

Maximin Strategies in Zero-Sum Games

- ▶ If a player adopts a maximin strategy, he believes that the opponent will always correctly predict their move.
- ▶ This means, the opponent will choose their best possible action based upon the player's act.
- ▶ In this case, player 1's expected payoff is:

$$R_{\text{maximin}}(d_i) = \min_j R(d_i, \delta_j)$$

- ▶ If this is the case, then player 2's payoff is:

$$-R_{\text{maximin}}(d_i) = \max_j -R(d_i, \delta_j)$$

- ▶ Hence $P1$ should play $d_{\text{maximin}}^* = \arg \max_{d_i} \min_j R(d_i, \delta_j)$.
- ▶ One could swap the two players to obtain a maximin strategy for player 2.

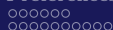
Example (RPS and Maximin)

- ▶ Let $M = (m_{ij})$ denote the payoff matrix for the RPS game.
- ▶ Then, $\min_j R(d_i, \delta_j) = \min_j m_{ij} = -1$ for all i .
- ▶ Thus any move is maximin for player 1.
- ▶ Player 1 expects to receive a payout of -1 whatever he does.
- ▶ If both players adopt a maximin view, then player 2 has the same expectation (by symmetry).
- ▶ How can we resolve this paradox?



What's Gone Wrong?

- ▶ The players aren't using all of the information available.
- ▶ They haven't used the fact that it is a zero sum game.
- ▶ They don't have compatible beliefs:
 - ▶ If P1 believes P2 can predict their move and P2 believes that P1 can predict their move then things inevitably go wrong.
 - ▶ It cannot be common knowledge that *both* players will adopt a maximin strategy!
- ▶ If a player really believes their opponent can predict their move then they can use randomization to make their action less predictable. . .



Mixed Strategies

- ▶ A *mixed strategy* for player 1 is a probability distribution over D .
- ▶ If a player has mixed strategy $\mathbf{x} = (x_1, \dots, x_n)$ then they will play move d_i with probability x_i .
- ▶ This can be achieved using a randomization device such as a spinner to select a move.
- ▶ A *pure strategy* is a mixed strategy in which exactly one of the x_i is non-zero (and is therefore equal to 1).
- ▶ A similar definition applies when considering player 2.



Expected Rewards and Mixed Strategies

What is player 1's expected reward if...

- ▶ Player 1 has mixed strategy \underline{x} and player 2 plays pure strategy δ_j ?
- ▶ Player 1 has pure strategy d_i and player 2 plays mixed strategy \underline{y} ?
- ▶ Player 1 has mixed strategy \underline{x} and player 2 has mixed strategy \underline{y} ?



Zero-Sum Games

In the first case, the uncertainty is player 1's own move, and his expectation is:

$$\sum_{i=1}^n x_i R(d_i, \delta_j)$$

In the second case, the uncertainty comes from player 2:

$$\sum_{j=1}^m y_j R(d_i, \delta_j)$$

Whilst both provide (independent) uncertainty in the third case:

$$\sum_{i=1}^n \sum_{j=1}^m x_i R(d_i, \delta_j) y_j = \underline{x}^T M \underline{y}$$



Maximin Revisited

- ▶ Player 1's maximin *mixed* strategy is the \underline{x} which maximises:

$$V_1 = \max_{\underline{x}} \min_{\underline{y}} \sum_i \sum_j x_i R(d_i, \delta_j) y_j$$

- ▶ Player 2's maximin *mixed* strategy is the \underline{y} which minimises:

$$\begin{aligned} & \max_{\underline{y}} \min_{\underline{x}} - \sum_i \sum_j x_i R(d_i, \delta_j) y_j \\ &= \min_{\underline{y}} \max_{\underline{x}} \sum_i \sum_j x_i R(d_i, \delta_j) y_j = V_2 \end{aligned}$$

What is the relationship between these two values?

Theorem (Fundamental Theorem of Zero Sum Two Player Games)

V_1 and V_2 as defined before satisfy:

$$V_1 = V_2$$

The unique value, $V = V_1 = V_2$ is known as the value of the game.

- ▶ The strategies \underline{x} and \underline{y} which achieve this value may not be unique.
- ▶ How can we find suitable strategies in general?

(Sketch of proof of theorem see later.)



Example (Maximin in a Simple Game)

- ▶ Consider a zero sum two player game with the following payoff matrix:

	δ_1	δ_2
d_1	1	3
d_2	4	2

- ▶ With a pure strategy maximin approach:
 - ▶ P1 plays d_2 expecting P2 to play δ_2 .
 - ▶ P2 plays δ_2 expecting P1 to play d_1 .
 - ▶ P1 expects to gain 2; P2 expects to lose 3.
 - ▶ This is not consistent.



