

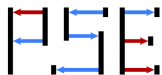
Ancestral lines under selection and recombination

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joint work with Frederic Alberti

1. The selection-recombination equation (forward in time)
2. Its solution via genealogical thinking
3. The dual process

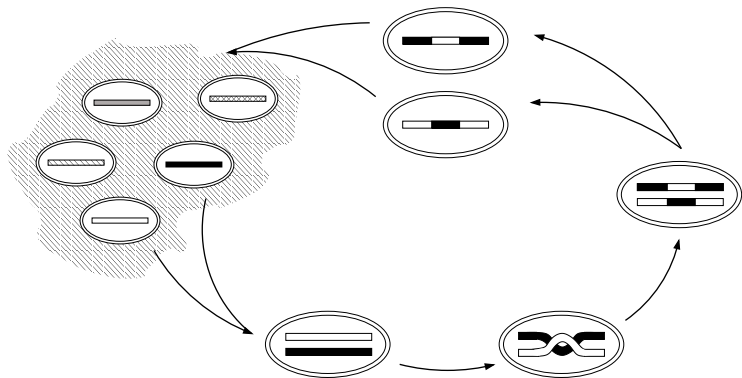


PROBABILISTIC STRUCTURES
IN EVOLUTION

DFG SPP 1590



Recombination



Sequences and selection

individual: **sequence** of n sites, $S = \{1, \dots, n\}$

types: $x := (x_1, \dots, x_n) \in \prod_{i \in S} X_i =: X, X_i = \{0, 1\}$

marginal types: $x_U := (x_i)_{i \in U}, U \subseteq S$

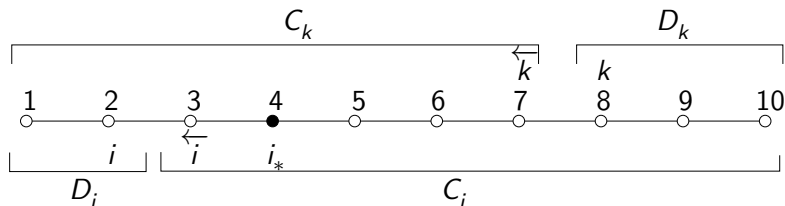
reproduction and selection:

- single **selected site**: $i_* \in S$ (all others 'neutral')
- types x with $x_{i_*} = 1$: reproduce at rate 1 ('bad, less fit')
- types x with $x_{i_*} = 0$: reproduce at rate $1 + s$ ('good, fit')

Ordering the sites

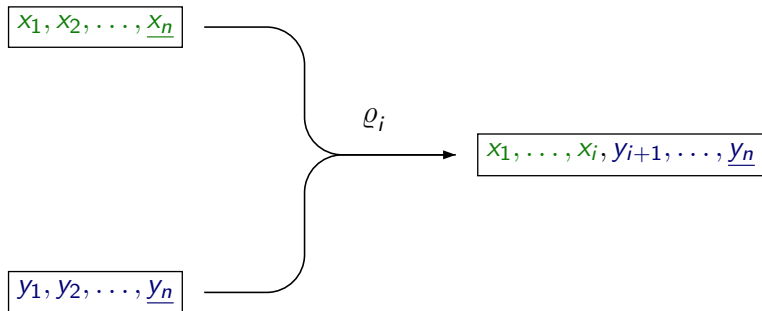
partial order on S :

- $i \preceq j$ if either $i_* \leq i \leq j$ or $i_* \geq i \geq j$
- $i \prec j$ if $i \preceq j$ and $i \neq j$
- $\leftarrow i$ maximal element that precedes i (predecessor)
- $D_i := \{j \in S \mid i \preceq j\}$ (i -tail), $C_i := S \setminus D_i$ (i -head), $i \in S \setminus i_*$

