y_\mu(w) = \sum_{i=1}^n y_i(\mu; q, t) w^i.$$

Then,

$$E_\mu(z; q, t) = \prod_{\nu \prec \mu} \frac{Y(w) - y_\nu(w)}{y_\mu(w) - y_\nu(w)} \cdot z^\mu.$$

► This can be seen as “Lagrange interpolation”.

Proof: Use the following two properties:

- For $\nu \prec \mu$, $\frac{Y(w) - y_\nu(w)}{y_\mu(w) - y_\nu(w)} \cdot E_\kappa = \begin{cases} 0, & \text{if } \kappa = \nu, \\ E_\mu, & \text{if } \kappa = \mu. \end{cases}$
- $E_\mu = z^\mu + \sum_{\nu \prec \mu} c_{\mu, \nu}(q, t) z^\nu \iff z^\mu = E_\mu + \sum_{\nu \prec \mu} d_{\mu, \nu}(q, t) E_\nu.$

□

Non-symmetric Macdonald Polynomials: Coefficients

Prop. 3.6 : Fix a composition μ . Let $m \in \mathbb{Q}_{>0}$ and $p \in \mathbb{N}$. If

$$\text{Coeff}_p [E_\mu, m] := \lim_{q \rightarrow t^{-m}} (1 - qt^m)^p E_\mu(z; q, t)$$

exists and is non-zero. Then,

$$\text{Coeff}_p [E_\mu, m] = \lim_{q \rightarrow t^{-m}} (1 - qt^m)^p \sum_{\nu \in \mathcal{E}_\mu} c_\nu(q, t) E_\nu(z; q, t),$$

where $\mathcal{E}_\mu = \{\nu : \nu \prec \mu, y_\nu(w) = y_\mu(w) \text{ at } q = t^{-m}\}$.

Recall that the eigenvalues are $y_i(\mu; q, t) = q^{\mu_i} t^{\rho(\mu)_i + n - i + 1}$.

Thus, $y_\mu(w) = y_\nu(w) \iff y_i(\mu) = y_i(\nu), \quad \forall 1 \leq i \leq n,$

$$\iff q^{\mu_i} t^{\rho(\mu)_i} = q^{\nu_i} t^{\rho(\nu)_i}, \quad \forall 1 \leq i \leq n.$$

Unique solution $\nu = \mu$ for generic q and t ; more solutions at $q = t^{-m}$.

Non-symmetric Macdonald Polynomials: Coefficients

Prop. 3.6 : Fix a composition μ . Let $m \in \mathbb{Q}_{>0}$ and $p \in \mathbb{N}$. If

$$\text{Coeff}_p [E_\mu, m] := \lim_{q \rightarrow t^{-m}} (1 - qt^m)^p E_\mu(z; q, t)$$

exists and is non-zero. Then,

$$\text{Coeff}_p [E_\mu, m] = \lim_{q \rightarrow t^{-m}} (1 - qt^m)^p \sum_{\nu \in \mathcal{E}_\mu} c_\nu(q, t) E_\nu(z; q, t),$$

where $\mathcal{E}_\mu = \{\nu : \nu \prec \mu, y_\nu(w) = y_\mu(w) \text{ at } q = t^{-m}\}$.

Proof: Use “Lagrange interpolation”,

$$\prod_{\substack{\nu \prec \mu \\ \nu \notin \mathcal{E}_\nu}} \frac{Y(w) - y_\nu(w)}{y_\mu(w) - y_\nu(w)} \cdot z^\mu = E_\mu + \sum_{\nu \in \mathcal{E}_\mu} d_{\mu, \nu}(q, t) E_\nu.$$

By taking $\lim_{q \rightarrow t^{-m}} (1 - qt^m)^p$ on both sides, LHS = 0 and RHS gives the proposition. \square

Non-symmetric Macdonald Polynomials: Coefficients

Prop. 3.7 : Assume that $\rho = |\mathcal{E}_\mu|$, then there exists a unique ν for which

$$E_\nu(z; t^{-m}, t) := \lim_{q \rightarrow t^{-m}} E_\nu(z; q, t)$$

is well-defined and $\text{Coeff}_\rho [E_\mu, m] \propto E_\nu(z; t^{-m}, t)$.