Example

- ▶ Consider, instead, a mixed strategy maximin approach:
 - ▶ P1 plays a strategy $(x, 1 - x)$ and player 2 plays $(y, 1 - y)$.
 - ▶ Player 1's expected payoff is:

$$[x \quad 1 - x] \begin{bmatrix} 1 & 3 \\ 4 & 2 \end{bmatrix} \begin{bmatrix} y \\ 1 - y \end{bmatrix} = -4xy + x + 2y + 2, \text{ use algebra...}$$

$$= -4\left(x - \frac{1}{2}\right)\left(y - \frac{1}{4}\right) + \frac{5}{2}$$

- ▶ Player 1 seeks to maximise this for the worst possible y .
- ▶ As the 2nd player can control the sign of the first term, his optimal strategy is to make it vanish by choosing $x = \frac{1}{2}$.
- ▶ Similarly, the 2nd player wants to prevent the first player from exploiting the first term and chooses $y = \frac{1}{4}$.
- ▶ Now, the expected reward for the first player is, consistently, 2.5 as both expect the same maximin strategies to be played.
- ▶ *Both* players have a higher expected return than they would playing pure strategies.



How do we determine maximin mixed strategies?

- ▶ We need a general strategy for determining strategies \underline{x}^* and \underline{y}^* which achieve the common maximin return for player 1.
- ▶ It's straightforward (if possibly tedious) to calculate, for payoff matrix M the expected return for player 1 as a function of the strategies:

$$V(\underline{x}, \underline{y}) = \underline{x}^T M \underline{y}$$

- ▶ We then seek to obtain $\underline{x}^*, \underline{y}^*$ such that:

$$V(\underline{x}^*, \underline{y}^*) = \max_{\underline{x}} \min_{\underline{y}} V(\underline{x}, \underline{y})$$

- ▶ In general, this is a problem which can be efficiently addressed by linear programming.
- ▶ If one player has only two possible decisions, however, a simple graphical method can be employed. *(Only 1 parameter!)*



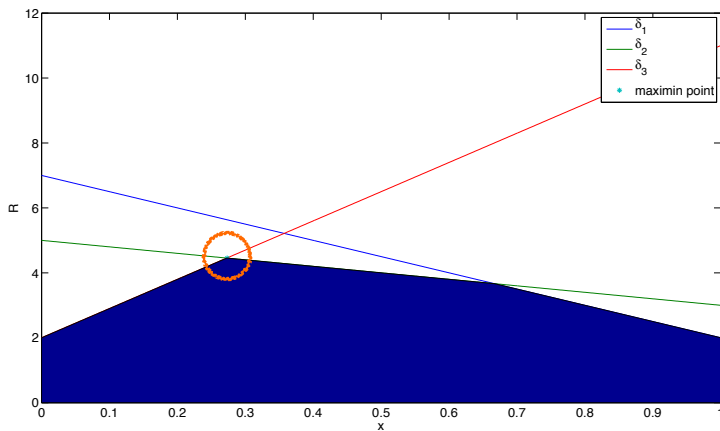
Graphical Solution, Part 1: Player 1's approach

- ▶ Consider a two player zero sum game with payoff matrix:

$$M = \begin{bmatrix} 2 & 3 & 11 \\ 7 & 5 & 2 \end{bmatrix}$$

- ▶ Consider a mixed strategy $(x, 1 - x)$ for player 1.
- ▶ For the three pure strategies available to player 2, player 1 has expected reward:
 - ▶ $\delta_1 : 2x + 7(1 - x) = 7 - 5x$
 - ▶ $\delta_2 : 3x + 5(1 - x) = 5 - 2x$
 - ▶ $\delta_3 : 11x + 2(1 - x) = 2 + 9x$
- ▶ For each value of x , the worst case response of player 2 is the one for which the expected reward of player 1 is minimised.
- ▶ Plotting the three lines as a function of x ...