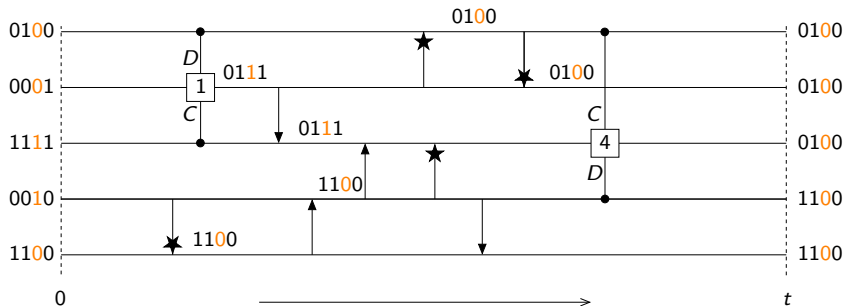


- $\rightsquigarrow i_* \in C_i$
- (single-crossover) recombination between C_i and D_i , $i \in S \setminus i_*$ (partitions S into C_i and D_i) at rate ρ_i

Sequences and recombination



Moran model with selection and recombination



$N(= 5)$ individuals, $n = 4$ sites, selected site $i_* = 3$

- neutral reproduction, rate 1 (for all individuals)
- ★ selective reproduction, rate s (for $(**0*)$ ind.'s)
- $\frac{C}{\boxed{i}} \frac{D}{\quad}$ reco at site i , rate ρ_i , $i \in S \setminus i_*$ (for all individuals)

Selection-recombination equation

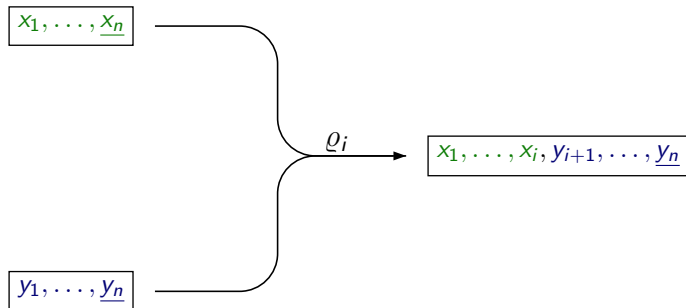
- $Z_t^N = (Z_t^N(x))_{x \in X}$ counting measure on X
- $Z_t^N(x)$ # individuals of type x at time t , $x \in X$
- deterministic limit (dynamical law of large numbers):
 $N \rightarrow \infty$ w/o rescaling of parameters or time \rightsquigarrow

$$\frac{1}{N}(Z_t^N)_{t \geq 0} \rightarrow (\omega_t)_{t \geq 0}$$

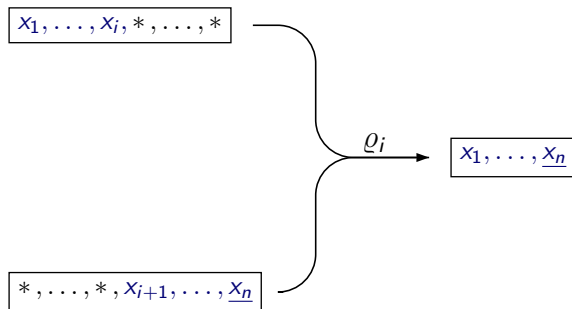
- $\omega_t = (\omega_t(x))_{x \in X}$ probability measure on X
solves selection-recombination equation (SRE)

$$\dot{\omega}_t(x) = s[(1 - x_{i_*}) - \omega_t(*, 0, *)] \omega_t(x) + \sum_{i \in S \setminus i_*} \varrho_i [\omega_t(x_{C_i}, *) \omega_t(*, x_{D_i}) - \omega_t(x)], \quad x \in X$$

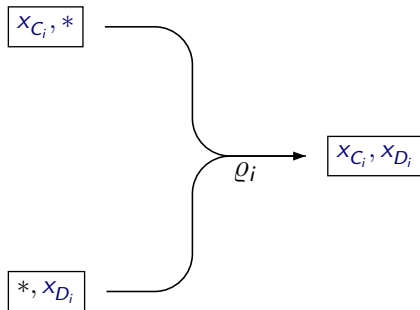
Recombination



Recombination



Recombination



Selection-recombination equation

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- $\omega_t = (\omega_t(x))_{x \in X}$ probability measure on X
solves **selection-recombination equation (SRE)**

$$\begin{aligned} \dot{\omega}_t(x) = & s[(1 - x_{i_*}) - \omega_t(*, 0, *)] \omega_t(x) \\ & + \sum_{i \in S \setminus i_*} \varrho_i [\omega_t(x_{C_i}, *) \omega_t(*, x_{D_i}) - \omega_t(x)], \quad x \in X \end{aligned}$$