- ▶ Conjecture 3.8 is true with an additional condition $\rho = |\mathcal{E}_\mu|$.
- ▶ Proof uses “Lagrange interpolation” and by induction.
- ▶ The composition ν corresponds to the minimal composition in \mathcal{E}_μ .

Non-symmetric Macdonald Polynomials: Coefficients

- ▶ Example of \mathcal{E}_μ for $\mu = (0, 0, 2)$ and $m = 1$.
- ▶ w_μ is the permutation with minimal length s.t. $\mu = w_\mu \cdot \mu^+$.
 $\rightsquigarrow w_\mu = (13)$ since the permutation (312) is longer than (13) .
- ▶ Define $S_m(\mu) := (m\mu_i - \rho(\mu)_i)_i = m(0, 0, 2) + (3, 2, 1) = (3, 2, 2m + 1)$. Then,

$$\nu \in \mathcal{E}_\mu \iff S_m(\mu) = S_m(\nu).$$

- ▶ $\sum \mu_i = \sum \nu_i \Rightarrow \nu$ must contain at least one 0.
- ▶ Possible candidates for ν are s.t. $\nu^- = (0, 1, 1)$ or $(0, 0, 2)$.
 - ▷ If ν is s.t. $\nu^- = (0, 1, 1)$, then $\nu \prec \mu$ and $S_1(\mu) = S_1(\nu)$.
 - ▷ $(0, 0, 2)$ is a trivial solution. Moreover, its permutations are not smaller than μ for the partial order \prec .

Non-symmetric Macdonald Polynomials: Coefficients

- ▶ Example of \mathcal{E}_μ for $\mu = (0, 0, 2)$ and $m = 1$.
- ▶ w_μ is the permutation with minimal length s.t. $\mu = w_\mu \cdot \mu^+$.
 $\rightsquigarrow w_\mu = (13)$ since the permutation (312) is longer than (13) .
- ▶ Define $S_m(\mu) := (m\mu_i - \rho(\mu)_i)_i = m(0, 0, 2) + (3, 2, 1) = (3, 2, 2m + 1)$. Then,

$$\nu \in \mathcal{E}_\mu \iff S_m(\mu) = S_m(\nu).$$

- ▶ $\sum \mu_i = \sum \nu_i \Rightarrow \nu$ must contain at least one 0.
- ▶ Possible candidates for ν are s.t. $\nu^- = (0, 1, 1)$ or $(0, 0, 2)$.
- ▶ Conclusion: $\mathcal{E}_{(0,0,2)} = \sigma((0, 1, 1)) := S_3 \cdot (0, 1, 1)$ for $m = 1$.
- ▶ Similarly, one can show that:

If $\delta = (0^{n-m}, r^m)$, then $\mathcal{E}_\delta = \sigma(\varepsilon)$ where $\varepsilon = (0^{n-rm}, 1^{rm})$.

- ▶ The notion of the duality: operator / matrix formalism.
- ▶ Multi-species ASEP process.
- ▶ Hecke algebra, tKZ equation and trivial duality function.
- ▶ Non-symmetric Macdonald polynomials.
- ▶ ASEP polynomials: construction of non-trivial duality functions.
- ▶ Matrix product Ansatz and examples.
- ▶ Global duality.

ASEP polynomials: Definition

- ▶ Another basis **ASEP polynomials** can be defined from NSMP,

$$\begin{aligned}f_{\delta}(z; q, t) &= E_{\delta}(z; q, t), & \forall \delta = (\delta_1 \leq \dots \leq \delta_n), \\f_{s_i \mu} &= T_i^{-1} f_{\mu}, & \mu_i < \mu_{i+1}.\end{aligned}$$

- ▶ Well-defined due to the Hecke algebra structure.
- ▶ $f_{\mu} \neq E_{\mu}$ in general: the recursion formula is different.

Recall that $E_{s_i \mu} = t^{-1} \left(T_i + \frac{1-t}{1-y_{i+1}(\mu)/y_i(\mu)} \right) E_{\mu}$ for $\mu_i < \mu_{i+1}$.