Zero-Sum Games



Zero-Sum Games

- ▶ The maximin response maximises the return in the worst case.
- ▶ In terms of our graph, this means we choose x to maximise the distance between the lowest of the lines and the ordinate axis.
- ▶ This is at the point where the lines associated with δ_2 and δ_3 intersect, at x^* which solves:

$$5 - 2x = 2 + 9x$$

$$11x = 3 \Rightarrow x^* = 3/11$$

- ▶ Hence player 1's maximin mixed strategy is $(3/11, 8/11)$.
- ▶ Playing this, his expected return is:

$$V_1 = 2 + 9 \times 3/11 = 49/11 = 5 - 2 \times 3/11 = 49/11$$

Graphical Solution, Part 2: Player 2's approach

- ▶ Player 2 only needs to consider the moves which optimally oppose player 1's maximin strategy (δ_2 and δ_3).
- ▶ They may consider a mixed strategy $(0, y, 1 - y)$.
- ▶ By the fundamental theorem, player 2's maximin strategy leads to the same expected payoff for player 1 as his own maximin strategy:

$$V_2 = V_1 = 49/11. \quad \begin{array}{r} 2 \quad 3 \quad 11 \\ 7 \quad 5 \quad 2 \end{array}$$

- ▶ They should play y^* to solve:

$$V_2 = 3y + 11(1 - y) = 49/11$$

$$8y = (121 - 49)/11 = 72/11 \Rightarrow y^* = 9/11$$

- ▶ Leading to a mixed strategy $(0, 9/11, 2/11)$.

Example (Spy Game)

- ▶ A spy has escaped and must choose to flee down a *river* or through a *forest*. Their guard must choose to chase them using a *helicopter*, a pack of *dogs* or a *jeep*.
- ▶ They agree that the probabilities of escape are as given in this payoff matrix:

	H	D	J
R	0.1	0.8	0.4
F	0.9	0.1	0.6

- ▶ Both players wish to adopt maximin strategies.

Example

- ▶ The spy plays strategy $(x, 1 - x)$: with probability x they escape via the river; with probability $1 - x$ they run through the forest.
- ▶ For given x , their probabilities of escaping for each of the guard's possible actions are:

$$p_H = 0.1x + 0.9(1 - x) \qquad p_D = 0.8x + 0.1(1 - x)$$

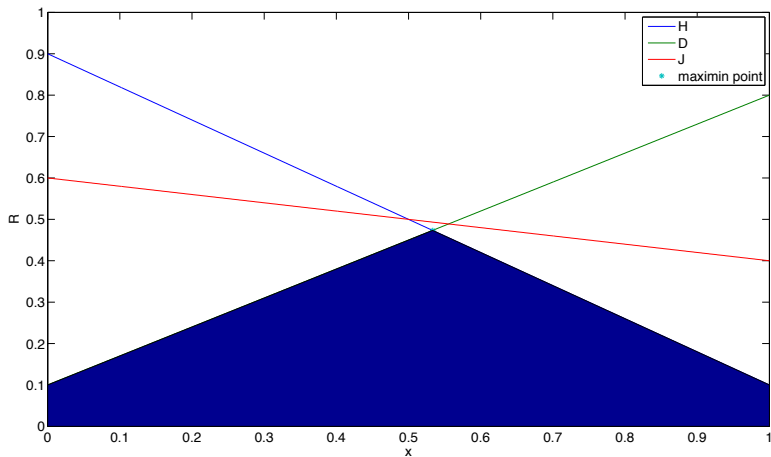
$$= \frac{9 - 8x}{10} \qquad = \frac{1 + 7x}{10}$$

$$p_J = 0.4x + 0.6(1 - x)$$

$$= \frac{6 - 2x}{10}$$

- ▶ Plotting these three lines as a function of x we obtain the following figure:

Zero-Sum Games



Example

- ▶ The maximin solution is the intersection of the lines for strategies D and H .
- ▶ This occurs at the solution, x^* of:

$$p_H = p_D \Rightarrow 9 - 8x = 1 + 7x$$

$$8 = 15x \quad \Rightarrow x^* = 8/15$$

- ▶ The value of the game is: $V = V_1 = \frac{9-8x^*}{10} = 71/150$

Example

- ▶ By the fundamental theorem of zero sum two player games, the guard needs to consider only H and D .
- ▶ Otherwise the spy's chance of escape will be better than V_1 if he plays his own maximin strategy.
- ▶ Consider a strategy $(y, 1 - y, 0)$.
- ▶ By the same theorem, $V_2 = V = V_1$, so:

$$V_2 = 0.1y^* + 0.8(1 - y^*) = 71/150$$

$$8 - 7y^* = 71/15$$

$$y^* = 7/15$$

On Zero Sum Two Player Games

- ▶ The “fundamental theorem” does not generalise to games of more than two players.
- ▶ The “fundamental theorem” does not generalise to non-zero-sum games.
- ▶ Games with an element of co-operation are much more interesting.