Recombinators

- canonical projection: for $\emptyset \neq U \subseteq S$,

$$\pi_U : X \rightarrow \prod_{i \in U} X_i = X_U, \quad \pi_U(x) = (x_i)_{i \in U} = x_U$$

- marginal measure wrt sites in U : for $\nu \in \mathcal{P}(X)$,

$$\pi_U.\nu = \nu \circ \pi_U^{-1} =: \nu^U$$

type distribution of sites in U

for $x_U \in X_U$: $\nu^U(x_U) = \nu(x_U, *)$

- recombinator: for $i \in S \setminus i_*$,

$$\mathcal{P}(X) \longrightarrow \mathcal{P}(X)$$

$$R_i(\nu) := \nu^{C_i} \otimes \nu^{D_i}$$

$$(R_i(\nu))(x) = \nu(x_{C_i}, *) \cdot \nu(*, x_{D_i})$$

distribution of sequences randomly pieced together from i -heads and i -tails

Selection-recombination equation

$$\dot{\omega}_t(x) = s[(1 - x_{i_*}) - \omega_t^{\{i_*\}}(0)] \omega_t(x) + \sum_{i \in S \setminus i_*} \varrho_i [\omega_t(x_{C_i}, *) \omega_t(*, x_{D_i}) - \omega_t(x)], \quad x \in X$$

or

$$\dot{\omega}_t = s(F - f(\omega_t))\omega_t + \sum_{i \in S \setminus i_*} \varrho_i (R_i - \mathbb{1})\omega_t$$

with

$$F\omega(x) = (1 - x_{i_*})\omega(x)$$

and

$$f(\omega) := \omega^{\{i_*\}}(0)$$

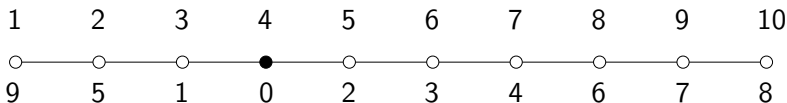
(proportion of fit individuals = proportion of individuals with 0 at selected site)

Research context

- **hitchhiking**: neutral x_i ($i \neq i_*$) selectively favoured if linked with $x_{i_*} = 0$
- recombination and selection on same time scale (unlike Cuthbertson, Etheridge, Yu 2012; Pfaffelhuber, Studeny 2007; Bossert, Pfaffelhuber 2017)
- **Stephan, Song, Langley 2006**: approximate solution forward in time ($n = 3$, wild approximations, technical result, not generalisable)

Recursive solution

- $(i_k)_{0 \leq k < n}$ non-decreasing permutation of S (sensu \Leftarrow)
- $i_0 = i_*$
- $C^{(k)} = C_{i_k}, D^{(k)} = D_{i_k}, \varrho^{(k)} := \varrho_{i_k}$



hierarchy of solutions:

- $(\omega_t^{(0)})_{t \geq 0}$ solution of SRE with $\varrho^{(\ell)} = 0$ for all ℓ
(pure selection equation)
- $(\omega_t^{(k)})_{t \geq 0}$ solution of **SRE truncated at k**
(with $\varrho^{(\ell)} = 0$ for all $\ell > k$), $1 \leq k < n$

Recursive solution

Theorem

The solution of the pure selection equation is given by

$$\omega_t^{(0)} = \frac{e^{stF} \omega_0}{e^{st} f(\omega_0) + 1 - f(\omega_0)},$$