ASEP polynomials: Definition

- ▶ Another basis **ASEP polynomials** can be defined from NSMP,

$$\begin{aligned} f_\delta(z; q, t) &= E_\delta(z; q, t), & \forall \delta = (\delta_1 \leq \dots \leq \delta_n), \\ f_{s, \mu} &= T_i^{-1} f_\mu, & \mu_i < \mu_{i+1}. \end{aligned}$$

- ▶ Well-defined due to the Hecke algebra structure.
- ▶ $f_\mu \neq E_\mu$ in general: the recursion formula is different.
- ▶ **ASEP polynomials** are triangular w.r.t. the canonical basis,

$$f_\mu = z^\mu + \sum_{\nu \prec \mu} c_{\mu, \nu}(q, t) z^\nu,$$

and satisfies (tKZ).

- ▶ $f_{\mu_n, \mu_1, \dots, \mu_{n-1}}(qz_n, z_1, \dots, z_{n-1}; q, t) = q^{\mu_n} f_\mu(z; q, t).$ (Cyclic-B.C.)
- ▶ Above two properties uniquely define **ASEP polynomials**.
- ▶ **ASEP** and **NSMP** are also related by a triangular change of basis.

ASEP polynomials: Properties

Prop. 3.10 : For any composition μ , the following expansions are unique,

$$E_\mu(z; q, t) = f_\mu(z; q, t) + \sum_{\substack{\nu \in \sigma(\mu) \\ \nu \prec \mu}} c_{\mu, \nu}(q, t) f_\nu(z; q, t);$$

$$f_\mu(z; q, t) = E_\mu(z; q, t) + \sum_{\substack{\nu \in \sigma(\mu) \\ \nu \prec \mu}} d_{\mu, \nu}(q, t) E_\nu(z; q, t).$$

Proof: For an anti-partition, $E_\delta = f_\delta$ so the proposition is true. By induction, assume μ is s.t. the proposition holds with $\mu_i < \mu_{i+1}$,

$$E_{s_i \mu} = t^{-1} \left(T_i + \frac{1-t}{1 - y_{i+1}(\mu)/y_i(\mu)} \right) \left(f_\mu + \sum_{\substack{\nu \in \sigma(\mu) \\ \nu \prec \mu}} c_{\mu, \nu}(q, t) f_\nu \right).$$

- ▶ $T_i f_\nu = (t-1)f_\nu + t f_{s_i \nu}$ if $\nu_i < \nu_{i+1}$;
- ▶ $\nu, s_i \nu \in \sigma(\mu) = \sigma(s_i \mu)$;
- ▶ $(\nu \prec \mu, \mu \prec s_i \mu) \Rightarrow \nu, s_i \nu \prec s_i \mu$.

ASEP polynomials: Properties

Thm 3.11 : Fix an anti-partition δ . Let $m \in \mathbb{Q}_{>0}$ and $p \in \mathbb{N}$ such that Conjecture 3.8 holds for δ . Then, there exists a unique anti-partition ε ,

$$f_\nu(z; t^{-m}, t) := \lim_{q \rightarrow t^{-m}} f_\nu(z; q, t)$$

is well-defined for all $\nu \in \sigma(\varepsilon)$ and

$$\text{Coeff}_p [f_\mu, m] = \sum_{\nu \in \sigma(\varepsilon)} \Psi(\nu, \mu; t) f_\nu(z; t^{-m}, t)$$

for all $\mu \in \sigma(\delta)$ and suitable coefficients $\Psi(\nu, \mu; t)$.

Proof:

- ▶ The proof is based on Prop. 3.10.
- ▶ Use the recurrence relation and the fact that T_i commutes with \lim .

ASEP polynomials: Local duality

Theorem 3.12 : Keep the notations from Theorem 3.11. $\Psi(\nu, \mu; t)$ defines a local duality function.

