

The Gentzen rules for propositional logic (\mathcal{G}_{PL})

\wedge rules

$$\frac{\varphi \quad \psi}{\varphi \wedge \psi} \wedge I \qquad \frac{\psi \quad \varphi}{\varphi \wedge \psi} \wedge I \qquad \frac{\varphi \wedge \psi}{\varphi} \wedge E \qquad \frac{\varphi \wedge \psi}{\psi} \wedge E$$

\vee rules

$$\frac{\varphi}{\varphi \vee \psi} \vee I \qquad \frac{\psi}{\varphi \vee \psi} \vee I \qquad \frac{\varphi \vee \psi \quad \begin{array}{c} [\varphi] \\ \vdots \\ \chi \end{array} \quad \begin{array}{c} [\psi] \\ \vdots \\ \chi \end{array}}{\chi} \vee E$$

Note: $\vee E$ discharges *both* the hypotheses φ and ψ .

\rightarrow rules

$$\frac{\begin{array}{c} [\varphi] \\ \vdots \\ \psi \end{array}}{\varphi \rightarrow \psi} \rightarrow I \qquad \frac{\varphi \rightarrow \psi \quad \varphi}{\psi} \rightarrow E$$

Note: Although in the form given, $\rightarrow I$ discharges the hypothesis φ , it also admits the form $\frac{[\psi]}{\varphi \rightarrow \psi} \rightarrow I$ in which ψ is *not* discharged.

\neg rules

$$\frac{\begin{array}{c} [\varphi] \\ \vdots \\ \perp \end{array}}{\neg \varphi} \neg I \qquad \frac{\varphi \quad \neg \varphi}{\perp} \neg E$$

\perp rules

$$\frac{\perp}{\varphi} \perp E \qquad \frac{\begin{array}{c} [\neg \varphi] \\ \vdots \\ \perp \end{array}}{\varphi} \text{RAA}$$

Note: RAA discharges the hypothesis $[\varphi]$.

\leftrightarrow rules

$$\frac{\begin{array}{c} [\varphi] \\ \vdots \\ \psi \end{array} \quad \begin{array}{c} [\psi] \\ \vdots \\ \varphi \end{array}}{\varphi \leftrightarrow \psi} \leftrightarrow I \qquad \frac{\varphi \leftrightarrow \psi \quad \varphi}{\psi} \leftrightarrow E \qquad \frac{\psi \leftrightarrow \psi \quad \varphi}{\varphi} \leftrightarrow E$$

Note: $\leftrightarrow I$ discharges both the hypotheses φ and ψ .

The Gentzen rules for first-order logic (\mathcal{G}_{FOL})

\mathcal{G}_{FOL} consists of the propositional rules of \mathcal{G}_{PL} together with the following first-order rules:

\forall rules

$$\frac{[\Gamma] \quad \vdots \quad \varphi(x)}{\forall x \varphi(x)} \forall I \quad \text{where } x \notin \text{FV}(\Gamma) \qquad \frac{\forall x \varphi(x)}{\varphi(x)[t/x]} \forall E \quad \text{where } t \text{ is free for } x \text{ in } \varphi(x)$$

\exists rules

$$\frac{\varphi(t)}{\exists x \varphi(x)} \exists I \quad \text{where } t \text{ is free for } x \text{ in } \varphi(x)$$

$$\frac{\exists x \varphi(x) \quad \begin{array}{c} [\varphi(x)] \\ \vdots \\ \psi \end{array}}{\psi} \exists E \quad \text{if } x \notin \text{FV}(\psi) \text{ and } x \text{ is also not free in any hypothesis on which the subderivation of } \psi \text{ depends other than } \varphi(x)$$

Identity axioms

$$\frac{}{x = x} RI_1 \qquad \frac{x = y}{y = x} RI_2 \qquad \frac{x = y \quad y = z}{x = z} RI_2$$

$$\frac{x_1 = y_1, \dots, x_n = y_n}{t(x_1, \dots, x_n) = t(y_1, \dots, y_n)} RI_4$$

$$\frac{x_1 = y_1, \dots, x_n = y_n \quad \varphi(x_1, \dots, x_n)}{\varphi(y_1, \dots, y_n)} RI_5$$

In RI_1, \dots, RI_5 it is assumed that the variables y_1, \dots, y_n are free for x_1, \dots, x_n in φ .