and the solutions $\omega_t^{(k)}$ can be computed via the recursion

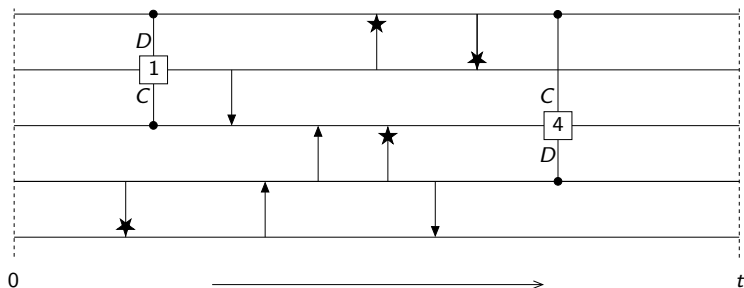
$$\omega_t^{(k)} = e^{-\varrho^{(k)} t} \omega_t^{(k-1)} + \pi_{C^{(k)}} \cdot \omega_t^{(k-1)} \otimes \int_0^t \varrho^{(k)} e^{-\varrho^{(k)} \tau} \pi_{D^{(k)}} \cdot \omega_\tau^{(k-1)} d\tau.$$

analytical proof ✓

geneological content?

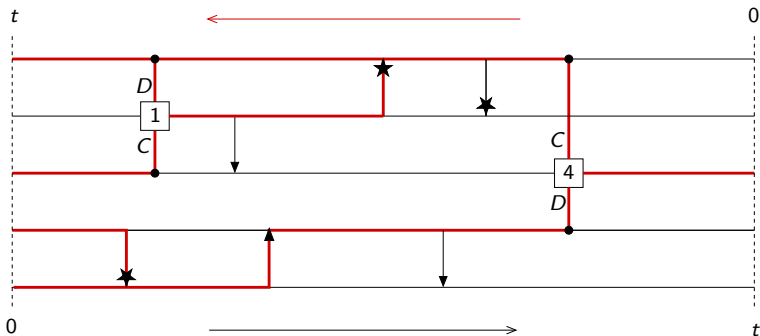
Genealogical approach

Moran model:



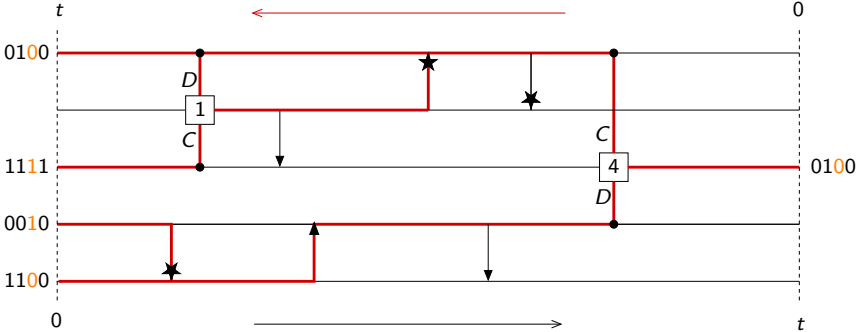
Genealogical approach

Ancestral selection-recombination graph (ASRG):

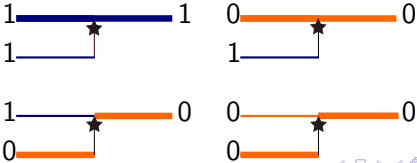


Genealogical approach

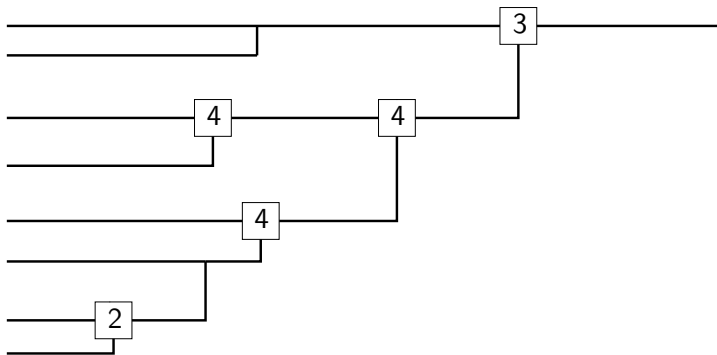
Ancestral selection-recombination graph (ASRG):



pecking order:

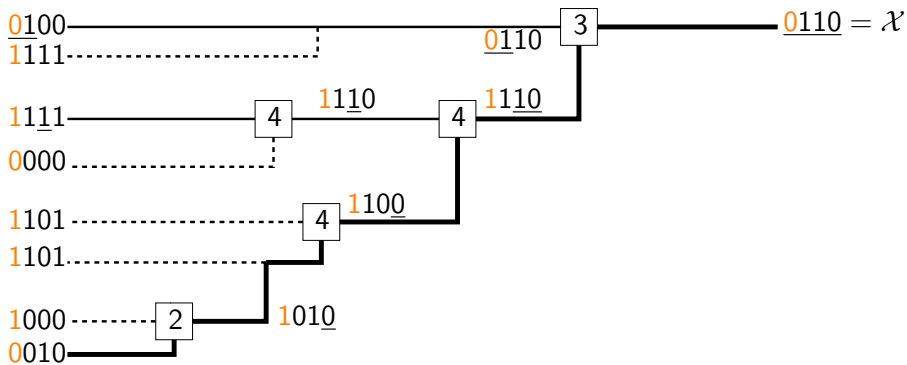


$N \rightarrow \infty$ limit of ASRG



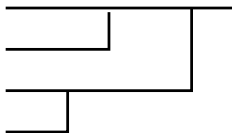
- no coalescence \rightsquigarrow branches conditionally independent
- every line branches at rate s , splits at rates ϱ_j , $i \in S \setminus i_*$ (here $i_* = 1$)
- pecking order and 'heads up' order

$N \rightarrow \infty$ limit of ASRG



- run ASRG backward
- sample types according to ω_0 (iid)
- propagate forward by pecking and glueing
- $\mathcal{X} \sim \omega_t$
- ancestral lines

Selection only: $\omega_t^{(0)}$



- $\varrho^{(\ell)} = 0$, for all $\ell \rightsquigarrow$ pure branching
 \rightsquigarrow Yule process at rate s ('ASG')

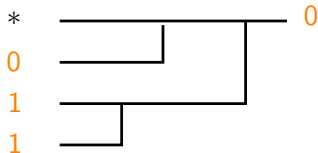
- $(K_t)_{t \geq 0}$ line-counting process

- $\mathbb{P}(K_t = k) = e^{-st}(1 - e^{-st})^{k-1}$

- $\mathbb{E}(K_t) = e^{st}$

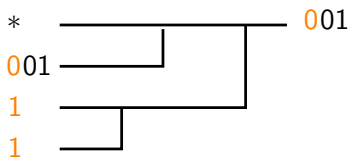
- $f(\omega_t) = 1 - \mathbb{E}[(1 - f(\omega_0))^{K_t}] = \frac{e^{st}f(\omega_0)}{e^{st}f(\omega_0) + 1 - f(\omega_0)}$

Selection only: $\omega_t^{(0)}$



$$\begin{aligned}\omega_t^{(0)} &= \mathcal{L}(\mathcal{X}) = \mathbb{E}[(1 - f(\omega_0))^{K_t}] \omega_0(\cdot \mid \tilde{\mathcal{X}}_{i_*} = 1) \\ &\quad + (1 - \mathbb{E}[(1 - f(\omega_0))^{K_t}]) \omega_0(\cdot \mid \tilde{\mathcal{X}}_{i_*} = 0) \\ &= \frac{e^{stF} \omega_0}{e^{st} f(\omega_0) + 1 - f(\omega_0)}\end{aligned}$$