Proof: From Prop 3.2 (trivial duality function), the following holds

$$\mathbb{L}_i |\mathcal{I}\rangle = \mathbb{M}_i |\mathcal{I}\rangle, \quad 1 \leq i \leq n-1, \quad (1)$$

where $|\mathcal{I}\rangle = \sum_{\mu} f_{\mu}(z; q, t) |\mu\rangle$. Taking the coefficient,

$$|\mathcal{I}_{p,m}\rangle := \text{Coeff}_p [|\mathcal{I}\rangle, m] = \sum_{\mu \in \sigma(\delta)} \sum_{\nu \in \sigma(\epsilon)} \Psi(\nu, \mu; t) f_{\nu}(z; t^{-m}, t) |\mu\rangle$$

also satisfies $\mathbb{L}_i |\mathcal{I}_{p,m}\rangle = \mathbb{M}_i |\mathcal{I}_{p,m}\rangle, \quad 1 \leq i \leq n-1.$

- ▶ The notion of the duality: operator / matrix formalism.
- ▶ Multi-species ASEP process.
- ▶ Hecke algebra, tKZ equation and trivial duality function.
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- ▶ Matrix product Ansatz and examples.
- ▶ Global duality.

Matrix product Ansatz

- ▶ The coefficients $\Psi(\nu, \mu)$ are difficult to compute in general.
- ▶ The matrix product Ansatz from [CGW15]¹ is useful for ASEP polynomials.

Theorem : Fix $r \geq 1$. For a composition μ with largest part equal to r ,

$$f_\mu(z; q, t) = \Omega_\mu(q, t) \times \text{Tr}\left(A_{\mu_1}(z_1) \dots A_{\mu_n}(z_n) S\right), \quad (\text{Mat-Ans})$$

where $\{A_i(z)\}_{0 \leq i \leq r}$ and S are explicit matrices.

- ▶ Idea : translate (tKZ) and (Cyclic-B.C.) into matrices.

¹Matrix product formula for Macdonald polynomials (2015)

Matrix product Ansatz

Theorem : Fix $r \geq 1$. For a composition μ with largest part equal to r ,

$$f_\mu(z; q, t) = \Omega_\mu(q, t) \times \text{Tr}\left(A_{\mu_1}(z_1) \dots A_{\mu_n}(z_n) S\right), \quad (\text{Mat-Ans})$$

where $\{A_i(z)\}_{0 \leq i \leq r}$ and S are explicit matrices.

► Matrices involved in (Mat-Ans) are infinite-dimensional matrices ϕ , ϕ^\dagger , k ,

$$[\phi]_{i,j} = \delta_{i+1,j}(1 - t^i), \quad [\phi^\dagger]_{i,j} = \delta_{i,j+1}, \quad [k]_{i,j} = \delta_{i,j} t^i, \quad i, j \in \mathbb{N}.$$

► For $r = 1$,

$$A_0(z) = 1,$$

$$A_1(z) = z.$$

► For $r = 2$,

$$A_0(z) = 1 + z\phi,$$

$$A_1(z) = zk,$$

$$A_2(z) = z\phi^\dagger + z^2.$$

Matrix product Ansatz: $\mu^- = (0^{n-m}, r^m)$

Thm 5.5 : Let $r, m \in \mathbb{N}$ such that $n - rm \geq 0$. Take $\delta = (0^{n-m}, r^m)$. Then, $\text{Coeff}_1 [f_\delta, m]$ exists and

$$\text{Coeff}_1 [f_\mu, m] = \sum_{\nu \in \sigma(\varepsilon)} \psi(\nu, \mu; t) z^\nu, \quad \forall \mu \in \sigma(\delta),$$

where $\varepsilon = (0^{n-rm}, 1^{rm})$.

Proof:

► Check $\text{Coeff}_1 [f_\delta, m]$ exists using the explicit formula.

►
$$\text{Coeff}_\rho [E_\mu, m] = \lim_{q \rightarrow t^{-m}} (1 - qt^m)^\rho \sum_{\nu \in \mathcal{E}_\mu} c_\nu(q, t) E_\nu(z; q, t),$$

where $\mathcal{E}_\mu = \{\nu : \nu \prec \mu, y_\nu(w) = y_\mu(w) \text{ at } q = t^{-m}\}$ contains only rank-one compositions.

► Use the fact that for any rank-one composition ν , f_ν is multilinear in variables.

