

All rules of Fitch have this feature.

Paradoxes of material implication

Adams (Inquiry 1965): “The following ... ‘inferences’ are all ones which, when symbolized in the ordinary way in the propositional calculus and analysed truth-functionally, are valid”:

1. John will arrive on the 10 o'clock plane. Therefore, if John does not arrive on the 10 o'clock plane, he will arrive on the 11 o'clock plane. [because Q entails $(P \rightarrow Q)$]
2. John will arrive on the 10 o'clock plane. Therefore, if John misses his plane in New York, he will arrive on the 10 o'clock plane. [because Q entails $(P \rightarrow Q)$]
3. If Brown wins the election, Smith will retire to private life. Therefore, if Smith dies before the election and Brown wins it, Smith will retire to private life. [because $(P \rightarrow R)$ entails $((P \wedge Q) \rightarrow R)$]
4. If Brown wins the election, Smith will retire to private life. If Smith dies before the election, Brown will win it. Therefore, if Smith dies before the election, then he will retire into private life. [because $(P \rightarrow Q)$ and $(Q \rightarrow R)$ entails $(P \rightarrow R)$]
5. It is not the case that if John passes history, he will graduate. Therefore, John will pass history. [because $\neg(P \rightarrow Q)$ entails P]
6. If you throw both switch S and switch T, the motor will start. Therefore, either if you throw switch S the motor will start, or if you throw switch T the motor will start. [because $((P \wedge Q) \rightarrow R)$ entails both $(P \rightarrow R)$ and $(Q \rightarrow R)$]

From rules to truth tables

On the assumption that \wedge Intro and \wedge Elim are sound, there is only one possible truth table for \wedge .

Material implication (\rightarrow) again

1.		
├		
└	2. \perp	
└	├	
└	3. \perp	\perp Elim: 2
4. $\perp \rightarrow \perp$		\rightarrow Intro: 2-3

1.		
├		
└	2. \perp	
└	├	
└	3. $P \vee \neg P$	\perp Elim: 2
4. $\perp \rightarrow (P \vee \neg P)$		\rightarrow Intro: 2-3

1.		
├		
└	2. $P \vee \neg P$	
└	├	
└	3. $P \vee \neg P$	\perp Elim: 2
4. $P \vee \neg P \rightarrow (P \vee \neg P)$		\rightarrow Intro: 2-3

Truth functional completeness

Example: for any truth function, we can find a logically equivalent one using only \neg , \vee , \wedge .

DeMorgan's Laws

$$\neg P \wedge \neg Q \equiv \neg(P \vee Q)$$

$$\neg P \vee \neg Q \equiv \neg(P \wedge Q)$$

Substitution Theorem

$$\text{If } A \equiv B \text{ then } S(A) \equiv S(B)$$

Exercises

Little bit more than usual because we have no lectures or seminars in reading week.

7.1–7.8

7.12.1–7.12,19 odd numbers only

FOR 7.12.* ALSO SUBMIT YOUR ANSWERS ON PAPER/BY EMAIL.

8.1

8.3, 8.5, 8.7, 8.9 odd nos only

8.12

8.16

8.17

8.18–8.25

8.26–30

FOR 8.26–30 YOU DO NOT HAVE PERMISSION TO USE TAUT CON TO PROVE $P \vee \neg P$