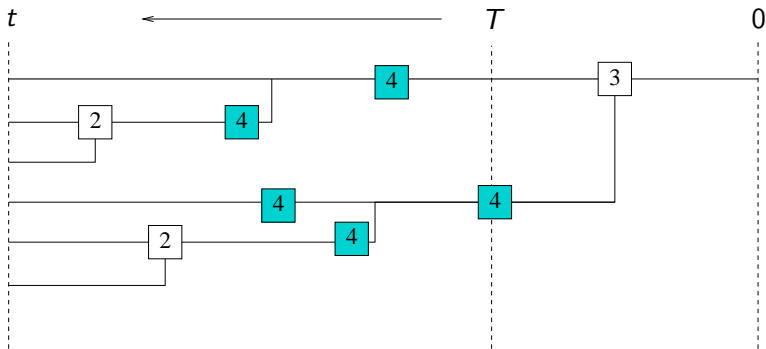
Selection only: $\omega_t^{(0)}$



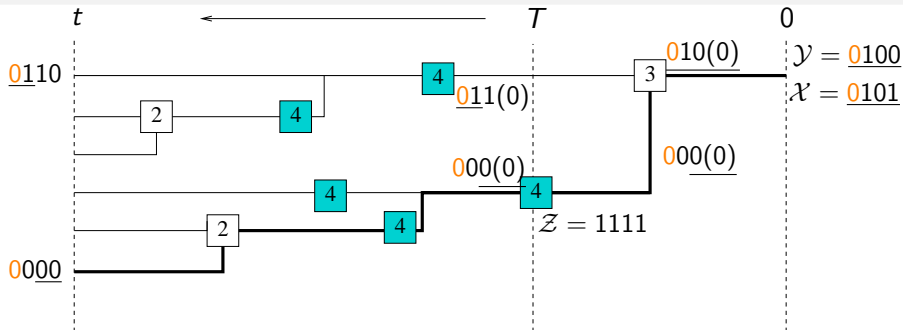
$$\begin{aligned}
 \omega_t^{(0)} = \mathcal{L}(\mathcal{X}) &= \mathbb{E}[(1 - f(\omega_0))^{K_t}] \omega_0(\cdot \mid \tilde{\mathcal{X}}_{i_*} = 1) \\
 &\quad + (1 - \mathbb{E}[(1 - f(\omega_0))^{K_t}]) \omega_0(\cdot \mid \tilde{\mathcal{X}}_{i_*} = 0) \\
 &= \frac{e^{stF} \omega_0}{e^{st} f(\omega_0) + 1 - f(\omega_0)}
 \end{aligned}$$

Truncated and collapsed ASRG

- $\text{ASRG}^{(k)}$: **ASRG truncated at k** ($\rho^{(\ell)} = 0$ for $\ell > k$)
- $\text{cASRG}^{(k)}$: **collapsed $\text{ASRG}^{(k)}$** ($\text{ASRG}^{(k-1)}$ decorated with i_k -squares at rate $\rho^{(k)}$ on every line)
- $\rightsquigarrow \text{ASRG}_t^{(k)} = \text{cASRG}_t^{(k)}$ with independent $\text{ASRG}^{(k)}$ attached to every i_k -square in a $\text{cASRG}_t^{(k)}$ for the remaining time

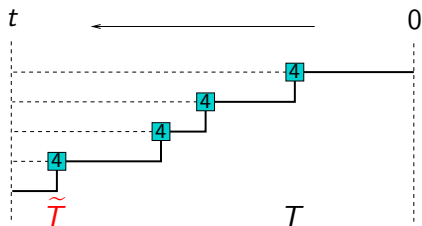


Induction step: $\omega_t^{(k-1)} \rightarrow \omega_t^{(k)}$



- \mathcal{Y} type at root of $\text{ASRG}_t^{(k-1)}$ ($\sim \omega_t^{(k-1)}$)
= type at root of $\text{cASRG}_t^{(k)}$ when ignoring the i_k -squares
- T time of rightmost i_k -square on line ancestral to $D^{(k)}$ in $\text{ASRG}_t^{(k-1)}$; if no such square, $T := \Delta$
- \mathcal{Z} type at root of $\text{ASRG}_{t-T}^{(k)}$ ($\sim \omega_{t-T}^{(k)}$)
- $\mathcal{X} = \mathbb{1}_{T=\Delta} \mathcal{Y} + \mathbb{1}_{T \neq \Delta} (\pi_{C^{(k)}}(\mathcal{Y}), \pi_{D^{(k)}}(\mathcal{Z}))$
type at root of $\text{ASRG}_t^{(k)}$ ($\sim \omega_t^{(k)}$)

Induction step: $\omega_t^{(k-1)} \rightarrow \omega_t^{(k)}$



$$\mathcal{X} = \mathbb{1}_{T=\Delta} \mathcal{Y} + \mathbb{1}_{T \neq \Delta} (\pi_{C^{(k)}}(\mathcal{Y}), \pi_{D^{(k)}}(\tilde{\mathcal{Z}}))$$

$$\rightsquigarrow \dots \rightsquigarrow \mathcal{X} = \mathbb{1}_{\tilde{T}=\Delta} \mathcal{Y} + \mathbb{1}_{\tilde{T} \neq \Delta} (\pi_{C^{(k)}}(\mathcal{Y}), \pi_{D^{(k)}}(\tilde{\mathcal{Z}}))$$