Matrix product Ansatz: $\mu^- = (0^{n-m}, r^m)$

Thm 5.6 : We have that, $\Psi(\nu, \mu) = d(t) \cdot t^{\Omega(\mu, \nu)} \cdot I(\mu, \nu)$, where

$$\blacktriangleright \Omega(\mu, \nu) = \sum_{1 \leq i < j \leq n} \mathbb{1}(\mu_i < \mu_j) \mathbb{1}(\nu_i = \nu_j = 1).$$

$$\blacktriangleright I(\mu, \nu) = \begin{cases} 0, & \exists k : (\mu_k, \nu_k) = (r, 0), \\ 1, & \text{otherwise.} \end{cases}$$

Proof: Check for $\mu = \delta^+ = (r^m, 0^{n-m})$ then by induction.

- $\blacktriangleright A_r(z)$ can be factorized by $z \implies f_{\delta^+}$ contains the factor $\prod_{i=1}^m z_i$.
- $\blacktriangleright f_{\delta^+}$ is homogeneous in (z_1, \dots, z_n) of degree rm .
- $\blacktriangleright f_{\delta^+}$ is symmetric in z_{m+1}, \dots, z_n .
- \blacktriangleright Thm 5.5: f_{δ^+} is multilinear in all the variables.

Matrix product Ansatz: $\mu^- = (0^{n-m}, r^m)$

Thm 5.6 : We have that, $\Psi(\nu, \mu) = d(t) \cdot t^{\Omega(\mu, \nu)} \cdot I(\mu, \nu)$, where

$$\blacktriangleright \Omega(\mu, \nu) = \sum_{1 \leq i < j \leq n} \mathbb{1}(\mu_i < \mu_j) \mathbb{1}(\nu_i = \nu_j = 1).$$

$$\blacktriangleright I(\mu, \nu) = \begin{cases} 0, & \exists k : (\mu_k, \nu_k) = (r, 0), \\ 1, & \text{otherwise.} \end{cases}$$

Proof: Check for $\mu = \delta^+ = (r^m, 0^{n-m})$ then by induction.

$$\begin{aligned} \Rightarrow \text{Coeff}_1 [f_{\delta^+}, m] &= d(t) \cdot \prod_{i=1}^m z_i \cdot e_{(rm-n)}(z_{m+1}, \dots, z_n) \\ &= d(t) \sum_{\nu \in \sigma(\varepsilon)} I(\delta^+, \nu) z^\nu, \end{aligned}$$

since $\Omega(\mu, \nu) = 0$ for all ν .

□

Matrix product Ansatz: $\mu^- = (0^{n-m}, 1^m)$

- ▶ Thm 5.6 recovers the duality function from [BCS+14]¹.

Thm 5.1 : Keeping the same notations, we have,

$$t^{\Omega(\mu, \nu)} \cdot I(\mu, \nu) = t^{-m(m-1)/2} \prod_{j=1}^m \left(\prod_{1 \leq i < x_j(\mu)} t^{\nu_i} \right) \nu_{x_j(\mu)}.$$

Proof:

- ▶ If $I(\mu, \nu) = 0$, then LHS = RHS = 0.
- ▶ If $I(\mu, \nu) = 1$, i.e. $\forall k, \mu_k \leq \nu_k$.

The set $A := \{(i, j) : i < j, \nu_i = 1, \mu_j = 1\}$ can be partitioned into:

- ▷ $A_1 := \{(i, j) : i < j, \mu_i = \mu_j = 1\}$; $|A_1| = \frac{1}{2}m(m-1)$.
- ▷ $A_2 := \{(i, j) : i < j, \mu_i = 0, \mu_j = 1\}$; $\Omega(\mu, \nu) = |A_2|$.

¹From duality to determinants for q-TASEP and ASEP (2014)

Matrix product Ansatz: $\mu^- = (0^{n-m}, 1^m)$

- ▶ Thm 5.6 recovers the duality function from [BCS+14]¹.