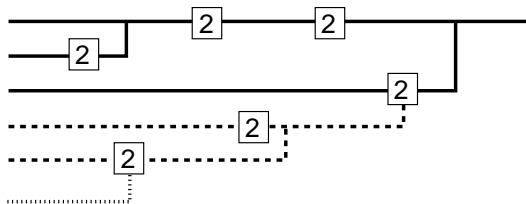
$\tilde{\mathcal{Z}}$ type at root of ASRG $_{t-\tilde{T}}^{(k-1)}$ ($\sim \omega_{t-\tilde{T}}^{(k-1)}$)

$$\begin{aligned} \omega_t^{(k)} &= \mathcal{L}(\mathcal{X}) = \mathbb{E}[\mathcal{L}(\mathcal{X} \mid \tilde{T})] \\ &= e^{-\rho^{(k)}t} \omega_t^{(k-1)} + \pi_{C^{(k)}} \cdot \omega_t^{(k-1)} \otimes \pi_{D^{(k)}} \cdot \int_0^t \rho^{(k)} e^{-\rho^{(k)}\tau} \omega_\tau^{(k-1)} d\tau \end{aligned}$$

(via symmetry of Poisson process)

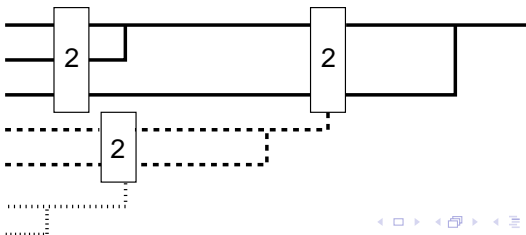
Synchronising the ASRG

- expand $cASRG^{(1)}$ along ancestral line

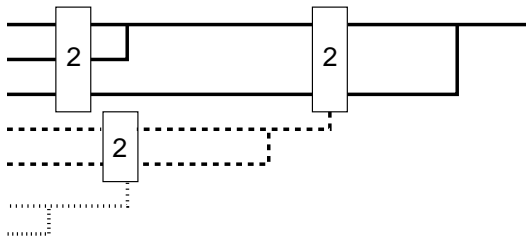


until new copy has no i_1 -square on ancestral line

- only i_1 -squares on the ancestral line of $D^{(1)}$ relevant
 \rightsquigarrow synchronise i_1 -squares across lines:



Synchronised ASRG



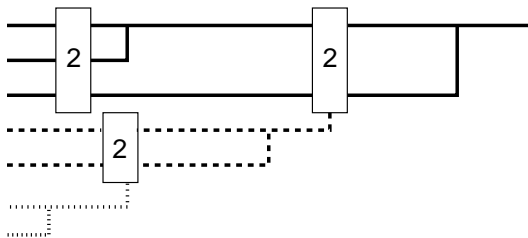
only i_1 -squares on the ancestral line of $D^{(1)}$ relevant

↪ identical distribution across lines matters,
but independence does not

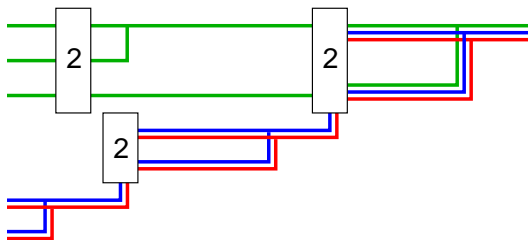
↪ ASG decorated by i_1 -bars laid down at rate $\rho^{(1)}$

↪ new tail attached to whichever 'head-line' is ancestral

Labelling and pruning

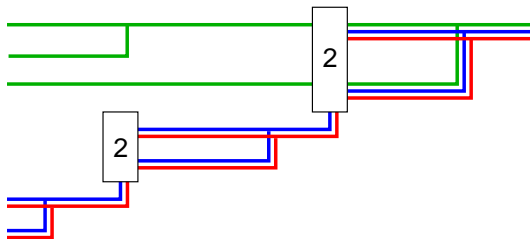


Labelling and pruning



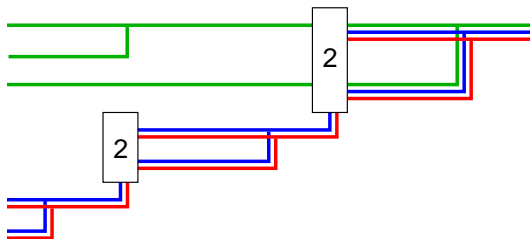
- label every line with the set A of sites to which it is potentially ancestral $\rightsquigarrow A$ splits into $A \cap C_{i_1}$ and $A \cap D_{i_1}$ at i_1 -bar

Labelling and pruning

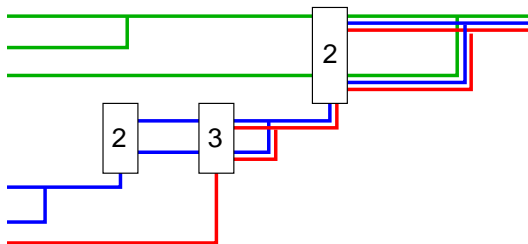


- label every line with the set A of sites to which it is potentially ancestral $\rightsquigarrow A$ splits into $A \cap C_{i_1}$ and $A \cap D_{i_1}$ at i_1 -bar
- delete i_1 -bar when $A \subseteq C_{i_1}$

Iteration

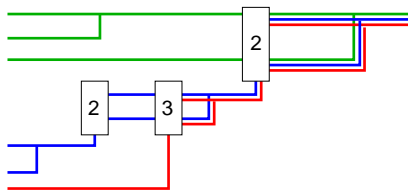


Iteration



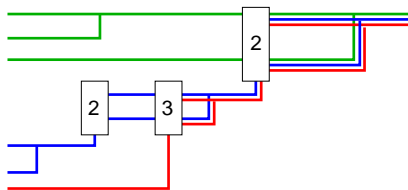
add i_2 -bars at rate ρ_{i_2} on every ASG, label, prune ...
... and so forth

Dual process: Weighted partitioning process



- $(\Sigma_t, V_t)_{t \geq 0}$ **weighted partitioning process**
- $(\Sigma_t)_{t \geq 0}$ partitioning process ('ARG')
on set of interval (or ordered) partitions of S
- V_t assigns an integer to every part of Σ_t
(the number of lines in its ASG)

Dual process: Weighted partitioning process



if $(\Sigma_t, V_t) = (\mathcal{A}, v)$,

$$\mathcal{A} = \{A_1, \dots, A_m\}, v = (v_A)_{A \in \mathcal{A}}:$$

- $\mathcal{A} \rightarrow \{A_1, \dots, A_j \cap C_i, A_j \cap D_i, \dots, A_m\}$,
 $v \rightarrow (v_{A_1}, \dots, v_{A_j}, 1, v_{A_{j+1}}, \dots, v_{A_m})$, at rate ϱ_i , $i \in S \setminus i_*$
- independently of Σ_t :
 - $v_A \rightarrow v_A + 1$ at rate sv_A for every $A \in \mathcal{A}$ ('Yule')
 - $v_A \rightarrow 1$ at rate $\sum_{\ell \preccurlyeq \min(A)} \varrho_\ell$ for every $A \in \mathcal{A}$
(**'reset'** at any split between A and i_*)

Duality

duality function:

$$H((\mathcal{A}, \nu), \omega) := \bigotimes_{A \in \mathcal{A}} \left[(1 - (1 - f(\omega))^{\nu_A}) \omega^A(\cdot \mid \tilde{\mathcal{X}}_{i_*} = 0) \right. \\ \left. + (1 - f(\omega))^{\nu_A} \omega^A(\cdot \mid \tilde{\mathcal{X}}_{i_*} = 1) \right]$$

Theorem

$$\mathbb{E}_{\omega_0} [H((\mathcal{A}, \nu), \omega_t)] = \mathbb{E}_{(\mathcal{A}, \nu)} [H((\Sigma_t, V_t), \omega_0)]$$

in particular,

$$\omega_t = \mathbb{E}_{(\{S\}, 1)} [H((\Sigma_t, V_t), \omega_0)]$$