Thm 5.1 : Keeping the same notations, we have,

$$t^{\Omega(\mu, \nu)} \cdot I(\mu, \nu) = t^{-m(m-1)/2} \prod_{j=1}^m \left(\prod_{1 \leq i < x_j(\mu)} t^{\nu_i} \right) \nu_{x_j(\mu)}.$$

Proof:

- ▶ If $I(\mu, \nu) = 0$, then LHS = RHS = 0.
- ▶ If $I(\mu, \nu) = 1$, i.e. $\forall k, \mu_k \leq \nu_k$.

$$\implies \sum_{j=1}^m \sum_{i=1}^{x_j(\mu)-1} \nu_i = \Omega(\mu, \nu) + \frac{1}{2}m(m-1).$$

□

¹From duality to determinants for q-TASEP and ASEP (2014)

Matrix product Ansatz: $\mu^- = (0^{n-m_1-m_2}, 1^{m_1}, 2^{m_2})$

Prop. 6.4 & Cor. 6.5 : Let μ be a rank-two decomposition with

$$\vec{x}(\mu) = (x_1 < \cdots < x_{m_1}),$$

$$\vec{y}(\mu) = (y_1 < \cdots < y_{m_2}),$$

labelling the positions of 1 and 2-particles. Then,

$$\Psi(\nu, \mu) = \prod_{x \in \vec{x}(\mu)} \left(\prod_{i < x} t^{\nu_i} \right) \nu_x \cdot \prod_{y \in \vec{y}(\mu)} \left(\prod_{i < y} t^{\nu_i} \right) \nu_y \cdot t^{-\chi(\vec{x}, \vec{y})}$$

is a (local) duality function, where

$$\chi(\vec{x}, \vec{y}) := \#\{(x_i, y_j) \in (\vec{x}, \vec{y}) : x_i > y_j\}.$$

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Global duality: Rank-one

Prop. 7.1 : Let ν be an infinite rank-one composition. Let μ also be a rank-one composition and denote its particle positions to be $\vec{x}(\mu) = (x_1 < \dots < x_m)$. Then the observable

$$H(\nu, \mu) = \prod_{x \in \vec{x}(\mu)} \left(\prod_{i \leq x} t^{\nu_i} \right)$$

is a (global) duality function.

- ▶ It is not a local duality function.
- ▶ Given in [BCS+14]¹.
- ▶ No constructive proof. Can be checked once the formula guessed.

¹From duality to determinants for q-TASEP and ASEP (2014)

Global duality: Rank-two

Prop. 7.2 : Let ν be an infinite rank-one composition. Let μ be a rank-two composition with 1 and 2-particle positions given by,

$$\vec{x}(\mu) = (x_1 < \cdots < x_{m_1}),$$

$$\vec{y}(\mu) = (y_1 < \cdots < y_{m_2}),$$

Then the observable

$$H(\nu, \mu) = \prod_{x \in \vec{x}(\mu)} \left(\prod_{i \leq x} t^{\nu_i} \right) \cdot \prod_{y \in \vec{y}(\mu)} \left(\prod_{i \leq y} t^{\nu_i} \right) \cdot t^{-\chi(\vec{x}, \vec{y})}$$

is a (global) duality function, where

$$\chi(\vec{x}, \vec{y}) := \#\{(x_i, y_j) \in (\vec{x}, \vec{y}) : x_i > y_j\}.$$

- ▶ Direct check using the rank-one global duality Prop. 7.1.
- ▶ Proof by looking at blocks of particles + color-blind argument (rank-one composition).

Global duality: Arbitrary rank

Cor. 7.3 : Let ν be an infinite rank-one composition. Let μ be a rank- r composition with positions of j -particles ($1 \leq j \leq r$) labelled by,

$$\bar{x}^{(j)}(\mu) = (x_1^{(j)} < \cdots < x_{m_j}^{(j)}).$$

Then the observable

$$H(\nu, \mu) = \prod_{j=1}^r \left(\prod_{x \in \bar{x}^{(j)}} \left(\prod_{i \leq x} t^{\nu_i} \right) \right) t^{-\chi(\bar{x}^{(1)}, \dots, \bar{x}^{(r)})}$$

is a (global) duality function, where

$$\chi(\bar{x}^{(1)}, \dots, \bar{x}^{(r)}) := \#\left\{ (x, y) \in (\bar{x}^{(i)}, \bar{y}^{(j)}) : i < j, x > y \right\}.$$

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